Dealing with Constraints via Random Permutation

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Motivation

Optimization for Large-scale Problems

- How to solve large-scale constrained problems?
- Popular idea: solve small subproblems
 - CD (Coordinate Descent)-type: min $f(x_1, \ldots, x_N)$. $x_1 \rightarrow x_2 \rightarrow \cdots \rightarrow x_N$

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 - SGD (Stochastic Gradient Descent): min $\sum_i f_i(x)$. $f_1 \rightarrow f_2 \rightarrow \cdots \rightarrow f_N$
- Widely (and wildly) used in practice: deep learning, glmnet for LASSO, libsvm for SVM, recommendation systems, EM

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- Widely (and wildly) used in practice: deep learning, glmnet for LASSO, libsvm for SVM, recommendation systems, EM
- Compared to other ideas, e.g., first-order methods and sketching:
 - Similar cheap iteration idea
 - "Orthogonal" to other ideas, so can combine

Go Beyond Unconstrained Optimization

- Many problems have (linear) constraints
- Classical convex optimization, e.g., linear programming.
 - · Combinatorial optimization (this workshop)
 - Operations research problems
- Machine learning applications, e.g., structured sparsity and deep learning
- Can we apply the *decomposition* idea? Turn out to be tricky!
- Algorithm: CD + multiplier —> ADMM (Alternating Direction Method of Multipliers)

Multi-block ADMM

• Consider a linearly constrained problem

$$\min_{x \in \mathbb{R}^{N}} f(x_{1}, x_{2}, \dots, x_{n})$$

s.t. $Ax \triangleq A_{1}x_{1} + \dots + A_{n}x_{n} = b,$ (1)
 $x_{j} \in \mathcal{X}_{j} \subseteq \mathbb{R}^{d_{j}}, j = 1, \dots, n.$

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Augmented Lagrangian function:

$$L_{\gamma}(x_1,\ldots,x_n;\lambda) = f(x) - \langle \lambda, \sum_i A_i x_i - b \rangle + \frac{\gamma}{2} \|\sum_i A_i x_i - b\|^2$$

• Multi-block ADMM (primal CD, dual ascent)

$$\begin{cases} x_{1} \longleftarrow \arg\min_{x_{1} \in \mathcal{X}_{1}} L_{\gamma}(x_{1}, \dots, x_{n}; \lambda), \\ \vdots \\ x_{n} \longleftarrow \arg\min_{x_{n} \in \mathcal{X}_{n}} L_{\gamma}(x_{1}, \dots, x_{n}; \lambda), \\ \lambda \longleftarrow \lambda - \gamma(Ax - b), \end{cases}$$
(2)

- 2-block ADMM converges [Glowinski-Marroco-1975], [Gabay-Mercier-1976].
- 3-block ADMM may diverge [Chen-He-Ye-Yuan-13].
- Example: solve 3 × 3 linear system

$$\begin{array}{l} \min_{x_1, x_2, x_3} & 0, \\ \text{s.t.} & \begin{bmatrix} 1 & 1 & 1 \\ 1 & 1 & 2 \\ 1 & 2 & 2 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = 0, \quad (3)$$

Random Permutation Helps

 RP-ADMM: Randomly permute update order (312), (123), (213), ...



• New outlet?



Motivation

Background

Convergence Analysis of RP-ADMM

Main Results

Proof Sketch

Variants of ADMM

Convergence Rate: Related Result and Discussion

Background

Two-block ADMM

 ADMM usually refers to 2-block ADMM [Glowinski-Marroco-75], [Gabay-Mercier-76], [Boyd-Parikh-Chu-Peleato-Eckstein-11] (5800 citations)

$$\min_{x,y} \quad f(x) + g(y)$$
s.t. $Ax + By = c$. (4)

Two-block ADMM

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$$\min_{\substack{x,y \\ \text{s.t.}}} f(x) + g(y)$$
(4)
s.t. $Ax + By = c$.

