# Charting the Landscape of Memory/Data Tradeoffs in Continuous Optimization: A Survey of Open Problems



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# How do *memory constraints* influence the speed of learning/optimization?

Today: Linear Regression min x'Ax - bx(or solving Ax = b)

Access to i.i.d. data samples

Distribution over (a,b) pairs

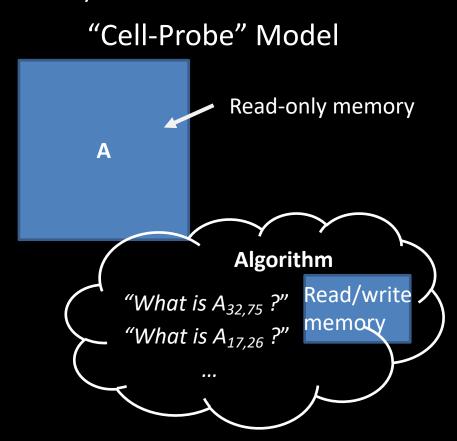
(e.g. a  $\leftarrow$ -- Gaussian, b = <a, x>

Algorithm

"Give me a datapoint"

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"...



|   | i.i.d. data samples  | "Cell-Probe" Model  |
|---|--|---|
| Discrete setting  Solve Ax=b  over a finite field                 | Unknown $x \in \{0,1\}^d$ chosen uniformly at random.  Each datapoint: $(a_i, < a_i, x > \text{mod } 2),$ $a_i \in \{0,1\}^d$ chosen unif. rand  Goal: find $x$  | Same as i.i.d. data setting,<br>but datapoints stored in<br>read-only memory. |
| Continuous setting Solve Ax=b over R <sup>d</sup> (or regression) | Unknown $x \in \mathbb{R}^d$ with $ x =1$ chosen uniformly at random.  Each datapoint: $(a_i, < a_i, x >),$ $a_i \in \mathbb{R}^d$ chosen e.g. from N(0,I <sub>d</sub> ) [or some other distribution]  Goal: approximate $x$ | Same as i.i.d. data setting, but datapoints stored in read-only memory.       |

Many other interesting models...

No direct access to datapoints, but instead interact via specific types of oracles

- Statistical Query access
- Function evaluation queries
- Gradient queries
- Etc.

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| Discrete setting    |
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| over a finite field |
|                     |

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# Continuous setting Solve Ax=b over R<sup>d</sup> (or regression)

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Goal: approximate x

Unknown  $x \in \{0,1\}^d$  chosen uniformly at random. Given access to stream of examples  $(a_i, < a_i, x > \text{mod } 2)$ ,  $a_i \in \{0,1\}^d$  chosen uniformly at random.

Gaussian elimination:  $O(d^2)$  memory , O(d) examples Brute force guess/check: O(d) memory , O(d) examples.

Conjecture [Steinhardt, Valiant, Wager'15]: Any algorithm with  $< d^2/4$  memory needs exponential number of samples to learn x. Unknown  $x \in \{0,1\}^d$  chosen uniformly at random. Given access to stream of examples  $(a_i, < a_i, x > \text{mod } 2)$ ,  $a_i \in \{0,1\}^d$  chosen uniformly at random.

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Thm [Raz'16]: Conjecture [Steinhardt, Valiant, Wager'15]: Any algorithm with  $< d^2/4$  memory needs exponential number of samples to learn x.

Subsequent work extended this to a broad class of learning problems over finite fields

Kol-Raz-Tal'17: sparse parities

Raz'17, Moshkovitz-Moshkovitz'17,18, Beame-Ovies Gharan-Yang'18, Garg-Raz-Tal'18: Large class of learning problems over finite fields satisfying combinatorial/mixing properties.

...and more recent papers

See Sumegha Garg's talk Thursday!!

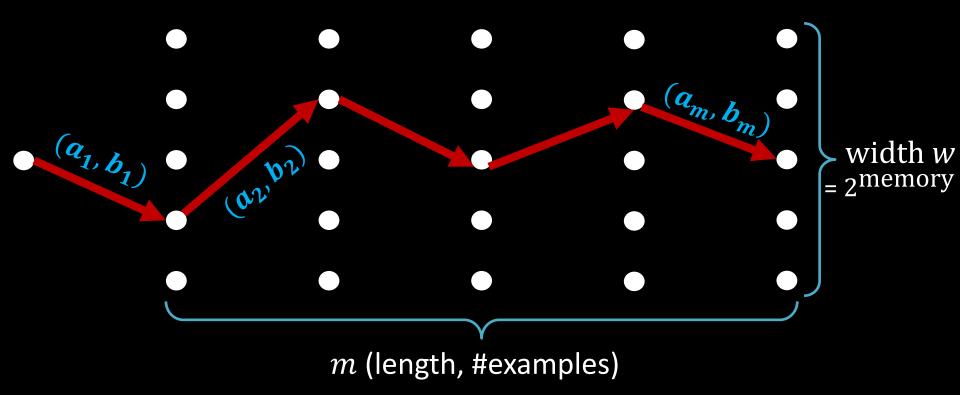


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# Branching program for learning



Each layer corresponds to a time step

Each vertex corresponds to a memory state

Every memory state has a transition function (an 'edge') which is a mapping from an example (a, b) to a vertex in the next layer.

