

Identifying Mixture Models

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Based on joint works with Chaitanya Swamy (Waterloo), Yuval Rabani (Hebrew U), Jian Li (Tsinghua), Spencer Gordon (Caltech) and Bijan Mazaheri (Caltech)

We are interested in Bayesian Networks $\mathcal{G} = (\mathcal{V} \sqcup \mathcal{U}, \mathcal{E})$ with **visible** vertices \mathcal{V} and **hidden** or **latent** confounders \mathcal{U} .

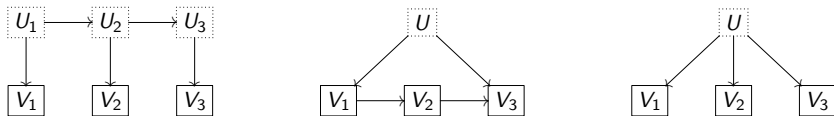


Figure 1: Hidden Markov Model; “Front Door”; MixProd

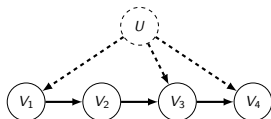
Write \mathcal{P} for the restriction of the joint distribution to \mathcal{V} . This is what we can learn (up to sampling noise) from data. \mathcal{P} is Markovian on the graph: factors as

$$\Pr(v_1, \dots, v_n) = \prod_{i=1}^n \Pr(V_i = v_i \mid \mathbf{pa}(V_i))$$

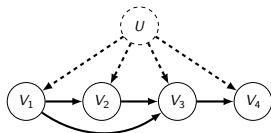
where $\mathbf{pa}(V_i)$ is the assignment to the parents of V_i . These conditionals are the **parameters** of the model.

Source Identification / Parameter learning

If we're given the distribution on all variables (including U), we can easily identify all the parameters of the model. But we're actually only given \mathcal{P} (or empirical $\hat{\mathcal{P}}$). So what *can* we determine? In some cases [Pearl/ Tian/ Shpitser/ Huang/ Valtorta] can make remarkable deductions. E.g., in:



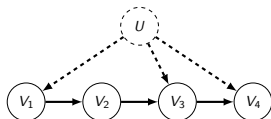
can deduce effect of an intervention at V_1 on V_3 , despite confounder U .
But in most cases, there's little we can determine from \mathcal{P} . E.g., if single U can affect all visible variables:



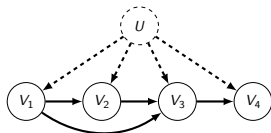
U can generate **any** distribution on \mathcal{V} .

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U can generate **any** distribution on \mathcal{V} . **But in order to do so, U needs to range over a large set.** (Size 2^n for binary V_i 's.)

Cardinality/dimension bounds on hidden variables

Cardinality or dimension bounds on hidden variables are a long-standing assumption e.g., for Hidden Markov Models. In our “non-parametric” context, natural assumption is cardinality.

$$k = \text{cardinality}(\text{range}(U))$$

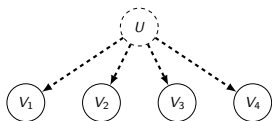


Figure 2: k -MixProd

Tower of increasingly general problems:

$$k\text{-MixIID} < k\text{-MixProd} < k\text{-MixBND}$$

In k -MixProd, the V_i are independent conditional on U .

In k -MixIID, they are moreover iid conditional on U .

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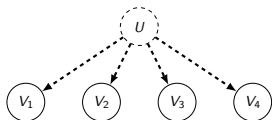


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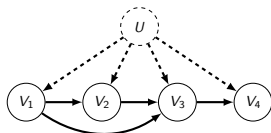


Figure 3: More general k -MixBND.

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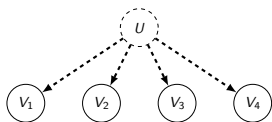


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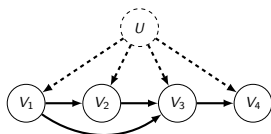


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Tower of increasingly general problems:

$$\underbrace{k\text{-MixIID} < k\text{-MixProd}}_{\text{this lecture}} < k\text{-MixBND}$$

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Talk outline

① k -MixIID and the classical moment problem.