• Augmented Lagrangian function:

$$L(x, y; \lambda) = f(x) + g(y) - \langle \lambda, Ax + By - c \rangle + \frac{\gamma}{2} \|Ax + By - c\|^2.$$

Two-block ADMM:

$$\begin{cases} x \longleftarrow \arg\min_{x} L(x, y; \lambda), \\ y \longleftarrow \arg\min_{y} L(x, y; \lambda), \\ \lambda \longleftarrow \lambda - \gamma (Ax + By - c). \end{cases}$$
(5)

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Variants of multi-block ADMM

- Multi-block cyclic ADMM may diverge
- Question: How to make multi-block ADMM converge?
- Approach 1: Change algorithm.
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- Approach 2: Change algorithm + problem.
 - Strong convexity + small stepsize $\gamma = O(\sigma/N)$ [Han-Yuan-12].
- And many other related works [Deng-Lai-Peng-Yin-13], [Lin-Ma-Zhang-14],

[Lin-Ma-Zhang-15], [Sun-Toh-Yang-14], [Li-Sun-Toh-15] ,etc.

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- And many other related works [Deng-Lai-Peng-Yin-13], [Lin-Ma-Zhang-14], [Lin-Ma-Zhang-15], [Sun-Toh-Yang-14], [Li-Sun-Toh-15], etc.
- What is a minimal modification + stepsize 1?

Apply Randomization Trick to ADMM

- We know:
 - 1) ADMM may diverge;
 - 2) Randomization helps CD/SGD [Strohmer-Vershynin-08],

[Leventhal-Lewis-10], [Nesterov-11], [Roux et al-12], [Blatt et al-07]

- First idea: (independently) randomized ADMM (x₃x₁x₁λ), (x₁x₃x₂λ),...
- Bad news: can diverge!
 - Diverge for Gaussian data
 - Converge for the counter-example in [Chen-He-Ye-Yuan-13]
- Second idea: random permutation

 $(x_3x_1x_2\lambda), (x_2x_1x_3\lambda), \ldots$

It always converges in the simulation.

Summarize ADMM Variants

- **Cyclic**: $(x_1x_2x_3\lambda), (x_1x_2x_3\lambda), \dots$
- Random permutation (RP): $(x_3x_1x_2\lambda), (x_2x_1x_3\lambda), \ldots$
- Independently random (IR): $(x_3x_1x_1\lambda), (x_2x_1x_2\lambda), \ldots$

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- **Cyclic**: $(x_1x_2x_3\lambda), (x_1x_2x_3\lambda), \dots$
- Random permutation (RP): $(x_3x_1x_2\lambda), (x_2x_1x_3\lambda), \ldots$
- Independently random (IR): $(x_3x_1x_1\lambda), (x_2x_1x_2\lambda), \ldots$
- Simulation: RP always converges, other two can diverge. $\label{eq:RP} \mathsf{RP} > \mathsf{IR}, \mathsf{Cyclic}.$
- Wait...practitioners may not care? (divergence of cyclic ADMM is just worst-case?)

Numerical Experiments: Cyc-ADMM Often Diverges

Table 1: Solve Linear Systems by Cyc-ADMM, RP-ADMM and GD

N	Diverg. Ratio for Cyc-ADMM	Iterations for $\epsilon = 0.001$			
		CycADMM ⁱ	RPADMM	GD	
Gaussian N(0, 1)					
3	0.7%	3.2e01	8.8e01	1.4e02	
100	3%	1.0e03	7.4e03	6.5e03	
Uniform [0, 1]					

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Uniform [0, 1]						
3	3.2%	7.0e01	2.6e02	6.0e02		
100	100%	N/A	1.4e04	9.7e04		

- Cyc-ADMM can diverge often; sometimes diverges w.p. 100%.
 - In fact, easy to diverge if off-diagonal entries are large. Cyc-ADMM is somewhat similar to Cyc-BCD.
- RP-ADMM converges faster than GD.

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 - Not worst-case example; happen often.
 - Stepsize does not help (at least constant).
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- Order (123) fails; maybe (231) works?
- Fact: Any fixed order diverges.

Theoretical Curiosity + Practical Need.

- **First**, decomposition idea can be useful for solving constrained problems
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Theoretical Curiosity + Practical Need.

- **First**, decomposition idea can be useful for solving constrained problems
 - cyclic ADMM may not converge.
 - RP-ADMM: a simple solution
- Second, help understand RP-rule, e.g. RP-CD, RP-SGD.
 - Many people write IR papers.
 - Many people run RP experiments (default choice in deep learning package e.g. Torch)

Convergence Analysis of RP-ADMM

• Solve a square linear system of equations ($f_i = 0, \forall i$).

$$\min_{x \in \mathbb{R}^N} \quad 0,
s.t. \quad A_1 x_1 + \dots + A_n x_n = b,$$
(6)

where $A = [A_1, \dots, A_n] \in \mathbb{R}^{N \times N}$ is full-rank, $x_i \in \mathbb{R}^{d_i}$ and $\sum_i d_i = N$.

