

Sparse Random Graphs-I

Souvik Dhara

Graph Limits and Processes on Networks:
From Epidemics to Misinformation Boot Camp

Lets start with a few questions...



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What causes **INTERNET** to **breakdown**?

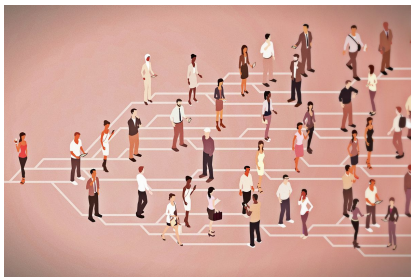
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When does **MISINFORMATION** reach a **large population**?

What are Random Graphs for?

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➤ Random Graphs provide a *simplified probabilistic representation* to model these complex system.

- ➔ Capture structural properties (degree distribution, communities)
- ➔ Provide insight into emergence of different types of behavior such as phase transition

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➤ Random Graphs serve to get provable **guarantees for graph algorithms**

Example: Heuristic algorithms for NP-hard problems such as graph partitioning, coloring

Plan

Today:

- Local Branching Process approximation technique on random graphs
- Explore its relation to *Giant Component Problem* on different models

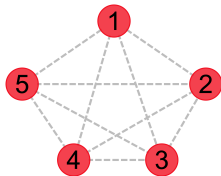
Tomorrow:

- Applications to Percolation, Epidemics
- Using Stochastic Process convergence in Random Graphs

Let's start with the most elementary yet fundamental model...

Erdős-Rényi Random Graph

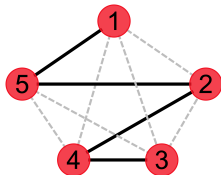
Erdős-Rényi Random Graph



Definition

- Given n nodes $\{1, 2, \dots, n\}$

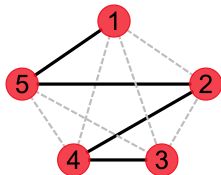
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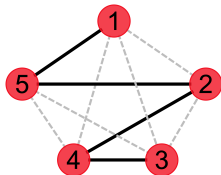
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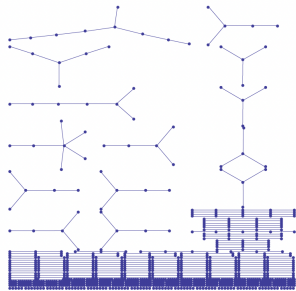
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Historical note:

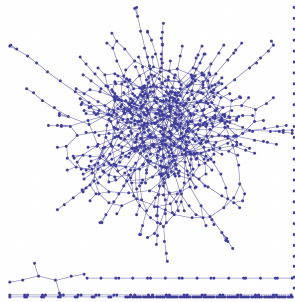
- This model was actually studied by Gilbert (1959) and heuristically by Solomonoff & Rapoport (1951)
- Erdős & Rényi (1959) initially worked with a slightly different model where fixed number of edges sampled uniformly. In a sequence of eight papers between 1959-1968 they laid the foundation of Random Graph theory

What are we after?

$ER_n(p)$ with $n = 1000$



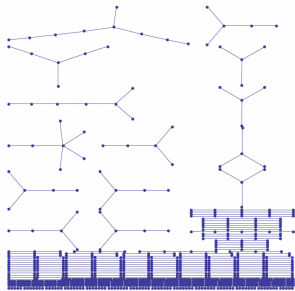
$$p = \frac{0.5}{n}$$



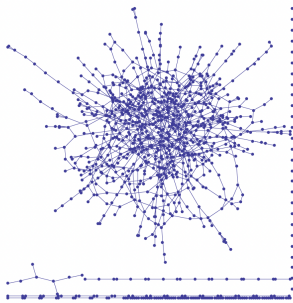
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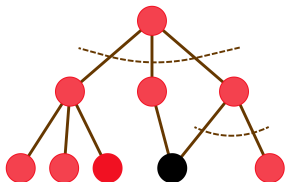
- If $p = \frac{\lambda}{n}$, then there is *phase transition around $\lambda = 1$*
- ➔ $\lambda < 1$: All components are small
 - ➔ $\lambda > 1$: There is a unique *giant component*

