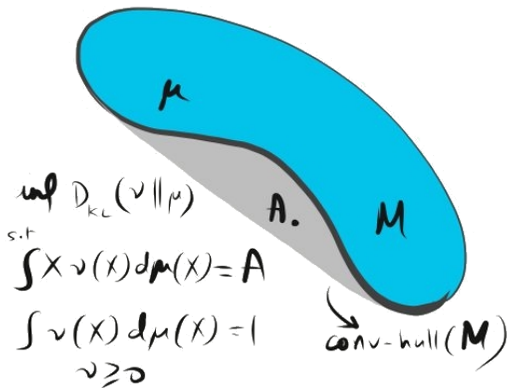


# Continuous Maximum Entropy Distributions



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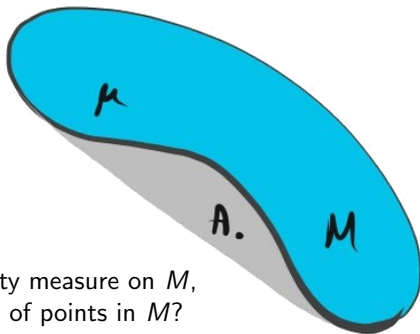
joint with

Nisheeth Vishnoi (Yale University)

## A First Question

Suppose we have:

- $M$ , a subset of  $\mathbb{R}^d$ ,
- $K$ , the convex hull of  $M$ , and
- $A$ , a point in  $K$ .



### Question:

Is there a canonical way to choose a probability measure on  $M$ , which expresses  $A$  as a “convex combination” of points in  $M$ ?

### When $M$ is discrete:

Choose the distribution on  $M$  which **maximizes Shannon entropy** [Jaynes '57].

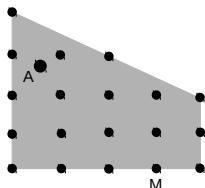
- Called **maximum entropy distributions**.
- Generalizable to minimization of **Kullback-Leibler divergence**.
- Applications to **combinatorial inequalities and discrete approximation**, algorithms for **metric TSP** approximation, **group orbit optimization** and **scaling algorithms**, and **statistical models**.

### How do we actually compute this?

# Computing Discrete Max-Entropy Distributions

Suppose we have:

- $M$ , a subset of  $\mathbb{R}^d$ ,
- $K$ , the convex hull of  $M$ , and
- $A$ , a point in  $K$ .



Primal	Dual
$\sup_{\substack{\text{supp}(\nu) \subseteq M \\ \mathbb{E}[\nu]=A}} - \sum_{x \in M} \nu(x) \log \nu(x)$	$\inf_{y \in \mathbb{R}^d} \left( \langle y, A \rangle + \log \sum_{x \in M} e^{-\langle y, x \rangle} \right)$

Dual program efficiently computable via **ellipsoid method**, given a **strong counting oracle** for the objective function [Singh-Vishnoi '15, Straszak-Vishnoi '17].

## Observations:

- Dual search space dim.  $\ll$  primal search space dim.
- Maximum entropy distribution  $\implies \nu^* \propto e^{-\langle y^*, x \rangle}$ .
- Dual program is **polynomial capacity**  $\inf_{x>0} \frac{p(x)}{x^A}$  (see [Straszak-Vishnoi '17]).

# An Aside: Polynomial Capacity

**Last slide:** Discrete maximum entropy  $\iff$  polynomial capacity.

**Polynomial capacity:**  $\inf_{x>0} \frac{p(x)}{x_1^{\alpha_1} \dots x_d^{\alpha_d}}$  for  $p \in \mathbb{R}[x]$ ,  $\alpha \in \mathbb{R}^d$  due to [Gurvits '04].

