

Extended formulations (II): semidefinite programming lifts

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Semidefinite programming extended formulations

\mathbf{S}_+^d = cone of $d \times d$ positive semidefinite matrices

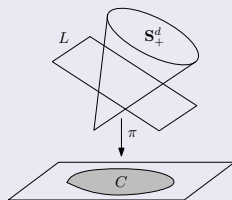
Definition

Let C be a convex set. We say that C has an **SDP extended formulation** (or SDP lift) of size d if we can write

$$C = \pi(\mathbf{S}_+^d \cap L)$$

where

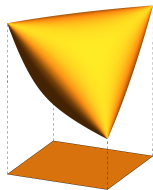
- π is a linear map;
- and L a linear subspace of \mathbf{S}^d



$x_{SDP}(C)$ = smallest size of an SDP lift of C

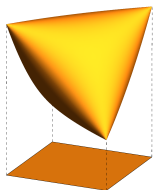
Examples of SDP lift

$$[-1, 1]^2 = \left\{ (x, y) \in \mathbb{R}^2 : \exists u \in \mathbb{R} \text{ s.t. } \begin{bmatrix} 1 & x & y \\ x & 1 & u \\ y & u & 1 \end{bmatrix} \succeq 0 \right\}$$



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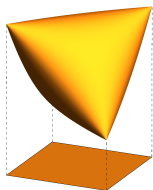


- Nuclear norm ball (cf. Pablo's talk)

$$\left\{ M \in \mathbb{R}^{n \times m} : \exists X, Y \text{ s.t. } \begin{bmatrix} X & M \\ M^T & Y \end{bmatrix} \succeq 0 \right. \\ \left. \frac{1}{2}(\text{Tr}(X) + \text{Tr}(Y)) \leq 1 \right\}$$

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- $STAB(G)$ for perfect graph G (cf. Michel's talk)

Existential question

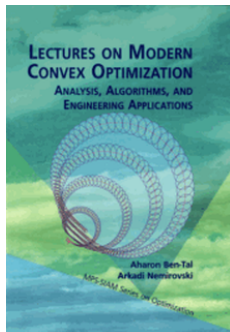
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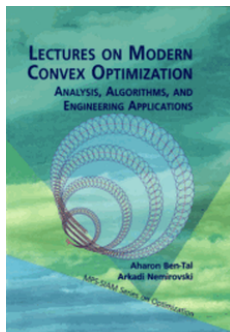
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Gives SDP representations for various convex sets and functions:

- ℓ_p norm balls for $p \geq 1$ rational
- Nuclear norm / Schatten ℓ_p norms
- Sum of k largest eigenvalues/singular values
- ...

Implemented in modeling tools like CVX and Yalmip

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- Scheiderer (2016): there are (many) convex semialgebraic sets that **do not** have an SDP representation

Slack matrix

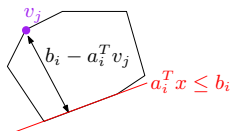
P polytope in \mathbb{R}^n

Slack matrix of P : Nonnegative matrix M of size $\#\text{facets}(P) \times \#\text{vertices}(P)$:

$$M_{i,j} = b_i - a_i^T v_j$$

where

- $a_i^T x \leq b_i$ are the facet inequalities of P
- v_j are the vertices of P



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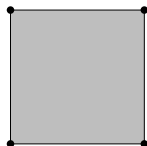
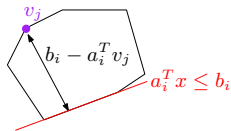
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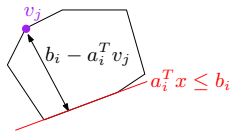
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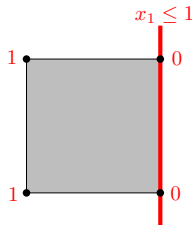
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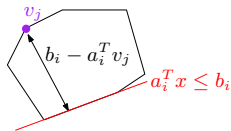
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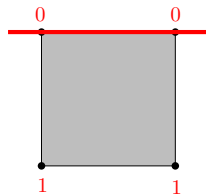
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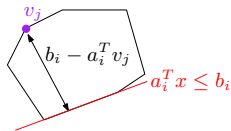
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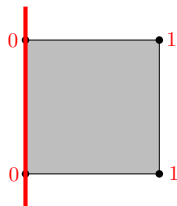
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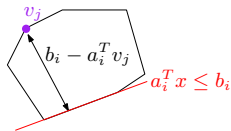
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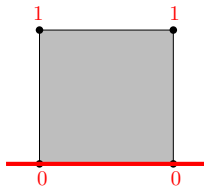
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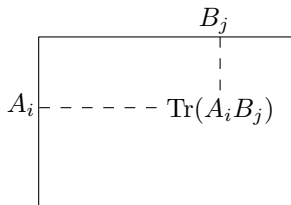
Positive semidefinite rank

$M \in \mathbb{R}^{p \times q}$ with nonnegative entries

- Positive semidefinite factorization:

$$M_{ij} = \langle A_i, B_j \rangle \quad \text{where} \quad A_i, B_j \in \mathbf{S}_+^d$$

- $\text{rank}_{\text{psd}}(M)$ = size of smallest psd factorization



Example of positive semidefinite factorization

Consider $M_{ij} = (i - j)^2$ for $1 \leq i, j \leq n$:

$$M = \begin{bmatrix} 0 & 1 & 4 & 9 & 16 \\ 1 & 0 & 1 & 4 & 9 \\ 4 & 1 & 0 & 1 & 4 \\ 9 & 4 & 1 & 0 & 1 \\ 16 & 9 & 4 & 1 & 0 \end{bmatrix}$$

- $\text{rank}_{\text{psd}}(M) = 2$ (independent of n): Let

$$A_i = \begin{bmatrix} 1 & i \\ i & i^2 \end{bmatrix} = \begin{bmatrix} 1 \\ i \end{bmatrix} \begin{bmatrix} 1 \\ i \end{bmatrix}^T \quad \text{and} \quad B_j = \begin{bmatrix} j^2 & -j \\ -j & 1 \end{bmatrix} = \begin{bmatrix} -j \\ 1 \end{bmatrix} \begin{bmatrix} -j \\ 1 \end{bmatrix}^T.$$

One can verify that $M_{ij} = \text{Tr}(A_i B_j)$.

SDP lifts and PSD rank

Theorem

Let P be a polytope with slack matrix M . Then $\chi_{\text{SDP}}(P) = \text{rank}_{\text{psd}}(M)$.

Connection with sums of squares

Theorem

Let $P = \text{conv}(X)$ be a polytope.

- If P has a SDP lift of size d , then there exists a subspace \mathcal{V} of \mathbb{R}^X such that the following holds:
 - (i) $\dim \mathcal{V} \leq d^2$
 - (ii) Any facet $b - a^T x \geq 0$ of P has a s.o.s. certificate from \mathcal{V} i.e., there exist $h_\alpha \in \mathcal{V}$ s.t.

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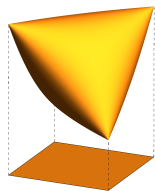
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Lasserre hierarchy: $\mathcal{V} =$ subspace of **polynomials of degree at most k**

Connection with sums of squares: example 1

- $X = \{-1, 1\}^2$, $\text{conv}(X) = [-1, 1]^2$.
- Four facet inequalities:

$$1 - x_1 \geq 0, \quad 1 + x_1 \geq 0, \quad 1 - x_2 \geq 0, \quad 1 + x_2 \geq 0$$



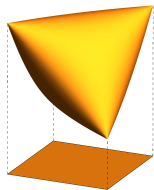
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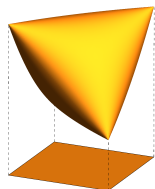
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- Similar certificate holds for the other facets
- Subspace

$$\mathcal{V} = \text{span}(1, x_1, x_2)$$

has dimension 3.