Local neighborhood structure of $ER_n(\frac{\lambda}{n})$

- **To analyze component sizes:** Gradually explore graph in BFS starting from any node, e.g., node 1

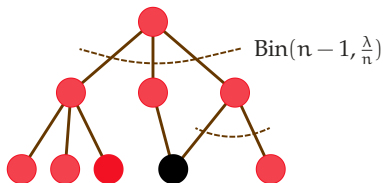
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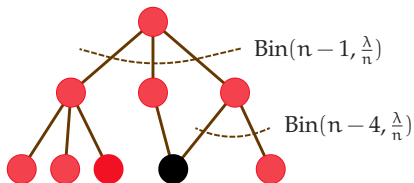
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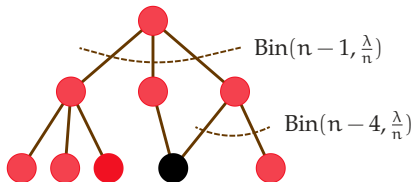
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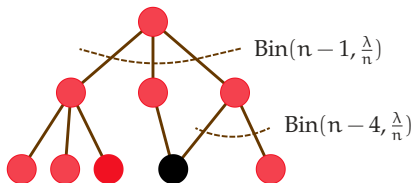
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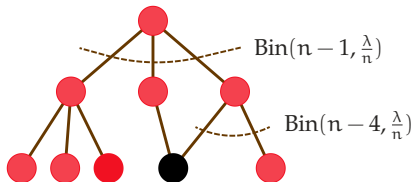


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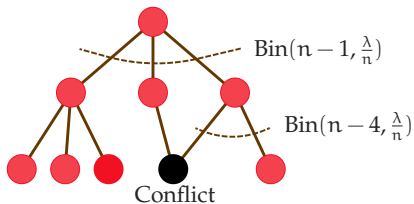
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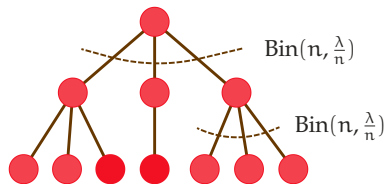
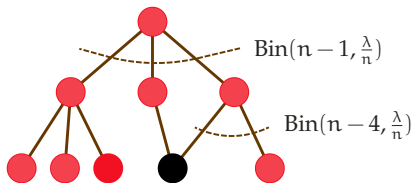
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- ➔ Depletion of vertices
- ➔ Conflicts among new vertices

Domination by Branching Process

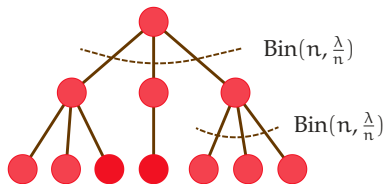
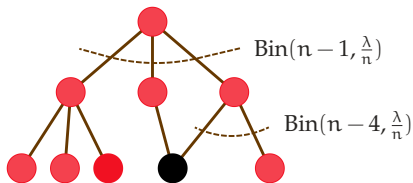
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Simple Fact

Let $N_k = \#$ vertices at depth k for $\text{ER}_n(\frac{\lambda}{n})$ exploration, and \bar{N}_k denotes same for Branching process. There is a coupling such that w.p. 1

$$N_k \leq \bar{N}_k \quad \forall k \geq 0$$

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Theorem: Subcritical $ER_n(\frac{\lambda}{n})$

If $\lambda < 1$, then

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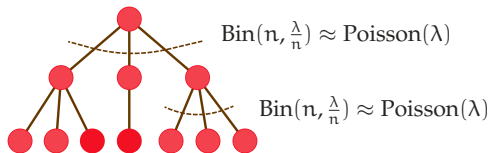
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➤ Proof uses **Large Deviation** estimates for branching process survival prob