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## Polynomial capacity / discrete max-entropy applications:

- Approximate/bound **permanent** of non-negative matrices, **mixed discriminant** of PSD matrices, **mixed volume** of convex bodies [Gurvits '00s].
- Approximate/bound **integer points/volume** of polytopes [Barvinok '00s, Barvinok-Hartigan '09, Gurvits '15, Gurvits-L '20, Csikvári-Schweitzer '20, Brändén-L-Pak '21].
- **Matroid** counting/optimization [Straszak-Vishnoi '17, Anari-Oveis Gharan '17, Anari-Liu-Oveis Gharan-Vinzant '19].
- **Metric TSP** approximation [Karlin-Klein-Oveis Gharan '21]

## Capacity, beyond polynomials:

- Matrix, operator, tensor **scaling** [Garg-Gurvits-Oliveira-Wigderson '15, Franks '18, Bürgisser-Franks-Garg-Oliveira-Walter-Wigderson '18, Allen-Zhu-Garg-Li-Oliveira-Wigderson '18, van Apeldoorn-Gribling-Li-Nieuwboer-Walter-de Wolf '20, Bürgisser-Li-Nieuwboer-Walter '20].
- **Non-commutative optimization** and **orbit intersection** [Bürgisser-Garg-Oliveira-Walter-Wigderson '17, Allen-Zhu-Garg-Li-Oliveira-Wigderson '18, Bürgisser-Franks-Garg-Oliveira-Walter-Wigderson '19, Franks-Walter '20, Franks-Reichenbach '21].
- **Stat models** [Améndola-Kohn-Reichenbach-Seigal '21, Franks-Oliveira-Ramachandran-Walter '21]

# Back to Maximum Entropy Distributions

**Discrete case:** Efficiently computable, with many applications / generalizations.

## What about non-discrete support?

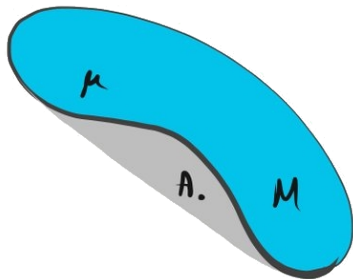
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Suppose we have:

- $M$ , a **manifold** embedded in  $\mathbb{R}^d$ ,
- $K$ , the convex hull of  $M$ , and
- $A$ , a point in  $K$ .

## Maximum entropy program:

$$\sup_{\substack{\text{supp}(\nu) \subseteq M \\ \mathbb{E}[\nu] = A}} - \int_M \nu(x) \log \nu(x) dx$$



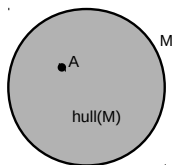
## New problems:

- What is  $dx$  here? **No canonical notion of entropy.**
- Search space now **infinite-dimensional.**
- $M$  can only be **implicitly expressed** (e.g., “all unit vectors”).

# Computing Continuous Max-Entropy Distributions

Suppose we have:

- $M$ , a **manifold** embedded in  $\mathbb{R}^d$ ,
- $\mu$  a **probability measure** on  $M$ ,
- $K$ , the convex hull of  $M$ , and
- $A$ , a point in  $K$ .



Primal	Dual
$\sup_{\substack{\text{supp}(\nu) \subseteq M \\ \mathbb{E}[\nu] = A}} - \int_M \nu(x) \log \nu(x) d\mu(x)$	$\inf_{y \in \mathbb{R}^d} \left( \langle y, A \rangle + \log \int_M e^{-\langle y, x \rangle} d\mu(x) \right)$

Dual search space is **finite-dimensional** and by strong duality [L-Vishnoi '20], the maximum entropy distribution is of the form  $d\nu^*(x) \propto e^{-\langle y^*, x \rangle} d\mu(x)$ .

## Questions remain:

- Why do we care about this?
- Efficient computability? (Now relies on an **integration oracle**.)
- Efficient sampling?

## Motivating Example: Quantum Entropy

**Von Neumann entropy:** Derived from **discrete** distribution on eigenvector pure states:

$$\mathcal{H}_{\text{vN}}(A) = -\text{tr}(A \log A) = -\sum_{i=1}^m \lambda_i \log \lambda_i,$$

where  $\lambda_i$  are the eigenvalues of the **density matrix**  $A$  (PSD,  $\text{tr}(A) = 1$ ).

### Issues:

- Critiqued by [Band-Park '76] as an indicator of uncertainty of density matrix  $A$ .
- Can be viewed as the entropy of a **minimum** entropy distribution (“most terse” description of  $A$ ).