Key concepts:

- ① Prony's algorithm.
- ② Hankel matrices.

Theorems: Hankel condition number, sample complexity

lower bound $\sim \exp(\Omega(k))$.

(And upper bounds also for transportation distance reconstruction.)

② k -MixProd.

Key concepts:

- ① Method of synthetic bits.
- ② Hadamard Extensions.

Theorems: Hadamard Extension condition number, sample complexity

upper bound $\sim \exp O(k \log k)$.

1. k -MixIID and the classical moment problem

In k -MixIID: U is distributed on $\{1, \dots, k\}$ according to an unknown prob. dist. π . Focus on case that V_i 's are all binary.

For each $u \in \{1, \dots, k\}$ there is an $0 \leq \mathbf{m}_u \leq 1$ s.t.

$$\Pr(V_i | U = u) = \mathbf{m}_u$$

so by conditional independence of the V_i , $V_R = \bigwedge_{i \in R} V_i$

$$\Pr(V_R = 1 | U = u) = \mathbf{m}_u^{|R|}.$$

So for the rv $Y = \# \text{ Heads}$

$$Y = |\{j : V_j = 1\}|$$

the moments of Y are linear combinations of the moments of the “ k -spike” atomic probability distribution on $[0, 1]$

$$p = \sum_{u=1}^k \pi_u \delta_{\mathbf{m}_u} \quad (\text{here } \delta_x \text{ is unit measure at } x)$$

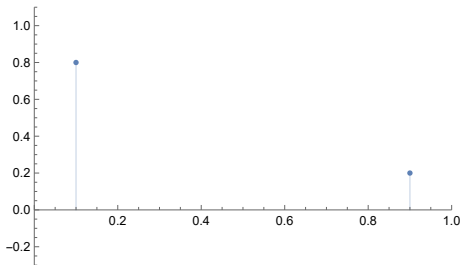


Figure 4: 2-spike dist. p with $\pi_1 = 0.8$ at $\mathbf{m}_1 = 0.1$, and $\pi_2 = 0.2$ at $\mathbf{m}_2 = 0.9$

Let $X \sim p$ and let $\mu_j = E(X^j)$. Then

$$E(Y) = n \sum_u \pi_u \mathbf{m}_u = n\mu_1$$

$$E(Y^2) = n\mu_1 + n(n-1)\mu_2$$

$$E(Y^3) = n\mu_1 + 3n(n-1)\mu_2 + n(n-1)(n-2)\mu_3 \quad \dots \text{etc.}$$

Triangular linear system with nonzero diagonal coefficients. So the moments of Y (0 through n), which we learn from \mathcal{P} , determine the moments μ_j of the k -spike dist. p .

The Moment Problem

Classical question: given μ_j ($j \geq 0$), are they the moments of a measure on \mathbb{R} ?

Classical answer: yes iff for every $K \geq 1$, the **Hankel matrix**

$$H_K = \begin{pmatrix} \mu_0 & \mu_1 & \dots & \mu_{K-1} \\ \mu_1 & \mu_2 & \dots & \mu_K \\ \dots & \dots & \dots & \dots \\ \mu_{K-1} & \mu_K & \dots & \mu_{2K-2} \end{pmatrix}$$

is nonnegative-definite.

Furthermore, the measure is **unique** provided the μ_j do not grow too quickly. For a distribution supported on $[0, 1]$ (Hausdorff moment problem), such as p , this is guaranteed.

For a k -spike distribution p , (1) How many moments are required to identify p , (2) How do we do so algorithmically?

Answers:

(1) μ_1, \dots, μ_{2k-1} suffice. (Easy to see necessary.)

(And can verify dist. is k -spike if we're also given μ_{2k} .)

Consequently sufficient to have $n = 2k - 1$ observable rv's.

