

# Pseudorandomness and regularity in graphs I

David Conlon

January 17, 2017



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In this talk, we will focus on discussing the second and third cases.

# The dense case

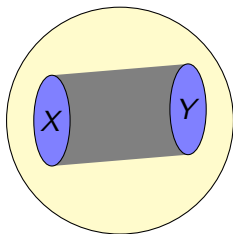
# The dense case

## $p$ -quasirandom sequence

A sequence of graphs  $(G_n)_{n \in \mathbb{N}}$  with  $|G_n| = n$  is said to be  $p$ -quasirandom if

$$|e(X, Y) - p|X||Y|| = o(pn^2)$$

for all  $X, Y \subseteq V(G_n)$ .





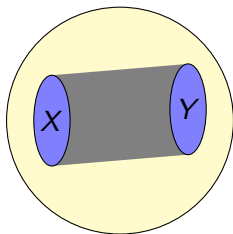
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Definition motivated by Szemerédi's regularity lemma.

# The adjacency matrix

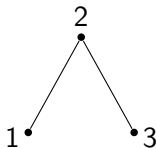
The adjacency matrix  $A$  of a graph  $G$  on vertex set  $\{1, 2, \dots, n\}$  is the  $n \times n$  matrix with entries given by

$$A_{uv} = \begin{cases} 0 & \text{if } uv \notin E(G); \\ 1 & \text{if } uv \in E(G). \end{cases}$$

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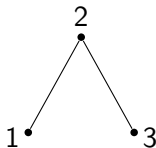


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Denote the eigenvalues of the adjacency matrix of  $A$  by  $\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_n$ .

# The Chung–Graham–Wilson theorem

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- (iv) the number of labelled copies of  $C_4$  is  $(1 + o(1))p^4n^4$ .



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## Expander Mixing Lemma

If  $G$  is a graph for which all eigenvalues of the adjacency matrix, save the largest, have absolute value at most  $\lambda$ , then

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Applying this lemma with  $\lambda = o(pn)$  gives the required result.

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Since  $\lambda_1 = (1 + o(1))pn$ , we must have  $\lambda_i = o(pn)$  for all  $i \neq 1$ .

## Counting Lemma

If  $G$  is a graph which satisfies DIS with  $p$  fixed, then, for any fixed graph  $H$ , the number of labelled copies of  $H$  in  $G$  is

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The required result is just the case where  $H = C_4$ .

To prove the counting lemma when  $H = K_3$ , write

$$e(S, T) = \sum_{s \in S, t \in T} 1_G(s, t),$$

where  $1_G$  is the indicator function for edges of  $G$ .

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$$\left| \sum_{s \in S, t \in T} (1_G(s, t) - p) \right| = o(n^2).$$

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In turn, this implies that for any functions  $u, v : V(G) \rightarrow [0, 1]$ ,

$$\left| \sum_{x, y \in V(G)} (1_G(x, y) - p)u(x)v(y) \right| = o(n^2).$$

This follows since the function we wish to optimise is linear in  $u(x)$  and  $v(y)$  for each  $x$  and  $y$ , so the maximum occurs when  $u$  and  $v$  are  $\{0, 1\}$ -valued, i.e., indicator functions.

By telescoping, the deviation between the number of labeled triangles in a set  $S \subseteq V(G)$  and its expected value is

$$\begin{aligned} & \sum_{x,y,z \in S} (1_G(x,y)1_G(y,z)1_G(z,x) - p^3) \\ &= \sum_{x,y,z \in S} (1_G(x,y) - p)1_G(y,z)1_G(z,x) + \\ & \quad \sum_{x,y,z \in S} p(1_G(y,z) - p)1_G(z,x) + \sum_{x,y,z \in S} p^2(1_G(z,x) - p). \end{aligned}$$

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Each term on the right-hand side of this equation may be written as a sum over terms of the form  $\sum_{x,y \in V(G)} (1_G(x,y) - p)u(x)v(y)$  for some appropriate  $u$  and  $v$ , thus implying that the deviation we are interested in is  $o(n^3)$ .

# Forcing graphs

## Forcing graph

A graph  $H$  is said to be forcing if any sequence of graphs  $(G_n)_{n \in \mathbb{N}}$  of density  $p$  with  $|G_n| = n$  containing  $(1 + o(1))p^{e(H)}n^{v(H)}$  labelled copies of  $H$  is  $p$ -quasirandom.