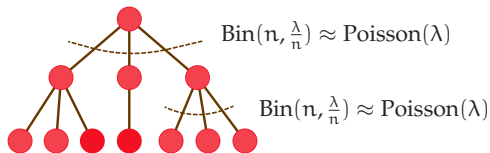
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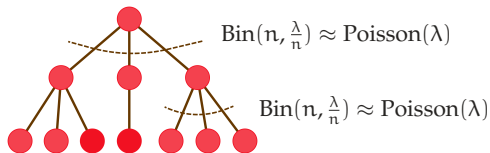
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➤ ζ_λ satisfies is a positive solution of $1 - \zeta = e^{-\lambda\zeta}$

Existence of a Giant for $\lambda > 1$

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Theorem: Supercritical $\text{ER}_n(\frac{\lambda}{n})$

Let $C_{(i)} := i$ -th largest component of $\text{ER}_n(\frac{\lambda}{n})$. If $\lambda > 1$, then as $n \rightarrow \infty$

$$\frac{C_{(1)}}{n} \xrightarrow{\mathbb{P}} \zeta_\lambda \quad \text{and} \quad \frac{C_{(2)}}{n} \xrightarrow{\mathbb{P}} 0$$

A unique giant component emerges...

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Next, prove two lemmas but before that...

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Will show: $N_k \approx \bar{N}_k$ until s vertices explored for $s = n^a$, $a < 1$

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➤ If k is a *large constant*, then $\mathbb{E}[\# \text{ conflicts}] \approx 0$ *Exploration = BP w.p. ≈ 1*

➤ Let $k = a \log_{\lambda} n$ and $0 \leq a < 1$. Then $\lambda^k = n^a$ and

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$$\# \text{ conflicts} = o(\bar{N}_k) \text{ whp} \implies \boxed{N_k \approx \bar{N}_k \approx \lambda^k}$$

When exploration survives for long time, growth rate of N_k becomes exponential in λ

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Next, lets prove two lemmas

Proving first moment lemma

To prove: $\sum_{i \geq 1} \frac{C_{(i)}}{n} \mathbb{1}\{C_{(i)} \geq L\} \approx \zeta_\lambda$

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Key fact 1: Local neighborhood of \mathbf{u} is approximately BP whp

➤ *Theory of approximating local neighborhood of graphs is called **Local-weak convergence** (Christian's talk)*

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$$= \sum_{i \geq 1} \frac{C_{(i)}^2}{n^2} \mathbb{1}\{C_{(i)} \geq L\} + \sum_{i \neq j} \frac{C_{(i)} C_{(j)}}{n^2} \mathbb{1}\{C_{(i)} \geq L, C_{(j)} \geq L\}$$

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Enough to show:

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Two large components cannot be disjoint...

Two large components cannot be disjoint, why?

To show $\mathbb{P}(C(u_1) \geq L, C(u_2) \geq L, u_1 \not\leftrightarrow u_2) \approx 0$, suffices to prove

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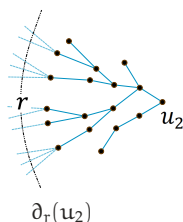
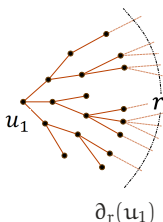
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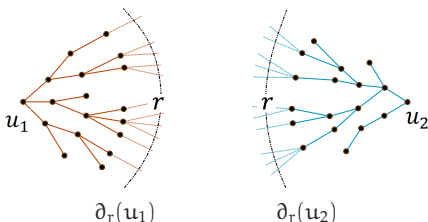
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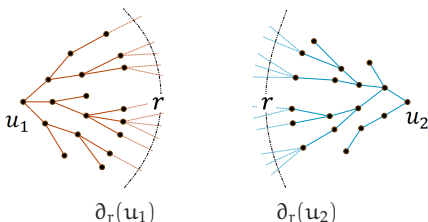
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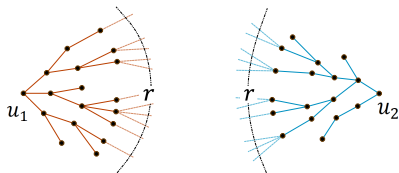
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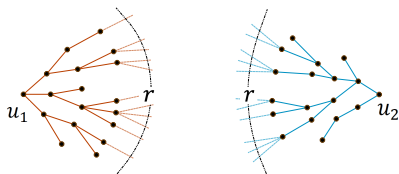
➤ Advantage with conditioning on $\{\partial_r(u_1), \partial_r(u_2) \neq \emptyset\}$ is we can now explore rest of the graph