(2) [Algorithm of Prony \(1795\)](#). Relies on the Hankel matrix which for k -spike dists is:

$$H_{k+1} = V_{k+1}^\perp \cdot \text{diag}(\pi) \cdot V_{k+1} \quad (1)$$

where V_ℓ is the $k \times \ell$ Vandermonde matrix of the spike sites:

$$V = \begin{pmatrix} 1 & \mathbf{m}_1 & \mathbf{m}_1^2 & \dots & \mathbf{m}_1^{\ell-1} \\ \dots & \dots & \dots & \dots & \dots \\ 1 & \mathbf{m}_k & \mathbf{m}_k^2 & \dots & \mathbf{m}_k^{\ell-1} \end{pmatrix}$$

$$\text{diag}(\pi) = \begin{pmatrix} \pi_1 & 0 & 0 \\ 0 & \dots & 0 \\ 0 & 0 & \pi_k \end{pmatrix}$$

This solves k -MixIID if you have perfect statistics, i.e., exact H_{k+1} .
“Living in Asymptotia”

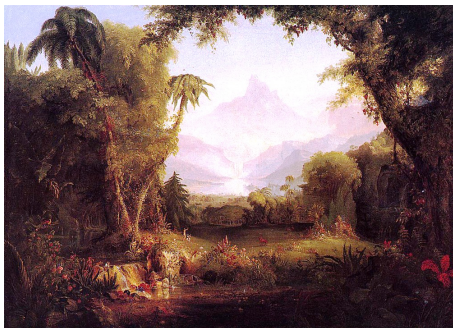


Figure 5: Thomas Cole, The Garden of Eden, 1828

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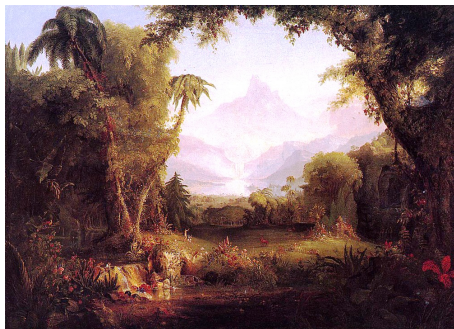


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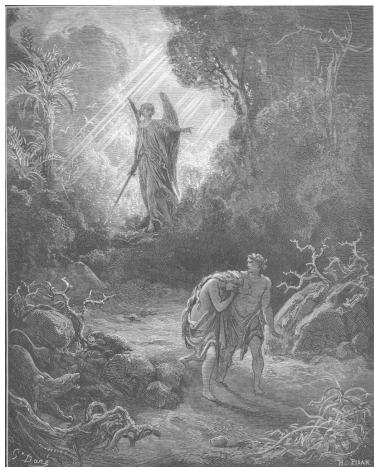


Figure 6: Gustave Doré, Adam and Eve Driven out of Eden, 1865

Sample size

Prony's alg. notoriously unstable as a function of empirical dist. $\hat{\mathcal{P}}$. Is this a property of the algorithm or of the problem? When spikes collide, model parameters not identifiable, so accuracy of $\hat{\mathcal{P}}$ (hence sample size) must depend on: **separation parameter**

$$\zeta = \min_{i \neq j} |\mathbf{m}_i - \mathbf{m}_j|.$$

Theorem 1 (Rabani S Swamy '14)

*For any $n \in O(k)$, $\|\hat{\mathcal{P}} - \mathcal{P}\|_\infty \leq \zeta^{O(k)}$ (therefore sample size $\geq (1/\zeta)^{\Omega(k)}$) is **necessary** even to determine parameters within $\pm 1/k$. (Neglecting dependence on mixture weights.)*

That paper also gave sample size upper bound of $(1/\zeta)^{O(k^2)}$. Since improved [Li Rabani S Swamy '15], [Kim, Koehler, Moitra, Mossel, Ramnarayan '19], [Gordon Mazaheri Rabani S '20] to $(1/\zeta)^{O(k)}$; also give reconstruction in Weierstrass-1 (transportation) distance. Key is an upper bound on condition number of Hankel H_k .

2. k -MixProd

Recall k -MixProd, a much more general problem than k -MixIID.

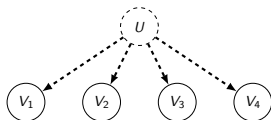


Figure 7: k -MixProd

Parameters: prior π on hidden variable U , and an $n \times k$ matrix

$$\mathbf{m}_{iu} = \Pr(V_i = 1 | U = u)$$

Prior work focused on **learning** rather than **identifying** the model.

