On the  $O(1/k)$  Convergence of Asynchronous Distributed Alternating Direction Method of Multipliers (ADMM)

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#### **Motivation**

- Many networks are large-scale and comprise of agents with local information and heterogeneous preferences.
- This motivated much interest in developing distributed schemes for control and optimization of multi-agent networked systems.



Routing and congestion control in wireline and wireless networks



Parameter estimation in sensor networks

<span id="page-1-0"></span>Multi-agent cooperative control and coordination



Smart grid systems

#### Distributed Multi-agent Optimization Distributed Optimization for General Objective Functions

- Many of these problems can be represented within the general formulation:
- A set of agents (nodes)  $\{1,\ldots,N\}$  connected through a network. A set of agents (nodes)  $\{1,\ldots,N\}$
- The goal is to cooperatively solve

 $\min_{x}$   $\sum_{i=1}^{N}$  $i=1$  $f_i(x)$ s.t.  $x \in \mathbb{R}^n$ ,  $\sum f_i(x)$ s.t.  $x \in \mathbb{R}^n$ ,

<span id="page-2-0"></span>

 $f_i(x) : \mathbb{R}^n \to \mathbb{R}$  is a convex *•*<sub>1</sub>(*×*) • <del><sub>■</sub>∗ • / ■∗ is a convex</del><br>(possibly nonsmooth) function, known only to agent *i*. optimal solution set by *X*<sup>∗</sup> (assumed

Since such systems often lack a centralized processing unit, algorithms for this problem should involve each agent performing computations locally and communicating this information according to the underlying network.

#### Machine Learning Example

- A network of 3 sensors, supervised passive learning.
- <span id="page-3-0"></span> $\bullet$  Data is collected at different sensors: temperature t, electricity demand d.



#### Machine Learning General Set-up

- A network of agents  $i = 1, \ldots, N$ .
- Each agent  $i$  has access to local feature vectors  $A_i$  and output  $b_i$ .
- $\bullet$  System objective: train weight vector x to

$$
\min_x \sum_{i=1}^{N-1} L(A'_i x - b_i) + p(x),
$$

for some loss function L (on the prediction error) and penalty function  $p$  (on the complexity of the model).

Example: Least-Absolute Shrinkage and Selection Operator (LASSO):

<span id="page-4-0"></span>
$$
\min_{x} \quad \sum_{i=1}^{N-1} ||A'_i x - b_i||_2^2 + \lambda ||x||_1.
$$

Other examples from ML estimation, low rank matrix completion, image recovery [Schizas, Ribeiro, Giannakis 08], [Recht, Fazel, Parrilo 10], [Steidl, Teuber, 10]

#### Existing Distributed Algorithms

• Given an undirected connected graph  $G = \{V, E\}$  with M nodes, we reformulate the problem as



- Distributed gradient/subgradient methods for solving these problems:
	- Each agent maintains an local estimate, updates it by taking a (sub)gradient step and averaging with neighbors' estimates.
	- Best known convergence rate:  $O(1/\sqrt{k})$ .[Nedic, Ozdaglar 08], [Lobel, Ozdaglar 09], [Duchi, Agarwal, Wainwright 12].

<span id="page-5-0"></span> $x_3 = x_4$ 

#### Faster ADMM-based Distributed Algorithms

- **Classical Augmented Lagrangian/Method of Multipliers and Alternating** Direction Method of Multipliers (ADMM) methods: fast and parallel [Glowinski, Marrocco 75], [Eckstein, Bertsekas 92], [Boyd et al. 10]:
- Known convergence rates for synchronous ADMM type algorithm:
	- [He, Yuan 11] General convex  $O(1/k)$ .
	- [Goldfarb et al. 10] Lipschitz gradient  $O(1/k^2)$ .
	- [Deng, Yin 12] Lipschitz gradient, strong convexity linear rate.
	- [Hong, Luo 12] Strong convexity linear rate.
- Highly decentralized nature of the problem calls for an asynchronous algorithm. Almost all known distributed algorithms are synchronous.<sup>1</sup>
- <span id="page-6-0"></span>• In this talk, we present asynchronous ADMM-type algorithms for general convex problems and show that it converges at the best known rate of  $O(1/k)$  [Wei, Ozdaglar 13].