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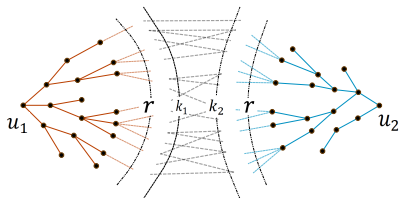
➤ Condition on $\{\partial_r(u_1), \partial_r(u_2) \neq \emptyset\}$ for large r

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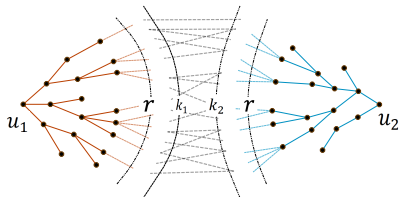
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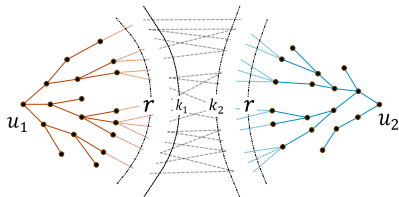
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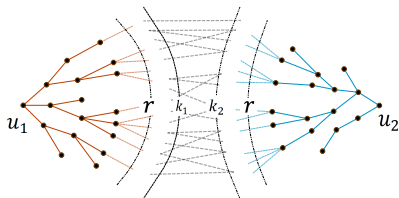
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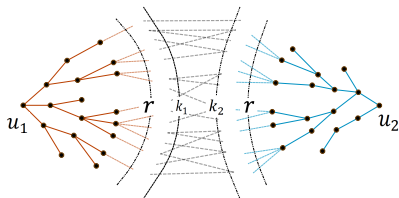


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Takeaway: If there are two components with large boundary, we can grow them until boundary has size \sqrt{n} and then they intersect

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Key fact 2: $\lim_{L \rightarrow \infty} \limsup_{n \rightarrow \infty} \mathbb{P}(C(u_1) \geq L, C(u_2) \geq L, u_1 \not\leftrightarrow u_2) = 0$

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Summary: Giant for Erdős-Rényi

We proved

Theorem: Supercritical $ER_n(\frac{\lambda}{n})$

Let $C_{(i)} := i$ -th largest component of $ER_n(\frac{\lambda}{n})$. If $\lambda > 1$, then

$$\frac{C_{(1)}}{n} \xrightarrow{\mathbb{P}} \zeta_\lambda \quad \text{and} \quad \frac{C_{(2)}}{n} \xrightarrow{\mathbb{P}} 0$$

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➤ van der Hofstad (2021) proved this for general graphs that converge in local-weak convergence sense

Before moving on to other models, lets see another useful application of the above ideas...

Typical distances in Erdős-Rényi

Typical distances in the giant of Erdős-Rényi

Typical distance: Graph distance between two uniform vertices u_1, u_2

Typical distances in the giant of Erdős-Rényi

Typical distance: Graph distance between two uniform vertices u_1, u_2

Theorem: Typical distances in $ER_n(\frac{\lambda}{n})$

Let $\lambda > 1$. Conditionally on u_1, u_2 in same component (i.e., $\text{dist}(u_1, u_2) \neq \infty$)

$$\frac{\text{dist}(u_1, u_2)}{\log_{\lambda} n} \xrightarrow{\mathbb{P}} 1$$

Typical distances in the giant of Erdős-Rényi

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