“Learning” = reconstruct any model (π, \mathbf{m}) creating statistics close to the observed statistics.

“Identifying” = learning in regions of parameter space (π, \mathbf{m}) where there is a stable invertibility guarantee:

$$\forall \epsilon \exists \delta \text{ s.t. if } \text{dist}((\pi, \mathbf{m}), (\pi', \mathbf{m}')) > \delta \text{ then } |\mu(\pi, \mathbf{m}) - \mu(\pi', \mathbf{m}')| > \epsilon.$$

Identification gives stronger output guarantees than *Learning*, under stronger assumptions.

Identification as a goal goes back at least to [Koopmans, Reiersol 1950], [Koopmans 1950], [Teicher 1963], [Blischke 1964], [Yakowitz, Spragins 1968]

Since more is assumed, runtime might be better.

For our motivations, identification is the right problem, since it tells you how the system will function if you **intervene** (set some of the random variables).

Literature on k -MixProd

Mixture models began with [Newcomb 1886], [Pearson 1894]. See [Everitt, Hand 1981], [Titterton et al. 1985], [Lindsay 1995], [McLachlan et al. 2019]. Abundant literature for discrete variables thanks to disparate motivations, e.g., astronomy, population genetics, bioinformatics, image recognition, text classification; see [Pritchard et al. 2000], [Ji et al. '05], [Juan, Vidal '02, '04]. Iterative methods (EM) often used [Juan et al. '04], [Li et al. '16], [Palmer et al. '16], [Carrerira-Perpiñán, Renals '00], [Najafi et al. '20] ...

Algorithms with provable guarantees, some for Gaussians: $k = 2$: [Kearns et al. '94], [Freund, Mansour '99], [Dasgupta '99], [Cryan, Goldberg, Goldberg '02]. General k : [Feldman, O'Donnell, Servedio '08], [Chaudhuri, Rao '08], [Moitra, Valiant '10], [Arora et al. '12], [Anandkumar et al. '12ab], [Rabani et al. '14], [Hardt, Price '15], [Li et al. '15], [Kim et al. '19], [Chen, Moitra '19], [Wu, Yang '20], [Rabani et al. '20] ...

The provable “learning” algorithms use grid-and-search in parameter space. Due to grid search, this is very expensive: to learn a model which reproduces statistics within variation distance ε , [FOS'08] runtime is $(nk/\varepsilon)^{O(k^3)}$, [CM'19] $k^{O(k^3)}(n/\varepsilon)^{O(k^2)}$.

Identifying a k -MixProd model

We study k -MixProd under a ζ -separation assumption:

$$\forall i \forall u \neq u' : |\mathbf{m}_{iu} - \mathbf{m}_{iu'}| > \zeta > 0 \quad (2)$$

Comments:

(a) Separation for $\zeta = 0$ was shown by [Tahmasebi Motahari Maddah-Ali '18] to imply that the mapping $(\pi, \mathbf{m}) \rightarrow \mu$ is injective. (Provided $\pi > 0$, and up to the obvious symmetry of permuting columns.) Algebraic result, no algorithm.

(b) It is clear that some kind of separation guarantee is necessary: e.g., two identical columns make the model unidentifiable.

The separation assumption (2) is a little stronger than necessary. We provide a sufficient weaker assumption in [Gordon **S** '22]. However it is not algorithmic.

Full characterization and efficient algorithm beyond ζ -separation remain open problems.

We give an algorithm (different approach from [TMM'18] entirely) for ζ -separated k -MixProd. Near-optimal in sample complexity.

Theorem 2 (Gordon Mazaheri Rabani **S**, manuscript)

For a k -MixProd model on $n \geq 3k - 3$ bits, we can identify a model $(\hat{\pi}, \hat{\mathbf{m}})$ with all parameters within $\pm \varepsilon$ of true (π, \mathbf{m}) , in runtime and sample complexity

$$(1/\zeta)^{O(k \log k)} \varepsilon^{-2} n \log n.$$

What is the key challenge? In k -MixIID the observables V_1, \dots, V_n were iid conditional on the hidden variable U . So our **multilinear moments** (3)

$$E(V_1 V_2) = \sum_u \Pr(u) \Pr(V_1 = 1|u) \Pr(V_2 = 1|u) = \sum \pi_u \mathbf{m}_{1u} \mathbf{m}_{2u} \quad (3)$$

$$= \sum_u \pi_u \mathbf{m}_{1u}^2 \quad (4)$$

were actually higher moments (4) of the single k -spike distribution p .