 $^1$ Exceptions: [Ram, Nedic, Veeravalli 09], [lutzeler, Bianchi, Ciblat, and Hachem 13] without any rate results.

#### Standard ADMM

• Standard ADMM solves a separable problem, where decision variable decomposes into two (linearly coupled) variables:

> <span id="page-7-1"></span><span id="page-7-0"></span> $\min_{x,y} f(x) + g(y)$  (1) s.t.  $Ax + By = c$ .

**• Consider an Augmented Lagrangian function:** 

$$
L_{\beta}(x, y, p) = f(x) + g(y) - p'(Ax + By - c) + \frac{\beta}{2} ||Ax + By - c||_2^2.
$$

ADMM: approximate version of classical Augmented Lagrangian method.

- Primal variables: approximately minimize augmented Lagrangian through a single-pass coordinate descent (in a Gauss-Seidel manner).
- Dual variable: updated through gradient ascent.

### Standard ADMM

More specifically, updates are as follows:

<span id="page-8-0"></span>
$$
x^{k+1} = \operatorname{argmin}_{x} L_{\beta}(x, y^{k}, p^{k}),
$$
  
\n
$$
y^{k+1} = \operatorname{argmin}_{y} L_{\beta}(x^{k+1}, y, p^{k}),
$$
  
\n
$$
p^{k+1} = p^{k} - \beta(Ax^{k+1} - By^{k+1} - c).
$$

- **Each minimization involves (quadratic perturbations of) functions f and g** separately.
	- In many applications, these minimizations are easy (quadratic minimization,  $l_1$  minimization, which arises in Huber fitting, basis pursuit, LASSO, total variation denoising). [Boyd et al. 10]

#### ADMM for Multi-agent Optimization Problem

- Multi-agent optimization problem can be reformulated in the ADMM framework:
- Consider a set of agents  $V = \{1, \ldots, N\}$  connected through an undirected connected graph  $G = \{V, E\}.$
- We introduce a local copy  $x_i$  for each of the agents and impose  $x_i = x_i$  for all  $(i, j) \in E$ .

<span id="page-9-0"></span>min x X N i=1 fi(xi) s.t. x<sup>i</sup> = x<sup>j</sup> , for (i, j) ∈ E, ! "# \$# %# &# f2(x2) f1(x1) f3(x3) f4(x4) f5(x5) x<sup>1</sup> = x<sup>2</sup> x<sup>3</sup> = x<sup>4</sup> x<sup>1</sup> = x<sup>4</sup> x<sup>2</sup> = x<sup>3</sup>

Multi-agent optimization problem with two agents: special case of problem [\(1\)](#page-7-1):

> <span id="page-10-0"></span> $\min_{x_1, x_2}$   $f_1(x_1) + f_2(x_2)$ s.t.  $x_1 = x_2$ .

• ADMM applied to this problem yields:



 $x_1^{k+1} = \text{argmin}_{x_1} \quad f_1(x_1) + f_2(x_2^k) - (p_{12}^k)'(x_1 - x_2^k) + \frac{\beta}{2}||x_1 - x_2^k||$ 2 2

Multi-agent optimization problem with two agents: special case of problem [\(1\)](#page-7-1):

> <span id="page-11-0"></span> $\min_{x_1, x_2}$   $f_1(x_1) + f_2(x_2)$ s.t.  $x_1 = x_2$ .

• ADMM applied to this problem yields:



• 
$$
x_1^{k+1}
$$
 = argmin<sub>x\_1</sub>  $f_1(x_1) - (p_{12}^k)'x_1 + \frac{\beta}{2}||x_1 - x_2^k||_2^2$ 

Multi-agent optimization problem with two agents: special case of problem [\(1\)](#page-7-1):

<span id="page-12-0"></span>
$$
\min_{x_1, x_2} f_1(x_1) + f_2(x_2) \ns.t. x_1 = x_2.
$$

• ADMM applied to this problem yields:



 $x_2^{k+1} =$  $\operatorname{argmin}_{x_2} \quad f_1(x_1^{k+1}) + f_2(x_2) - (p_{12}^k)'(x_1^{k+1} - x_2) + \frac{\beta}{2}$   $x_1^{k+1} - x_2$  $\begin{array}{c} \begin{array}{c} \begin{array}{c} \end{array} \\ \begin{array}{c} \end{array} \end{array} \end{array}$ 2 2

Multi-agent optimization problem with two agents: special case of problem [\(1\)](#page-7-1):

<span id="page-13-0"></span>
$$
\min_{x_1, x_2} f_1(x_1) + f_2(x_2) \ns.t. x_1 = x_2.
$$

• ADMM applied to this problem yields:



• 
$$
x_2^{k+1}
$$
 = argmin<sub>x2</sub>  $f_2(x_2) + (p_{12}^k)'x_2 + \frac{\beta}{2}||x_1^{k+1} - x_2||_2^2$ 

Multi-agent optimization problem with two agents: special case of problem [\(1\)](#page-7-1):

<span id="page-14-0"></span>
$$
\min_{x_1, x_2} f_1(x_1) + f_2(x_2)
$$
  
s.t.  $x_1 = x_2$ .

• ADMM applied to this problem yields:



 $p^{k+1} = p^k - \beta(x_1^{k+1} - x_2^{k+1}).$ 

#### Multi-agent Asynchronous ADMM - Problem Formulation

<span id="page-15-0"></span>
$$
\min_{x} \sum_{i=1}^{N} f_i(x_i)
$$
  
s.t.  $x_i = x_j$ , for  $(i,j) \in E$ .

• Reformulate to decouple  $x_i$  and  $x_j$  by introducing the auxiliary z variable [Bertsekas, Tsitsiklis 89], which allows us to simultaneously update  $x_i$  and potentially improves performance.



#### Multi-agent Asynchronous ADMM - Algorithm



• Set 
$$
Z = \{ z \mid z_{ei} + z_{ej} = 0 \text{ for all } e = (i, j) \}.
$$

• Write constraint as  $Dx = z$ , set  $E(i)$ : the set of edges incident to node i.

- We associate an independent Poisson local clock with each edge.
- <span id="page-16-0"></span> $\bullet$  At iteration k, if the clock corresponding to edge  $(i, j)$  ticks:
	- The constraint  $x_i = z_{ei}$ ,  $-x_i = z_{ei}$  (subject to  $z_{ei} + z_{ei} = 0$ ) is active.
	- $\bullet$  The agents *i* and *j* are active.
	- The dual variables  $p_{ei}$  and  $p_{ei}$  associated with edge  $(i, j)$  are active.

- ${\sf A}$  Initialization: choose some arbitrary  $x^0$  in  $X$ ,  $z^0$  in  $Z$  and  $p^0=0.$
- B At time step k, an edge  $e = (i, j)$  and its end points become active.
	- a The active primal variables  $x_q$  for  $q = i, j$  are updated as

$$
x_q^{k+1} \in \mathop{\rm argmin}_{x_q \in X_q} f_q(x_q) - \sum_{e \in E(q)} (p_{eq}^k)' D_{eq} x_q + \frac{\beta}{2} \sum_{e \in E(q)} ||D_{eq} x_q - z_{eq}^k||^2.
$$

with  $x_w^{k+1} = x_w^k$  for w not active. **b** The active primal variables  $z_{ei}$  and  $z_{ei}$  are updated as

$$
z_{ei}^{k+1}, z_{ej}^{k+1} \in \mathop{\rm argmin}_{z_{ei}+z_{ej}=0} \sum_{q=i,j} (p_{eq}^k)' z_{eq} + \frac{\beta}{2} ||D_{eq} x_q^{k+1} - z_{eq}||^2.
$$

with  $z_j^{k+1} = z_j^k$  for *l* not active.

c The active dual variables  $p_{eq}$  for  $q = i, j$  are updated as

<span id="page-17-0"></span>
$$
p_{eq}^{k+1} = p_{eq}^{k} - \beta \left[ D_q x_q^{k+1} - z_{eq}^{k+1} \right].
$$

• Update in z is a quadratic programming with linear constraint: has closed form solution and can be easily computed in a distributed way.