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- Why linear system?
 - Basic constrained problem
 - Already difficult to analyze.

Theorem 1ⁱⁱ

The expected output of RP-ADMM converges to the solution of (6), i.e.

$$\{E_{\xi_k}(y^k)\}_{k\to\infty} \longrightarrow y^*.$$
(7)

Remark: Expected convergence \neq convergence, but is a strong evidence for convergence.

Denote *M* as the expected iteration matrix of RP-ADMM.

Theorem 2

 $\rho(M) < 1$, i.e. spectral radius of *M* is less than 1.

[&]quot;S, Luo, Yinyu Ye, "On the Expected Convergence of Randomly Permuted ADMM",

Meta-proof-frameworks in optimization don't work (or I don't know how).

Potential function.

- E.g. GD, C-CD or R-CD for *min_xx^TAx*, the potential function is the (expected) objective.
- Our system: E(y^{k+1}) = ME(y^k), but ||M|| > 2.3 for the counterexample. y^TMy is not a potential function.
- There exists *P* such that $P M^T PM$ is PSD, and $y^T Py$ is a potential function. Hard to compute *P*.

Contraction: can prove convergence of 2-block ADMM.

- Again, how to distinguish between cyclic ADMM and PR-ADMM?
- Not a big surprise. 2-block is very special.

RP-ADMM can be viewed as switched linear systems:

 $y_{k+1} = M_k y_k,$

where $M_k \in \{B_1, \ldots, B_m\}$. For RP-ADMM, m = n!.

Our problem: each single B_i is not stable (corresponding to a single order), but randomly picking from $\{B_1, \ldots, B_m\}$ makes the system stable.

Related to product of random matrices [Furstenberg-Kesten-60]; but hard to apply to our case.

A useful first step is to find a convex combination of B_i 's that is stable [Wicks et al.-94]

Theorem 2: a Pure Linear Algebra Problem

• Define matrix L_{σ} by deleting half off-diagonal entries of $A^{T}A$

$$L_{\sigma}[\sigma(i), \sigma(j)] \triangleq \begin{cases} A_{\sigma(i)}^{T} A_{\sigma(j)} & j \le i, \\ 0 & j > i, \end{cases}$$
(8)

• Example:

$$L_{(231)} = \begin{bmatrix} 1 & A_1^T A_2 & A_1^T A_3 \\ 0 & 1 & 0 \\ 0 & A_3^T A_2 & 1 \end{bmatrix}.$$

Theorem 2: a Pure Linear Algebra Problem

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• Define $Q = E(L_{\sigma}^{-1})$. Compare: $E(L_{\sigma}) = \frac{1}{2}(I + A^{T}A)$.

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- Define $Q = E(L_{\sigma}^{-1})$. Compare: $E(L_{\sigma}) = \frac{1}{2}(I + A^{T}A)$.
- Theorem 2 claims $\rho(M) < 1$, with *M* being a function of *A*:

$$M = \begin{bmatrix} I - QA^{T}A & QA^{T} \\ -A + AQA^{T}A & I - AQA^{T} \end{bmatrix}.$$
 (9)

• Step 1: Relate M to a symmetric matrix AQA^{T} .

Lemma 1

$$\lambda \in \operatorname{eig}(M) \iff \frac{(1-\lambda)^2}{1-2\lambda} \in \operatorname{eig}(AQA^T).$$
 (10)

When *Q* is symmetric, we have $\rho(M) < 1 \iff \text{eig}(AQA^T) \subseteq (0, \frac{4}{3}).$ (11)

• This lemma treats *Q* as a black box.

• Step 2: Bound eigenvalues of AQA^T.

Lemma 2

For any non-singular *A*, let $Q = E(L_{\sigma}^{-1})$ where L_{σ} is given by (8), then

$$\operatorname{eig}(AQA^{T}) \subseteq (0, \frac{4}{3}). \tag{12}$$

• **Remark**: 4/3 should be tight: we find examples > 1.33.

- AQA^T relates to **RP-CD** (quadratic): $x \leftarrow (I QA^T A)x$.
 - **RP-CD** converges \iff eig(AQA^T) \in (0, 2).