But now each V_i has a unique dependence on U . What good does it do to combine information between different V_i ?

Hadamard extension of \mathbf{m}

Given two row vectors \mathbf{m}_1 and \mathbf{m}_2 in \mathbb{R}^k , their **Hadamard product** is

$$\mathbf{m}_1 \odot \mathbf{m}_2 \in \mathbb{R}^k$$
$$(\mathbf{m}_1 \odot \mathbf{m}_2)_u = \mathbf{m}_{1u} \mathbf{m}_{2u}$$

For $n \times k$ matrix \mathbf{m} , its **Hadamard Extension** is the $2^n \times k$ matrix with rows indexed by $S \subseteq [n]$, **multilinear version of Vandermonde**:

$$\mathbb{H}(\mathbf{m})_S = \bigodot_{i \in S} \mathbf{m}_i$$

or explicitly:
$$\mathbb{H}(\mathbf{m})_{S,u} = \prod_{i \in S} \mathbf{m}_{iu}$$

(appearing first, not with this name, in [Chen Moitra'19].) E.g.,

$$\mathbf{m} = \begin{pmatrix} 1/2 & 1/3 & 1/5 \\ 1/7 & 1/11 & 1/13 \end{pmatrix} \Rightarrow \mathbb{H}(\mathbf{m}) = \begin{pmatrix} 1 & 1 & 1 \\ 1/2 & 1/3 & 1/5 \\ 1/7 & 1/11 & 1/13 \\ 1/14 & 1/33 & 1/65 \end{pmatrix}$$

A complete list of observable statistics of our model is $\Pr(V_R)$, where $V_R = \bigwedge_{i \in R} V_i$, ranging over all $R \subseteq [n]$. These probabilities are given by the vector

$$\left(\Pr(V_R) \right) = \mathbb{H}(\mathbf{m}) \begin{pmatrix} \pi_1 \\ \dots \\ \pi_k \end{pmatrix}$$

Of course, we know only the vector on the LHS, not $\mathbb{H}(\mathbf{m})$ or π .

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It turns out that we will be able to use $\mathbb{H}(\mathbf{m})$ in our algorithm, *without* knowing it.

In order to be able to use $\mathbb{H}(\mathbf{m})$ at finite sample size, though, we *also* need to understand something about its numerical stability (not just rank). Discuss this first; later the algorithm.

Condition number of Hadamard extensions

Lemma 3

Let A be any set of $k - 1$ ζ -separated rows of \mathbf{m} . Write $\mathbf{m}|_A = \mathbf{m}$ restricted to the rows $i \in A$. Then the k 'th-largest singular value of $\mathbb{H}(\mathbf{m}|_A)$ satisfies:

$$\sigma_k(\mathbb{H}(\mathbf{m}|_A)) \geq \zeta^{O(k)}.$$

Effectively a far generalization of the eigenvalue lower bound for Hankel matrices; here Vandermonde \leftrightarrow Hadamard Extension. Clearest using

Lemma 4 (Feldman O'Donnell Servedio '08)

Let M be an $r \times k$ matrix, $r \geq k$. Then \exists a set J of k rows s.t.

$$\sigma_k(M|_J) \geq \frac{\sigma_k(M)}{\sqrt{k(r-k)+1}}.$$

Corollary 5

$\mathbb{H}(\mathbf{m}|_A)$ has a $k \times k$ submatrix \mathcal{A} with $\sigma_k(\mathcal{A}) \geq \zeta^{O(k)}$.