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$H$  is not forcing if it is

- non-bipartite or
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Progress on this conjecture closely parallels the progress on Sidorenko's conjecture (Conlon–Fox–Sudakov, Szegedy–Li, Kim–Lee–Lee, Szegedy, Conlon–Kim–Lee–Lee).

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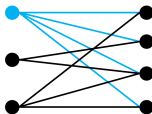
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## Example - Conlon–Fox–Sudakov; Szegedy–Li

Every bipartite graph which is not a tree and has a vertex which is complete to the other side is forcing.





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**HER** for every  $S \subseteq V(G_n)$ , the number of labelled triangles in  $G_n[S]$  is equal to  $p^3|S|^3 + o(p^3n^3)$ .

# Hereditary quasirandomness

## Theorem (Simonovits–Sós, 1997)

For any  $0 < p < 1$  and any  $\epsilon > 0$ , there exists  $\delta > 0$  such that if the number of labelled triangles in  $G[S]$  is equal to  $p^3|S|^3 \pm \delta n^3$ , then

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An alternative regularity-free proof given by Reiher and Schacht.

## Ramsey's theorem

For any graph  $H$ , there exists a natural number  $n$  such that if the edges of the complete graph  $K_n$  are two-coloured, there is always a monochromatic copy of  $H$ .

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## Bounds for complete graphs

$$\sqrt{2}^t \leq r(K_t) \leq 4^t$$



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Theorem (Erdős–Szekeres, 1935)

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- Show that this contradicts another property of the colouring.

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NO: DIS  $\not\Rightarrow$  EIG

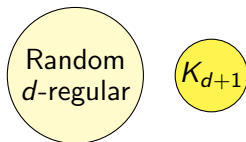
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# Some positive results

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## Theorem (Alon–Coja-Oghlan–Hàn–Kang–Rödl–Schacht, 2010)

Suppose that  $(G_n)_{n \in \mathbb{N}}$  with  $|G_n| = n$  is a sequence of graphs such that

$$|e(X, Y) - p|X||Y|| = o(pn^2)$$

for all  $X, Y \subseteq V(G_n)$ . Then one may remove a  $o(1)$ -fraction of the vertices to find a sequence of graphs  $(G'_n)_{n \in \mathbb{N}}$  satisfying EIG.

## Cayley graph

Suppose that  $G$  is a group and  $S$  is a subset of  $G$  satisfying  $S = S^{-1}$ . The Cayley graph  $\text{Cay}(G, S)$  is the graph with vertex set  $G$  and edge set  $\{(sg, g) : g \in G, s \in S\}$ .

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- Paley graph -  $G = \mathbb{Z}_p$ ,  $S = \{x^2 : x \in \mathbb{Z}_p \setminus \{0\}\}$

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- Lubotzky–Phillips–Sarnak -  $G = PSL(2, q)$ ,  $S = \dots$



# DIS $\Rightarrow$ EIG in Cayley graphs

Theorem (Kohayakawa–Rödl–Schacht, 2016+)

If  $G$  is an abelian group and  $\text{Cay}(G, S)$  satisfies

$$|e(X, Y) - \frac{d}{n}|X||Y|| \leq \epsilon dn$$

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## Jumbledness

A graph  $G$  on vertex set  $V$  is  $(p, \beta)$ -jumbled if, for all vertex subsets  $X, Y \subseteq V(G)$ ,

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## Example - random graphs

Let  $p = p(n) \leq 0.99$ . Then, asymptotically almost surely, the binomial random graph  $G(n, p)$  has the following property. For any two subsets  $X, Y \subseteq V(G)$ ,

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This is close to best possible, in that we must have  $\beta = \Omega(\sqrt{pn})$ .

# $(n, d, \lambda)$ -graphs

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A graph  $G$  is said to be an  $(n, d, \lambda)$ -graph if it has  $n$  vertices, every vertex has degree  $d$  and, if  $\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_n$  are the eigenvalues of the adjacency matrix of  $G$ ,  $|\lambda_i| \leq \lambda$  for all  $i \geq 2$ .

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By the expander mixing lemma,

$$\left| e(X, Y) - \frac{d|X||Y|}{n} \right| \leq \lambda \sqrt{|X||Y|},$$

so that  $(n, d, \lambda)$ -graphs are  $(d/n, \lambda)$ -jumbled.