- ${\sf A}$  Initialization: choose some arbitrary  $x^0$  in  $X$ ,  $z^0$  in  $Z$  and  $p^0=0.$
- B At time step k, an edge  $e = (i, j)$  and its end points become active.
	- a The active primal variables  $x_q$  for  $q = i, j$  are updated as

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x_q^{k+1} \in \mathop{\rm argmin}_{x_q \in X_q} f_q(x_q) - \sum_{e \in E(q)} (p_{eq}^k)' D_{eq} x_q + \frac{\beta}{2} \sum_{e \in E(q)} ||D_{eq} x_q - z_{eq}^k||^2.
$$

with  $x_w^{k+1} = x_w^k$  for w not active.

**b** The active primal variables  $z_{ei}$  and  $z_{ei}$  are updated as

$$
z_{ei}^{k+1}, z_{ej}^{k+1} \in \mathop{\rm argmin}_{z_{ei}+z_{ej}=0} \sum_{q=i,j} (p_{eq}^k)' z_{eq} + \frac{\beta}{2} ||D_{eq} x_q^{k+1} - z_{eq}||^2.
$$

with  $z_j^{k+1} = z_j^k$  for *l* not active.

c The active dual variables  $p_{eq}$  for  $q = i, j$  are updated as

<span id="page-18-0"></span>
$$
p_{eq}^{k+1} = p_{eq}^{k} - \beta \left[ D_q x_q^{k+1} - z_{eq}^{k+1} \right].
$$

Update in z is a quadratic programming with linear constraint: has closed form solution and can be easily computed in a distributed way.

- $\mathsf A$  Initialization: choose some arbitrary  $x^0$  in  $X$ ,  $z^0$  in  $Z$  and  $p^0=0.$
- B At time step k, an edge  $e = (i, j)$  and its end points become active.

a For  $q = i, j$ , the active primal variable  $x_q$  is updated as

$$
x_q^{k+1} \in \mathop{\rm argmin}_{x_q \in X_q} f_q(x_q) - \sum_{e \in E(q)} (p_{eq}^k)' D_{eq} x_q + \frac{\beta}{2} \sum_{e \in E(q)} \left| \left| D_{eq} x_q - z_{eq}^k \right| \right|^2.
$$

with  $x_w^{k+1} = x_w^k$  for w not active. b To compute z update,

$$
v^{k+1} = \frac{1}{2}(-p_{ei}^k - p_{ej}^k) + \frac{\beta}{2}(D_{ei}x_i^{k+1} + D_{ej}x_j^{k+1}),
$$
  

$$
z_{eq}^{k+1} = \frac{1}{\beta}(-p_{eq}^k - v^{k+1}) + D_{eq}x_q^{k+1}.
$$

c The active dual variables  $p_{eq}$  for  $q = i, j$  are updated as

<span id="page-19-0"></span>
$$
p_{eq}^{k+1}=-\nu^{k+1}.
$$

- $\mathsf A$  Initialization: choose some arbitrary  $x^0$  in  $X$ ,  $z^0$  in  $Z$  and  $p^0=0.$
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x_q^{k+1} \in \mathop{\rm argmin}_{x_q \in X_q} f_q(x_q) - \sum_{e \in E(q)} (p_{eq}^k)' D_{eq} x_q + \frac{\beta}{2} \sum_{e \in E(q)} \left| \left| D_{eq} x_q - z_{eq}^k \right| \right|^2.
$$

with  $x_w^{k+1} = x_w^k$  for w not active.

 $b$  To compute z update,

$$
v^{k+1} = \frac{1}{2}(-p_{ei}^k - p_{ej}^k) + \frac{\beta}{2}(D_{ei}x_i^{k+1} + D_{ej}x_j^{k+1})
$$
  

$$
z_{eq}^{k+1} = \frac{1}{\beta}(-p_{eq}^k - v^{k+1}) + D_{eq}x_q^{k+1}.
$$

c The active dual variables  $p_{eq}$  for  $q = i, j$ are updated as

$$
p_{eq}^{k+1}=-\nu^{k+1}.
$$

<span id="page-20-0"></span>

- $\mathsf A$  Initialization: choose some arbitrary  $x^0$  in  $X$ ,  $z^0$  in  $Z$  and  $p^0=0.$
- B At time step k, an edge  $e = (i, j)$  and its end points become active.
	- a For  $q = i$ , *i*, the active primal variable  $x_q$  is updated as