An alternative way of arranging the values $\Pr(V_R)$ is in the $2^n \times 2^n$ matrix

$$C = \mathbb{H}(\mathbf{m}) \text{diag}(\pi) \mathbb{H}(\mathbf{m})^\top$$

If $R \cap R' = \emptyset$ then $C_{R,R'} = \Pr(V_{R \cup R'})$. So we can observe some, **not all**, entries of this matrix. E.g., first column corresp. to \emptyset so fully observable:

$$(\mathbb{H}(\mathbf{m}) \text{diag}(\pi) \mathbf{1})_R = \Pr(V_R).$$

If for $A, B \subseteq [n]$, $A \cap B = \emptyset$ then we can observe the entire smaller matrix

$$\mathbb{H}(\mathbf{m} | _B) \text{diag}(\pi) \mathbb{H}(\mathbf{m} | _A)^\top. \quad \text{multilinear gen'l of Hankel matrix}$$

In particular if $A, B \subseteq [n]$, $|A| = |B| = k - 1$, $A \cap B = \emptyset$, then by Cor. 5, $\mathbb{H}(\mathbf{m} | _A)$ has a $k \times k$ submatrix \mathcal{A} , and $\mathbb{H}(\mathbf{m} | _B)$ has a $k \times k$ submatrix \mathcal{B} , such that we have good conditioning of the $k \times k$ matrix

$$C_{B,A} = \mathcal{B} \text{diag}(\pi) \mathcal{A}^\top$$

Lemma 6

$$\sigma_k(C_{B,A}) \geq \zeta^{O(k)}.$$

Reducing k -MixProd to k -MixIID: method of synthetic bits

Fix disjoint $A, B \subseteq [n]$ and well-conditioned $C_{B,A}$ as above. Let \mathbf{m}_1 be any row *outside of* $A \cup B$.

Strategy: we use the rows of B to **synthesize** a row equivalent to \mathbf{m}_1 ; we use A to determine the weights of this synthesis.

Recall that

$$E(V_1) = \mathbf{m}_1 \text{diag}(\pi) \mathbb{1}$$

We wish we had a variable V'_1 that was iid to V_1 conditional on U ; if so we'd be able to observe

$$E(V_1 V'_1) = (\mathbf{m}_1 \odot \mathbf{m}_1) \text{diag}(\pi) \mathbb{1}$$

We don't have such a V'_1 but the next-best thing is to construct $\mathbf{m}_1 \odot \mathbf{m}_1$. Concretely (marking in **violet** quantities we can compute):

(1) Let

$$v_1 := \mathbf{m}_1 \text{diag}(\pi) \mathcal{A}^\top$$

(We can observe v_1 because row 1 is not in A .)

In particular if $S = \emptyset$ is among the sets used in \mathcal{A} , then v_1 has an entry

$$(v_1)_\emptyset = \mathbf{m}_1 \text{diag}(\pi) \mathbb{1} = E(V_1).$$

(2) Let

$$u_1 := v_1 C_{\mathcal{B}, \mathcal{A}}^{-1}$$

u_1 is a set of weights that synthesize a copy of \mathbf{m}_1 out of \mathcal{B} :

$$\begin{aligned} u_1 \mathcal{B} &= [\mathbf{m}_1 \text{diag}(\pi) \mathcal{A}^\top] C_{\mathcal{B}, \mathcal{A}}^{-1} \mathcal{B} \\ &= \mathbf{m}_1 \text{diag}(\pi) \mathcal{A}^\top (\mathcal{A}^\top)^{-1} \text{diag}(\pi)^{-1} \mathcal{B}^{-1} \mathcal{B} \\ &= \mathbf{m}_1 \end{aligned}$$

(3) Since row 1 is not in B , we can replace every $R \in \mathcal{B}$ by $R \cup \{1\}$. Form the $k \times k$ matrix \bar{B} with these “upshifted” rows, then let

$$v_2 := u_1 C_{\bar{B}, \mathcal{A}}$$

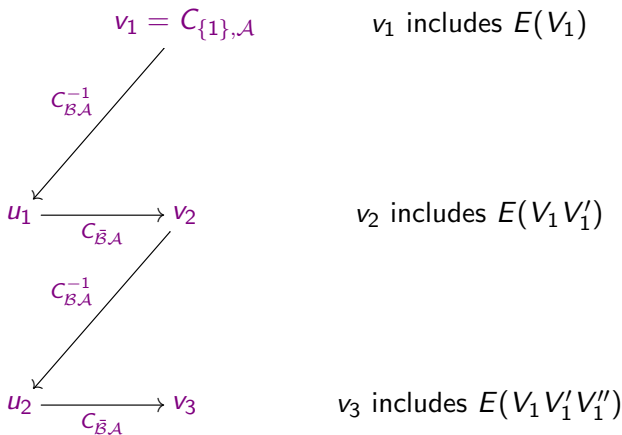
This gets us a second moment! v_2 has an entry