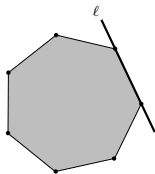
This yields an SDP lift of $[-1, 1]^2$ of size 3.



Example 2: polygons

- $P = \text{conv}(X)$ with $X = N$ roots of unity
- Facet inequality

$$\ell(x, y) = \cos(\pi/N) - \cos(\pi/N)x - \sin(\pi/N)y \geq 0$$

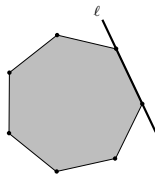


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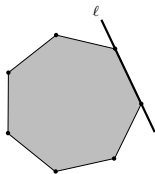
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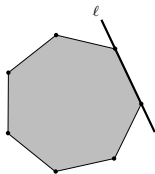
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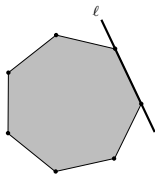
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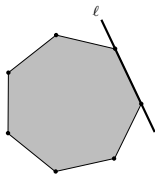
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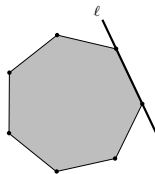
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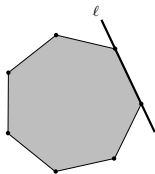
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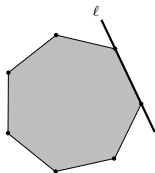
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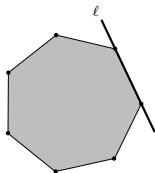
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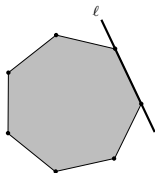
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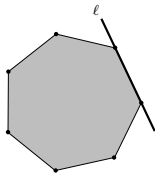
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$$V = \bigoplus_{i \in \mathcal{T}} V_i$$

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- $P = \text{conv}(X)$ with $X = N$ roots of unity
- Facet inequality

$$\ell(x, y) = \cos(\pi/N) - \cos(\pi/N)x - \sin(\pi/N)y \geq 0$$



- Any s.o.s. certificate of ℓ must have $\text{deg} \geq N/4$

$$\mathbb{R}^X = \boxed{\text{Pol}_0 \oplus \text{Pol}_1 \oplus \text{Pol}_2 \oplus \text{Pol}_3} \oplus \text{Pol}_4 \oplus \dots \oplus \text{Pol}_N$$

- Different choice of subspace?

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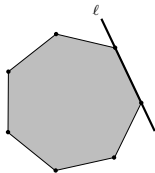
$$V = \bigoplus_{i \in \mathcal{T}} V_i$$

$$\ell = \sum_{j=0}^{n-2} \frac{\sin\left(\frac{\pi}{2^n}\right)}{2^j \sin\left(2^{j+1} \cdot \frac{\pi}{2^n}\right)} \left(\cos\left(\frac{\pi}{2^{n-j}}\right) c_0 - \cos\left(\frac{\pi}{2^{n-j}}\right) c_{2j} - \sin\left(\frac{\pi}{2^{n-j}}\right) s_{2j} \right)^2$$

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Subspace of dimension $\sim \log N$

Hierarchies and extended formulations

$P = \text{conv}(X)$. Certify nonnegativity of facets ℓ of P

	Hierarchies	Extended formulations
LP	$X = \{0, 1\}^n \cap \{g_i(x) \geq 0\}$ Handelman/Sherali-Adams $\ell = \text{nonneg. comb. of}$ $x^\alpha (1-x)^\beta \prod_i g_i(x)^{\gamma_i}$	$\ell = \text{nonneg. combinations of some}$ well-chosen $a_i : X \rightarrow \mathbb{R}_+$
SDP	$\ell = \text{s.o.s. of degree } \leq k$	$\ell = \text{s.o.s. from a well-chosen}$ subspace \mathcal{V}