$$
x_q^{k+1} \in \mathop{\rm argmin}_{x_q \in X_q} f_q(x_q) - \sum_{e \in E(q)} (p_{eq}^k)' D_{eq} x_q + \frac{\beta}{2} \sum_{e \in E(q)} \left| \left| D_{eq} x_q - z_{eq}^k \right| \right|^2.
$$

with  $x_w^{k+1} = x_w^k$  for w not active.

b To compute z update,

$$
v^{k+1} = \frac{1}{2}(-p_{ei}^k - p_{ej}^k) + \frac{\beta}{2}(D_{ei}x_i^{k+1} + D_{ej}x_j^{k+1}
$$

$$
z_{eq}^{k+1} = \frac{1}{\beta}(-p_{eq}^k - v^{k+1}) + D_{eq}x_q^{k+1}.
$$

c The active dual variables  $p_{eq}$  for  $q = i, j$ are updated as

$$
p_{eq}^{k+1}=-\nu^{k+1}.
$$

<span id="page-21-0"></span>

- $\mathsf A$  Initialization: choose some arbitrary  $x^0$  in  $X$ ,  $z^0$  in  $Z$  and  $p^0=0.$
- B At time step k, an edge  $e = (i, j)$  and its end points become active.
	- a For  $q = i, j$ , the active primal variable  $x_q$  is updated as

$$
x_q^{k+1} \in \mathop{\rm argmin}_{x_q \in X_q} f_q(x_q) - \sum_{e \in E(q)} (p_{eq}^k)' D_{eq} x_q + \frac{\beta}{2} \sum_{e \in E(q)} \left| \left| D_{eq} x_q - z_{eq}^k \right| \right|^2.
$$

with  $x_w^{k+1} = x_w^k$  for w not active.

- b To compute z update,  $v^{k+1} = \frac{1}{2}$  $\frac{1}{2}(-\rho_{\text{\rm ei}}^k - \rho_{\text{\rm ej}}^k) + \frac{\beta}{2}(D_{\text{\rm ei}}x_i^{k+1} + D_{\text{\rm ej}}x_j^{k+1})$  $=-p_{ei}^k+\frac{\beta}{2}$  $\frac{\beta}{2}(D_{ei}x_i^{k+1}+D_{ej}x_j^{k+1}),$  $z_{eq}^{k+1}=\frac{1}{\varphi}$  $\frac{1}{\beta}(-p_{eq}^k-v^{k+1})+D_{eq}x_q^{k+1}.$
- c The active dual variables  $p_{eq}$  for  $q = i, j$ are updated as

$$
p_{eq}^{k+1}=-\nu^{k+1}.
$$

<span id="page-22-0"></span>

- $\mathsf A$  Initialization: choose some arbitrary  $x^0$  in  $X$ ,  $z^0$  in  $Z$  and  $p^0=0.$
- B At time step k, an edge  $e = (i, j)$  and its end points become active.
	- a For  $q = i$ , *i*, the active primal variable  $x_q$  is updated as

$$
x_q^{k+1} \in \mathop{\rm argmin}_{x_q \in X_q} f_q(x_q) - \sum_{e \in E(q)} (p_{eq}^k)' D_{eq} x_q + \frac{\beta}{2} \sum_{e \in E(q)} \left| \left| D_{eq} x_q - z_{eq}^k \right| \right|^2.
$$

with  $x_w^{k+1} = x_w^k$  for w not active.