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**Quantum entropy:** Derived from **continuous** distribution on **all** pure states:

$$\mathcal{H}_q(A) = \sup_{\mathbb{E}[\nu]=A} - \int_{\mathcal{P}_1} \nu(X) \log \nu(X) d\mu_1(X),$$

where  $\mathcal{P}_1 =$  all Hermitian **rank-one PSD projections**,  $\text{hull}(\mathcal{P}_1) =$  all **density matrices**, and  $\mu_1 =$  the **Haar measure** on  $\mathcal{P}_1$  (via action of unitary group).

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**Efficient computability:** Low-dim. cases [Slater '91], general case\* [L-Vishnoi '20]

# Motivating Example: Interior Point Methods

**Entropic barrier function** for convex body  $K \subset \mathbb{R}^d$ :

$$B_K(v) = \sup_{y \in \mathbb{R}^d} \left[ \langle y, v \rangle - \log \int_K e^{\langle y, x \rangle} dx \right].$$

- Barrier for  $K$  with **optimal self-concordance parameter** [Bubeck-Eldan '15].
- $-B_K(v)$  is the optimal value of **maximum entropy dual program**.
- Open manifold  $K \subset \mathbb{R}^d$  with  $\mu =$  **Lebesgue measure** restricted to  $K$ .

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**Open problem:** Is the entropic barrier efficiently computable for  $K$ , given as a membership oracle?

**Partial answer:** Yes, given an efficient **integration oracle** for  $K$ .

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**Other connections** beyond quantum entropy and interior point methods:

- Isotropic constant [Klartag '06], see also [Gromov '90]
- Langevin/Bingham distributions on matrix spaces [Khatri-Mardia '77, Chikuse '03]
- SDP rounding, from matrix to vector [Goemans-Williamson '95]

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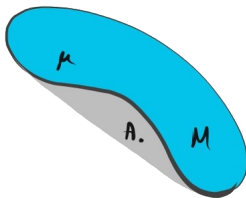
**Question:** What about efficient **computability** (and **sampling**)?



# Computability of Continuous Max-Entropy Distributions

Suppose we have:

- $M$ , a **manifold** embedded in  $\mathbb{R}^d$ ,
- $\mu$  a **probability measure** on  $M$ ,
- $K$ , the convex hull of  $M$ , and
- $A$ , a point in  $K$ .



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**Ellipsoid method-based algorithm** [L-Vishnoi '20] to efficiently approximate the corresponding maximum entropy distribution whenever:

- $\mu$  is “balanced” - intuitively,  $\mu$  is close to **uniform/symmetric**,
- $A$  is not too near the boundary of  $K$ , and
- we have a **strong integration oracle** for  $\mu$ ; i.e. oracles for

$$F(Y) = \log \int_M e^{-\langle Y, X \rangle} d\mu(X) \quad \text{and} \quad \nabla F(Y).$$

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**Question:** When can we obtain these **strong integration oracles**?

# Integration Oracles via Symmetry

**Focus:**  $M = \text{orbit of the unitary group}$  in the space of Hermitian matrices,

$$\mathcal{O}_\lambda := \{U \text{diag}(\lambda)U^* : U \in U(n)\} = \{\text{Hermitian matrices with e-vals } \lambda\}.$$

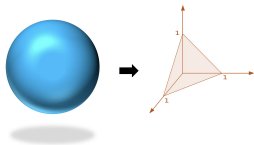
- **Invariant measure**  $\mu_\lambda$  on  $\mathcal{O}_\lambda$  derived from **Haar measure** on  $U(n)$ .
- **Question:** What about the oracle for  $\log \int_{\mathcal{O}_\lambda} e^{-\langle Y, X \rangle} d\mu_\lambda(X)$ ?

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**Quantum entropy case:**  $\lambda = e_1 = (1, 0, 0, \dots, 0) \implies \mathcal{O}_\lambda = \text{rank-1 PSDs}$ .