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 - RP-ADMM converges $\iff eig(AQA^T) \in (0, 4/3).$
- Cyc-CD (quadratic): $x \leftarrow (I L_{12...n}^{-1} A^T A) x$
 - Cyc-CD converges $\iff eig(AL_{12...n}^{-1}A^{T}) \in (0, 2).$

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- Cyc-CD (quadratic): $x \leftarrow (I L_{12...n}^{-1} A^T A) x$
 - Cyc-CD converges $\iff eig(AL_{12...n}^{-1}A^T) \in (0, 2).$
- Remark: spectrum of RP-CD is "nicer" than Cyc-CD.
 - "Pre-assigned" space for RP-ADMM.

- Step 2.1: Symmetrization \implies induction formula of $Q = E(L_{\sigma}^{-1}).$
- Step 2.2: Induction inequality of $\rho = \rho(QA^TA)$:

$$\rho \leq \boldsymbol{P}(\hat{\rho}, \rho) \triangleq \max_{\theta \geq 0} \hat{\rho} + \theta \left(\frac{\rho}{4\rho - 4 + \theta} - 1 \right), \quad (13)$$

where $\hat{\rho}$ is the (n-1)-block analog of $\rho(QA^TA)$.

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- **Remark**: $\rho = 4/3$ is the fixed point of $\rho = P(\rho, \rho)$.
 - $P(\frac{4}{3},\frac{4}{3}) = \frac{4}{3} + \max_{\theta \ge 0} \theta(\frac{\rho}{\rho+\theta} 1) = \frac{4}{3} \max_{\theta \ge 0} \frac{\theta^2}{\rho+\theta} = \frac{4}{3}.$

Variants of ADMM

Interesting Byproduct: New Randomization Rule

- Finding: 2-level symmetrization is enough.
- New algorithm: Bernolli randomization (BR).
 - Phase 1: sweep 1,..., n; for each block, update w.p. 1/2;
 - Phase 2: sweep *n*,..., 1; if previously not updated, now update.
- Examples of valid order: (2, 3; 4, 1), (1, 2, 4; 3).
 Non-examples: (3, 4, 1, 2)

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- Examples of valid order: (2, 3; 4, 1), (1, 2, 4; 3).
 Non-examples: (3, 4, 1, 2)
- Proposition: BR-ADMM converges in expectation.

The problem is still $\min_x f(x)$, s.t. Ax = b.

Original ADMM: each cycle is (x_1, x_2, x_3, λ) .

Primal-dual ADMM: each cycle is $(x_1, \lambda, x_2, \lambda, x_3, \lambda)$.

- Cyclic version still can diverge for the counter-example.
- Randomized version was proven to converge with high probability (e.g. [Xu-2017])

However, in simulation, randomized PD-ADMM is much slower than other versions (next page).

Comparison of Algorithms

Uniform [0,1] data:

- cyclic ADMM and primal-dual version of Bernolli randomization fail to converge.
- PD-rand-ADMM is much slower than others.



Comparison of Algorithms (cont'd)

Standard Gaussian data:

- PD-rand-ADMM is significantly slower than all other methods.
- Recall: randomized ADMM is the only method that diverges!

Strange issue: (independent) random rule is bad for Gaussian data.



Simple summary of different methods (no stepsize tuning):

Update Order	Original Version	Primal-Dual Version	
cyclic	Diverge	Diverge	
indep. random	Diverge	Converge but very slow	
Bernolli random	Converge	Diverge	
random permutation	Converge ⁱⁱⁱ	Converge? iv	

Observation: random permutation is a universal "stabilizer".

Open question: Any convergence analysis of P-D version of RP-ADMM?

ⁱⁱⁱOnly expected convergence for simple problems are proved. ^{iv}Based on extensive simulation

Convergence Rate: Related Result and Discussion

Convergence Rate of Cyclic CD

- Status: many results on (independently) random rule; little understanding of RP/cyclic/whatever rule
 - A few works [Recht-Re-12], [Gurbuzbalaban-Ozdaglar-Parrilo-15], [Wright-Lee-17] studied random permutation, but why RP is better than IR in general is still unknown
 - Mark Schmidt talked about Gauss-Southwell rule this morning.
- Classical literature says: they are "essentially cyclic" rule, all converge for CD
- However, their convergence speed can be quite different