 $b$  To compute z update,

$$
v^{k+1} = -p_{ei}^k + \frac{\beta}{2}(D_{ei}x_i^{k+1} + D_{ei}x_j^{k+1}),
$$

$$
z_{eq}^{k+1} = \frac{1}{\beta}(-p_{eq}^{k} - v^{k+1}) + D_{eq}x_{q}^{k+1}.
$$

c The active dual variables  $p_{eq}$  for  $q = i, j$ are updated as

$$
p_{eq}^{k+1}=-\nu^{k+1}.
$$

<span id="page-23-0"></span>

- $\mathsf A$  Initialization: choose some arbitrary  $x^0$  in  $X$ ,  $z^0$  in  $Z$  and  $p^0=0.$
- B At time step k, an edge  $e = (i, j)$  and its end points become active.
	- a For  $q = i, j$ , the active primal variable  $x_q$  is updated as

$$
x_q^{k+1} \in \mathop{\rm argmin}_{x_q \in X_q} f_q(x_q) - \sum_{e \in E(q)} (p_{eq}^k)' D_{eq} x_q + \frac{\beta}{2} \sum_{e \in E(q)} \left| \left| D_{eq} x_q - z_{eq}^k \right| \right|^2.
$$

with  $x_w^{k+1} = x_w^k$  for w not active.

b To compute z update,

$$
v^{k+1} = -p_{ei}^k + \frac{\beta}{2} (D_{ei}x_i^{k+1} + D_{ej}x_j^{k+1}),
$$
  

$$
z_{eq}^{k+1} = \frac{1}{\beta} (-p_{eq}^k - v^{k+1}) + D_{eq}x_q^{k+1}.
$$

c The active dual variables  $p_{eq}$  for  $q = i, j$ are updated as

$$
p_{eq}^{k+1}=-\nu^{k+1}.
$$

Generalizes to any linear constraint  $Dx + Hz = 0$ .

<span id="page-24-0"></span>

#### **Convergence**

#### Assumption

(a) (Infinitely often updates): For all  $k$  and all  $l$  in the set of linear constraints,  $\mathbb{P}(l$  is active at time  $k$ ) > 0.

#### Theorem

Let  $\{x^k, z^k, p^k\}$  be the iterates generated by the general asynchronous ADMM algorithm. The sequence  $\{x^k, z^k, p^k\}$  converges to a saddle point  $(x^*, z^*, p^*)$  of the Lagrangian, i.e.,  $(x^k, z^k)$  converges to a primal optimal solution  $(x^*, z^*)$  almost surely.

#### Proof Sketch

Define auxiliary full information iterates  $y^k$ ,  $v^k$  and  $\mu^k$ .

<span id="page-25-0"></span>
$$
y^{k+1} \in \operatorname*{argmin}_{y \in X} \sum_{i=1}^{N} f_i(y_i) - (p^k - \beta Hz^k)'D_iy + \frac{\beta}{2} ||D_iy||^2,
$$
  

$$
v^{k+1} \in \operatorname*{argmin}_{v \in Z} \sum_{l=1}^{W} -(p^k - \beta Dy^{k+1})'H_iv + \frac{\beta}{2} ||H_iv||^2,
$$
  

$$
\mu^{k+1} = p^k - \beta(Dy^{k+1} + Hv^{k+1}).
$$

#### Convergence Analysis – Idea

- Active components of asynchronous iterates take the same value as full information iterates, inactive components remain at their previous value.
- Using the Lyapunov function  $\frac{1}{2\beta}$   $||p^{k+1} p^*||$  $2 + \frac{\beta}{2} ||H(z^{k+1} - z^*)||$ 2 , we can show full information iterates converge to an optimal solution.
- To develop a Lyapunov function for the asynchronous iterates, define probabilities

<span id="page-26-0"></span> $\lambda_l = \mathbb{P}(l$  is active at time k)

and weighted norm induced by matrix  $\bar{\mathsf{\Lambda}}$  where  $\bar{\mathsf{\Lambda}}_{ll}=1/\lambda_l.$ 

Using supermartingale arguments, we show that the probability adjusted norm,

$$
\frac{1}{2\beta} \left| \left| p^{k+1} - p^* \right| \right|_{\bar{\Lambda}}^2 + \frac{\beta}{2} \left| \left| H(z^{k+1} - z^*) \right| \right|_{\bar{\Lambda}}^2
$$

serves as a Lyapunov function for the asynchronous iterates.

#### Rate of Convergence

#### **Assumption**

(a) (Compact constraint set): Sets  $X$  and  $Z$  are compact.