**Polytope connection:** “Moment map” on  $\mathcal{O}_{e_1}$ :

$$\pi : \mathcal{O}_{e_1} \rightarrow \mathbb{R}^n \quad \text{via} \quad \pi : X \mapsto \text{diag. entries of } X$$



- Image of  $\pi$  acting on  $\mathcal{O}_{e_1}$  is **the standard simplex**  $\Delta_n = \text{hull}(e_1, \dots, e_n)$
- That is,  $\pi(\mathcal{O}_{e_1}) = \Delta_n$ , but **even better**  $\pi(d\mu_{e_1}) = dx|_{\Delta_n}$ :

$$\int_{\mathcal{O}_{e_1}} e^{-\langle Y, X \rangle} d\mu_{e_1}(X) = \int_{\Delta_n} e^{-\langle y, x \rangle} dx \quad \text{for } Y = \text{diag}(y).$$

**Bonus:** Bijection between max-entropy distributions on  $\mathcal{O}_{e_1}$  and  $\Delta_n$ .

# Integration Oracles via Symmetry

**Focus:**  $M =$  orbit of the unitary group in the space of Hermitian matrices,

$$\mathcal{O}_\lambda := \{U \operatorname{diag}(\lambda) U^* : U \in U(n)\} = \{\text{Hermitian matrices with e-vals } \lambda\}.$$

**Last slide:**  $\mathcal{O}_{e_1} \iff \Delta_n$  connection: measures, max-entropy, integration.

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**What about other  $\lambda$ ?** Want to compute  $\log \int_{\mathcal{O}_\lambda} e^{-\langle Y, X \rangle} d\mu_\lambda(X)$ .

- Apply “moment map”  $\pi$  to  $\mathcal{O}_\lambda$  to get a **polytope**:

$$\pi(\mathcal{O}_\lambda) = \operatorname{hull}\{\sigma \cdot \lambda : \sigma \in S_n\}.$$

- **Problem:**  $\pi(d\mu_\lambda) =$  piecewise polynomial density [Duistermaat-Heckman '82].
- **Fix** is possible via different map/polytope (see **Colin's talk** next).
- **Solution:** The **Harish-Chandra-Itzykson-Zuber formula**:

$$\int_{U(n)} e^{\operatorname{tr}(AUBU^*)} dU = \left( \prod_{p=1}^{n-1} p! \right) \frac{\det([e^{\alpha_i \beta_j}]_{i,j=1}^n)}{\prod_{i < j} (\beta_j - \beta_i)(\alpha_j - \alpha_i)}.$$

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**Question:** Why do we care about these orbits? Other manifolds?

# Integration Oracles via Symmetry

**Previous slides:** Integration oracles for Hermitian matrix unitary orbits:

$$\mathcal{O}_\lambda := \{U \operatorname{diag}(\lambda) U^* : U \in U(n)\} = \{\text{Hermitian matrices with e-vals } \lambda\}.$$

Implies efficient algorithm for **computing max-entropy distributions** on  $\mathcal{O}_\lambda$ .

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**Question:** Where do these distributions appear?

- The  $\mathcal{O}_{e_1}$  case is equivalent to computing **quantum entropy**.
  - The  $\mathcal{O}_{e_1+e_2+\dots+e_k}$  case yields max-entropy distributions on the **complex Grassmanian**: utilized in **private low-rank approx.** [L-McSwiggen-Vishnoi '21]
  - Complex Langevin/Bingham matrix distributions [Khatri-Mardia '77, Chikuse '03]
  - **HCIZ densities** in physics, random matrix theory, ...
- 

**Question:** What about other manifolds?

- Need a **measure** (via symmetry) and an **integration oracle**.
- **E.g.:** HCIZ formula extends to other adjoint orbits of **compact Lie groups**, with measure via the **Haar measure**. (see [L-Vishnoi '20])

# Conclusion

## This talk:

- **Discrete** maximum entropy distributions and applications: combinatorics, discrete approximation, orbit problems, norm minimization, etc.
  - **Continuous** maximum entropy distributions and (**different**) applications: quantum entropy, interior point methods, low-rank approximation, etc.
  - Specific examples related to **conjugation orbits of unitary groups**.
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**Upshot:** In the continuous case, **manifold symmetries** give rise to measures and integration oracles which **make efficient computability possible**.

**Question:** But what about **sampling**?

- **Answer:** See next talk by Colin.
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## Open problems:

- Many **combinatorial** applications of **discrete** maximum entropy distributions. **Anything like this for the continuous case?**
  - **Entropy interpretation** of the norm minimization generalization?
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Thanks!