$$\begin{aligned}(v_2)_\emptyset &= (u_1 C_{\bar{B}, \mathcal{A}})_\emptyset \\ &= u_1 \bar{B} \operatorname{diag}(\pi) \mathbb{1} \\ &= (\mathbf{m}_1 \odot (u_1 \mathcal{B})) \operatorname{diag}(\pi) \mathbb{1} && \text{Hadamard prod. distributes} \\ &= (\mathbf{m}_1 \odot \mathbf{m}_1) \operatorname{diag}(\pi) \mathbb{1} \\ &= E(V_1 V_1')\end{aligned}$$

(4) Synthesize again! Weight rows of B to create u_2 s.t. $u_2 B = \mathbf{m}_1 \odot \mathbf{m}_1$.

$$u_2 := v_2 C_{\mathcal{B}, \mathcal{A}}^{-1}$$

and keep going!



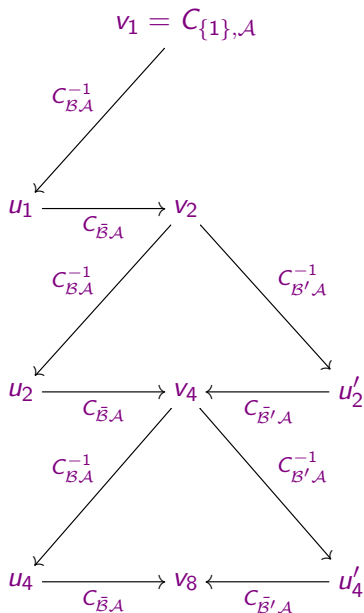
After $2k - 1$ levels, can apply k -MaxIID algorithm.

Operator norm of $C_{B,A}^{-1}$ is bounded by $(1/\zeta)^{O(k)}$ so after these $2k - 1$ levels, errors blow up by $\leq (1/\zeta)^{k^2}$. Improve this by:

Synthetic bits method with repeated squaring

Needs $n = 3k - 3$ instead of $n = 2k - 1$. Use disjoint sets A, B, B' each with $k - 1$ ζ -separated rows.

After these $\lg k$ levels, errors blow up by $\leq (1/\zeta)^{k \lg k}$.
 \Rightarrow sample size matches (almost) the $(1/\zeta)^k$ lower bound.

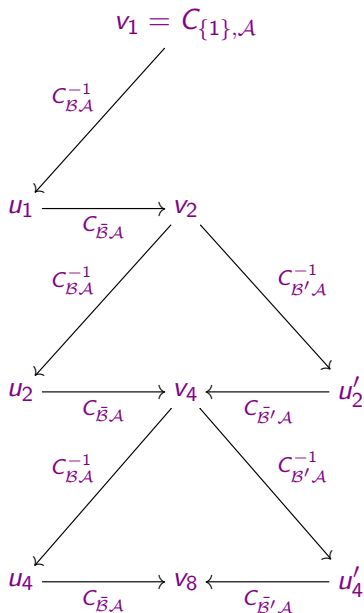


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Needs $n = 3k - 3$ instead of $n = 2k - 1$. Use disjoint sets A, B, B' each with $k - 1$ ζ -separated rows.

After these $\lg k$ levels, errors blow up by $\leq (1/\zeta)^{k \lg k}$.

\Rightarrow sample size matches (almost) the $(1/\zeta)^k$ lower bound. Proves Thm 2 (k -MixProd analysis).



Onwards:

1. “Learn” k -MixProd in Weierstrass (transportation) distance in time similar to identification? (Do have such results for k -MixIID.) I.e., $\sim \exp(k \lg k)$ rather than $\exp(k^3)$?
2. Parametric models.
3. Multiple cardinality- or dimension-bounded confounders.

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