• Ergodic average: 
$$
\bar{x}_i(k) = \frac{\sum_{t=1}^k x_i^t}{k}
$$
, for all *i*,  $\bar{z}_i(T) = \frac{\sum_{t=1}^k z_i^t}{k}$ , for all *l*.

#### Theorem

For  $F(x) = \sum_{i=1}^{N} f_i(x_i)$ , the iterates generated by the asynchronous ADMM algorithm satisfies

<span id="page-27-0"></span>
$$
||\mathbb{E}(F(\bar{x}(k))) - F(x^*)|| \leq \frac{\alpha}{k},
$$

where  $\alpha = ||p^*||_{\infty} \left[ \bar{Q} + \tilde{L}^0 + \frac{1}{2\beta} \left| \left|p^0 - p^* \right| \right|^2_{\bar{\Lambda}} + \frac{\beta}{2} \left| \left| H(z^0 - z^*) \right| \right|^2_{\bar{\Lambda}} \right] +$  $\Big[Q(p^*)+\tilde{L}(x^0,z^0,p^*)+\frac{1}{2\beta}\left|\left|p^0-\bar{\theta}\right|\right|_{\bar{\Lambda}}^2+\frac{\beta}{2}\left|\left|H(z^0-z^*)\right|\right|_{\bar{\Lambda}}^2\Big],$  for some scalar Q,  $\bar{Q}, \bar{\Lambda},$  $\bar{\theta}$ , related to p $^*$  and size of set  $X$  and Z.

A similar rate result holds for the constraint violation  $||E(D\bar{x}(k) + H\bar{z}(k))||$ .

#### [Simulations](#page-28-0)

#### **Simulations**

**• Sample Network:** 

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- Asynchronous ADMM algorithm is compared against a gradient based asynchronous gossip algorithm [Ram, Nedic, Veeravalli 09]
- Tested in three 5−node graphs: sample network, line graph and complete graph.

### Sample network



Figure: ADMM for the sample network.



<span id="page-29-0"></span>Figure: Asynchronous gossip for the sample network.

To reach 5% neighborhood of the optimal solution: asynchronous ADMM takes 80 iterations, asynchronous gossip takes 250 iterations.

#### Line Graph



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To reach 5% neighborhood of the optimal solution: asynchronous ADMM takes 70 iterations, asynchronous gossip takes 700 iterations.

#### Complete Graph



Figure: ADMM for the complete graph.

<span id="page-31-0"></span>Figure: Asynchronous gossip for the complete network.

To reach 5% neighborhood of the optimal solution: asynchronous ADMM takes 140 iterations, asynchronous gossip takes 380 iterations.



#### Image denoising

Given a noisy image measure  $b$ , recover the original image by solving the following problem:

$$
\min_{x} \frac{1}{2} ||x - b||_2^2 + \lambda ||x||_{TV},
$$
  
where  $||x||_{TV} = \sum_{i \sim j} |x_i - x_j|$ .



Figure: Original cameraman figure.



<span id="page-32-0"></span>Figure: Added white noise with standard deviation 25.

#### Image denoising

Image data  $b_i$  is available at two different sensors.



Figure: Original cameraman figure.

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Figure: Noisy image data in 2 parts.

#### [Simulations](#page-34-0)

#### Image denoising

Recover the original image by solving the following problem:

$$
\min_{x} \frac{1}{2} ||x - b_1||_2^2 + \frac{1}{2} ||x - b_2||_2^2 + \lambda ||x||_{TV},
$$

with asynchronous ADMM algorithm with 3 agents. Algorithm converged after 87 iterations, 35 seconds on laptop.



Figure: Original cameraman figure.



Figure: Noisy image data in 2 parts.

<span id="page-34-0"></span>Figure: Recovered using total variation denoising formula with  $\lambda = 20$ .

Recovered

### Conclusions and Future Work

- For general convex problems, we developed an asynchronous distributed ADMM algorithm, which converges at the best known rate  $O(1/k)$ .
- **•** Simulation results illustrate the superior performance of ADMM (even for network topologies with slow mixing).
- <span id="page-35-0"></span>**O** Ongoing and Future Work:
	- Online and dynamic distributed optimization problems.
	- ADMM type algorithm for time-varying graph topology.
	- Analyze network effects on ADMM algorithm.