# Strongly regular graphs

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A strongly regular graph  $srg(n, d, \eta, \mu)$  is a  $d$ -regular graph on  $n$  vertices in which every pair of adjacent vertices have exactly  $\eta$  common neighbours and every pair of nonadjacent vertices have exactly  $\mu$  common neighbours.

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## Paley graphs

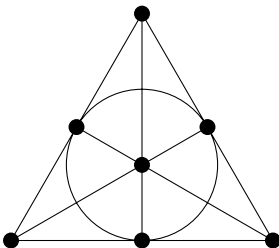
Suppose  $q \equiv 1 \pmod{4}$  is prime. The Paley graph  $P_q$  is the graph with vertex set  $\mathbb{Z}_q$ , where  $x$  and  $y$  are joined if and only if  $x - y$  is a quadratic residue. This graph is strongly regular, with parameters  $(q, (q - 1)/2, (q - 5)/4, (q - 1)/4)$ , and so is an  $(n, d, \lambda)$ -graph with  $n = q$ ,  $d = (q - 1)/2$  and  $\lambda = O(\sqrt{q})$ .

## Further examples

Let  $q$  be a prime power and  $PG(q, t)$  the projective space of dimension  $t$ , that is, each element is an equivalence class of non-zero vectors of length  $t + 1$  over the finite field of order  $q$ , where two vectors are taken as equivalent if one is a multiple of the other by an element in the field. This set has  $n = (q^{t+1} - 1)/(q - 1)$  elements.

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We define a graph  $G$  whose vertices are the points of  $PG(q, t)$  and where two vertices  $x = (x_0, x_1, \dots, x_t)$  and  $y = (y_0, y_1, \dots, y_t)$  are adjacent if and only if

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It is straightforward to check that  $G$  is  $d$ -regular with  $d = (q^t - 1)/(q - 1)$ , though there may be some loops. However, there are  $O(q^{t-1})$  vertices with loops. To calculate the eigenvalues of  $G$ , let  $I$  be the  $n \times n$  identity matrix and  $J$  the  $n \times n$  all-one matrix. Then

$$A^2 = \mu J + (d - \mu)I,$$

where  $\mu = (q^{t-1} - 1)/(q - 1)$ . This easily implies that all eigenvalues apart from the largest have absolute value  $\sqrt{d - \mu}$ .

## Theorem (Alon, 1994)

There is a triangle-free  $(n, d, \lambda)$ -graph with  $n = 2^{3k}$ ,  
 $d = 2^{k-1}(2^{k-1} - 1)$  and  $\lambda = O(2^k) = O(\sqrt{d})$ .



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Note that this has  $d = \Omega(n^{2/3})$ , which is significantly denser than the triangle-free graphs obtained from modifications of random graphs, which have  $d = \tilde{\Omega}(\sqrt{n})$ .

# Some properties

## Lemma

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## Tightness

In the Paley graph  $P_q$ , where  $q = p^2$  is the square of a prime, every element of the subfield  $\text{GF}(p)$  is a quadratic residue in  $\text{GF}(p^2)$  since  $\text{GF}(p^2)$  is a splitting field for  $X^2 - a$  for all  $a \in \text{GF}(p)$ . To form an independent set of order  $\sqrt{q}$ , consider  $\beta\text{GF}(p)$ , where  $\beta$  is a quadratic non-residue.

## Lemma

The chromatic number  $\chi(G)$  of a  $(p, \beta)$ -jumbled graph  $G$  with  $n$  vertices satisfies

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The Paley graph  $P_q$  with  $q = p^2$  again shows that this is tight.

## Lemma

For any  $\epsilon > 0$  and any graph  $H$  with maximum degree  $\Delta$ , there exists  $c > 0$  such that if  $G$  is a  $(p, \beta)$ -jumbled graph with  $\beta \leq cp^\Delta n$ , then the number of labelled copies of  $H$  in  $G$  is equal to

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For triangles, Alon's example shows that this is tight.



## Problem

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The MAXCUT of an  $m$ -edge triangle-free graph is  $m/2 + \Omega(m^{4/5})$ .

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## Theorem (Alon, 1996)

The MAXCUT of an  $m$ -edge triangle-free graph is  $m/2 + \Omega(m^{4/5})$ .

## Tightness

The MAXCUT of an  $(n, d, \lambda)$ -graph with  $m = dn/2$  edges is at most  $m/2 - \lambda_n n/4$ . In particular, the MAXCUT of Alon's example is  $m/2 + O(m^{4/5})$ .

Thank you for listening!