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| Continuous setting Solve Ax=b over R <sup>d</sup> (or regression) | Unknown $x \in \mathbb{R}^d$ with $ x =1$ chosen uniformly at random.  Each datapoint: $(a_i, < a_i, x > +noise)$ $a_i \in \mathbb{R}^d$ chosen e.g. from $N(0,I_d)$ [or some other distribution]  Goal: approximate $x$        | Same as i.i.d. data setting,<br>but datapoints stored in<br>read-only memory. |

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Goal: approximate x

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#### "Cell-Probe" Model

Same as i.i.d. data setting, but datapoints stored in read-only memory.

Easy: linear space and poly(d) queries.

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[or some other distribution]

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Linear regression: core ML and convex optimization problem

#### **Broader Context:**

1<sup>st</sup> order methods (linear memory, more samples) vs 2<sup>nd</sup> order methods (quadratic memory, less samples)

Huge effort to find optimization algorithms with linear memory, that behave like quadratic-memory algorithms (e.g. conjugate gradient...)

Largely unexplored frontier of continuous optimization research:

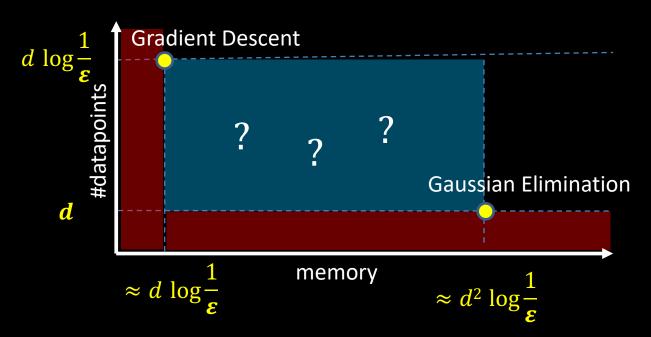
- Vast literature on lower bounds on #queries required from oracle
- Relatively recent work with *linear* memory: [Dagan-Kur-Shamir'19] on memory lower bounds in streaming model, sparse linear regression [Steinhardt-Duchi'15]

Little on memory/sample tradeoffs for optimization with super-linear memory.

# Memory/Data Tradeoffs for Linear Regression

Unknown  $x \in \mathbb{R}^d$  with ||x|| = 1, chosen uniformly at random. Given access to stream of examples  $(a_i, b_i)$ ,

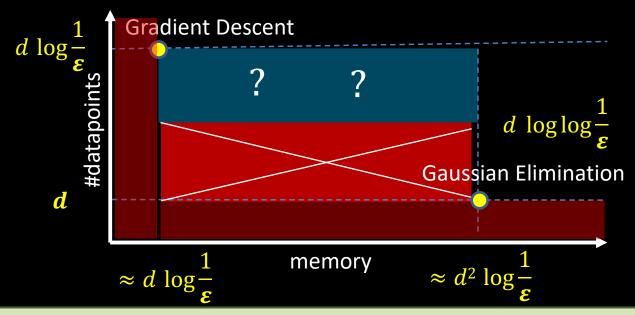
$$a_i \sim N(0, I_d)$$
  
 $b_i = \langle a_i, x \rangle + \eta_i$  noise:  $\eta_i \sim Unif[-2^{-d}, 2^{-d}]$ 



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Theorem [Sharan, Sidford, Valiant '19]:

Any algorithm with  $o(d^2)$  memory needs at least  $\Omega\left(d\log\log\frac{1}{\varepsilon}\right)$  samples to approximate x with  $L_2$  error  $\varepsilon$ .

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Either memory  $d^2$  or exp(d) datapoints.

#### Each datapoint:

 $(a_i, < a_i, x > +noise)$  $a_i \in \mathbb{R}^d$  chosen e.g. from N(0,I<sub>d</sub>) Goal: approximate x to within  $\varepsilon$ 

Thm:  $o(d^2)$  space implies need d log log  $1/\epsilon$  datapoints (with noise)

Conjecture I:  $o(d^2)$  space implies need d log  $1/\epsilon$  datapoints

Conj. II: o(d²) space implies need poly(condition number) datapoints

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Conjecture: Need super-linear space.

[Might be *very* hard to prove!!!!]

Approach of the discrete setting fails because numerical errors grow without poly(d) precision arithmetic