Convergence Rate of Cyclic CD

- **Question**: "true" convergence rate of cyclic CD or Gauss-Seidal method (Gauss 1823, Seidel 1874)?
- Why care cyclic order?
 - Understanding "non-independently-randomized" rule
 - Almost all convergence rate results on cyclic rule immediately apply to RP-rule
 - Randomization not available sometimes
- Puzzle: known rates can be sometimes n² times worse than R-CD for quadratic case [Beck-Tetruashvili-13], [Sun-Hong-15]
- Some claim cyclic order must be bad; an example given by Strohmer and Richtarik (independently) showed this.
 - Only *O*(*n*) gap between C-CD and R-CD;
 - Only fails for some particular orders. Randomly pick order and fix, then becomes fast.

Rate of Cyclic CD

• **Answer** ^v: up to *n*² times worse than R-CD, for equal-diagonal quadratic case.

Table 2: Complexity for equal-diagonal case (divided by $n^2 \kappa \log \frac{1}{\epsilon}$ and ignoring constants. $\tau = \lambda_{\max} / \lambda_{avg} \in [1, n]$)

	C-CD	GD	R-CD	SVRG
Lower bound	au	1	$\frac{1}{\tau}$	$\frac{1}{\tau}$
Upper bound	$\min\{\tau \log^2 n, n\}$	1	$\frac{1}{\tau}$	$\frac{1}{\tau}$

- Lower bound is based on analyzing one example. Steven Wright mentioned the example in two talks starting from 2015 summer. We independently discover the example.
- Analysis: tricky issue on non-symmetric matrix update. (even more tricky than ADMM case)

^vSun, Ye, "Worst-case Convergence Rate of Cyclic Coordinate Descent Method: $O(n^2)$ Gap with Randomized Versions", 2016.

Relation to Other Methods

- Same gap exists for Kaczmarz method and POCS (Projection onto Convex Sets).
- POCS, dating back to Von Neumann in 1930's, has been studied extensively. See a survey [Bauschke-Borwein-Lewis-1997]
- Convergence rate given by Smith, Solmon and Wagner in 1977. Still in textbook.
- Translate to CD: a rate dependent on all eigenvalues.
 - Turn out to be ∞ -times worse than our bound for the example.
 - Always worse than our bound (up to $O(\log^2 n)$ factor)

Convergence Speed of RP-CD and AM-GM inequality

Random permutation was studied in [Recht-Re'2012], mainly for RP-SGD.

Conjecture: Matrix AM-GM inequality ([Recht-Re'2012]) Suppose $A_1, \ldots, A_n \succeq 0$, then

$$\|\frac{1}{n!}\sum_{\sigma \text{ is a permutation}} A_{\sigma_1} \dots A_{\sigma_n}\| \leq \|\frac{1}{n}(A_1 + \dots + A_n)\|^n.$$

If this inequality holds, then the convergence rate of RP-CD for quadratic problems is faster than R-CD.

Zhang gave a proof for n = 3; Duchi gave a proof for a variant, again for n = 3.

Another variant of matrix AM-GM inequality

Conjeture (variant of matrix AM-GM inequality): If P_i is a projection matrix, i = 1, ..., n, then

$$\frac{1}{n!} \sum_{\sigma \text{ is a permutation}} P_{\sigma_1} \dots P_{\sigma_n} \leq \frac{1}{n} (P_1 + \dots + P_n).$$
(14)

Claim: If matrix AM-GM inequality (14) holds, then combining with our result $eig(QA^TA) \in (0, 4/3)$, RP-CD has better convergence rate than that of R-CD for convex quadratic problems.

We know $\operatorname{eig}(I - QA^T A) = \operatorname{eig}(M_{RP-CD}) \in (-1, 1).$

- Our result is about the left end by improving -1 to -1/3.
- Matrix AM-GM inequality (14) is about the right end near 1

We have some results on the expected convergence rate of RP-CD and RP-ADMM. Skip here.

Summary

- Main result: convergence analysis of RP-ADMM.
 - Implication 1 (problem): solver for constrained problems.
 - Implication 2 (algorithm): RP better. Even much better than independently randomized rule.
- Implication for RP-CD: "truncate" one side spectrum.
- Tight analysis of "non-independent-randomization": worst-case understanding of cyclic order, but more works are needed.
- Lots of open questions:
 - convergence of PD version of RP-ADMM
 - AM-GM inequality
 - Jacobi preconditioning

Thank You!