



# Synthesis of coordination programs from temporal specifications

**Suguman Bansal**, Rice University

Kedar Namjoshi, Nokia Bell Labs, Murray Hill

Yaniv Sa'ar, Nokia Bell Labs, Kfar Saba

# Coordination programs

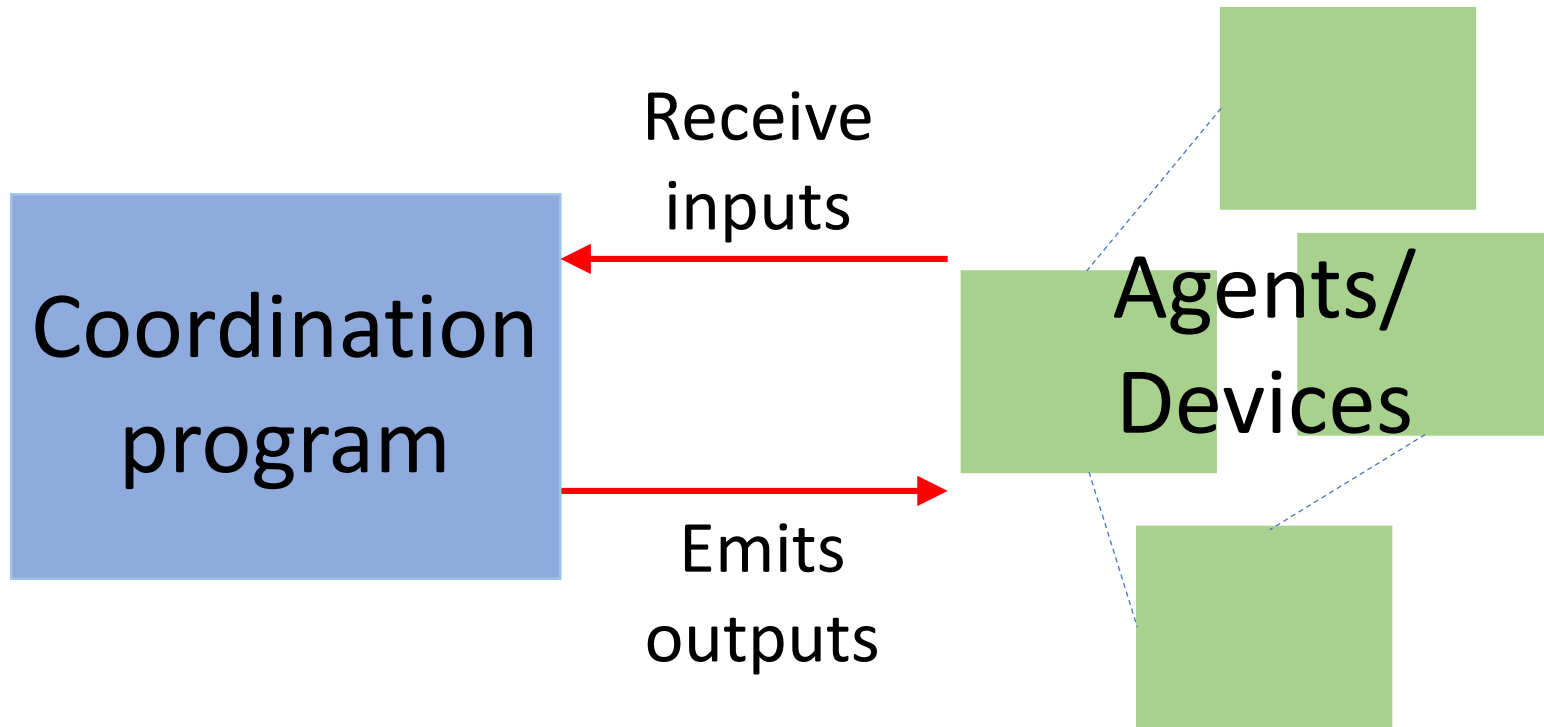


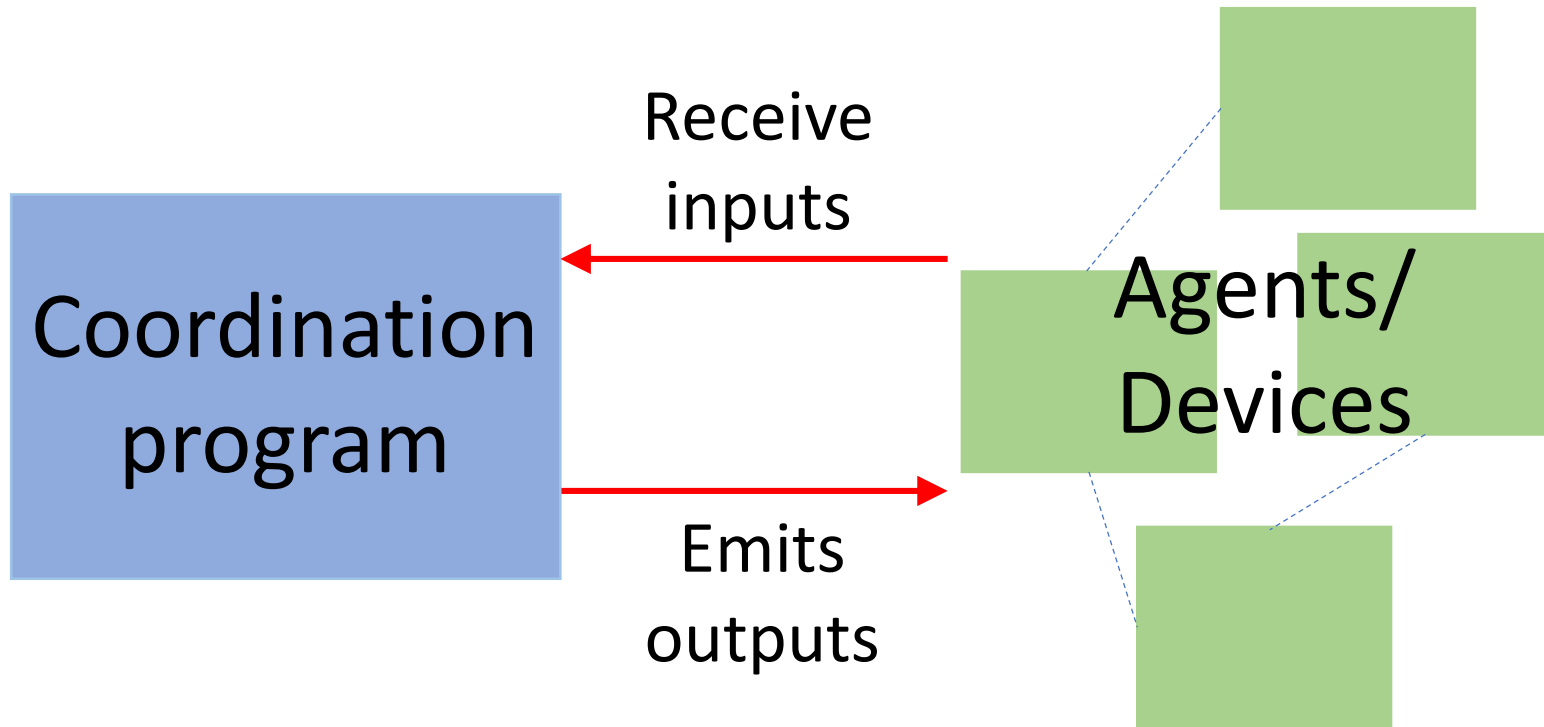




Coordination  
program

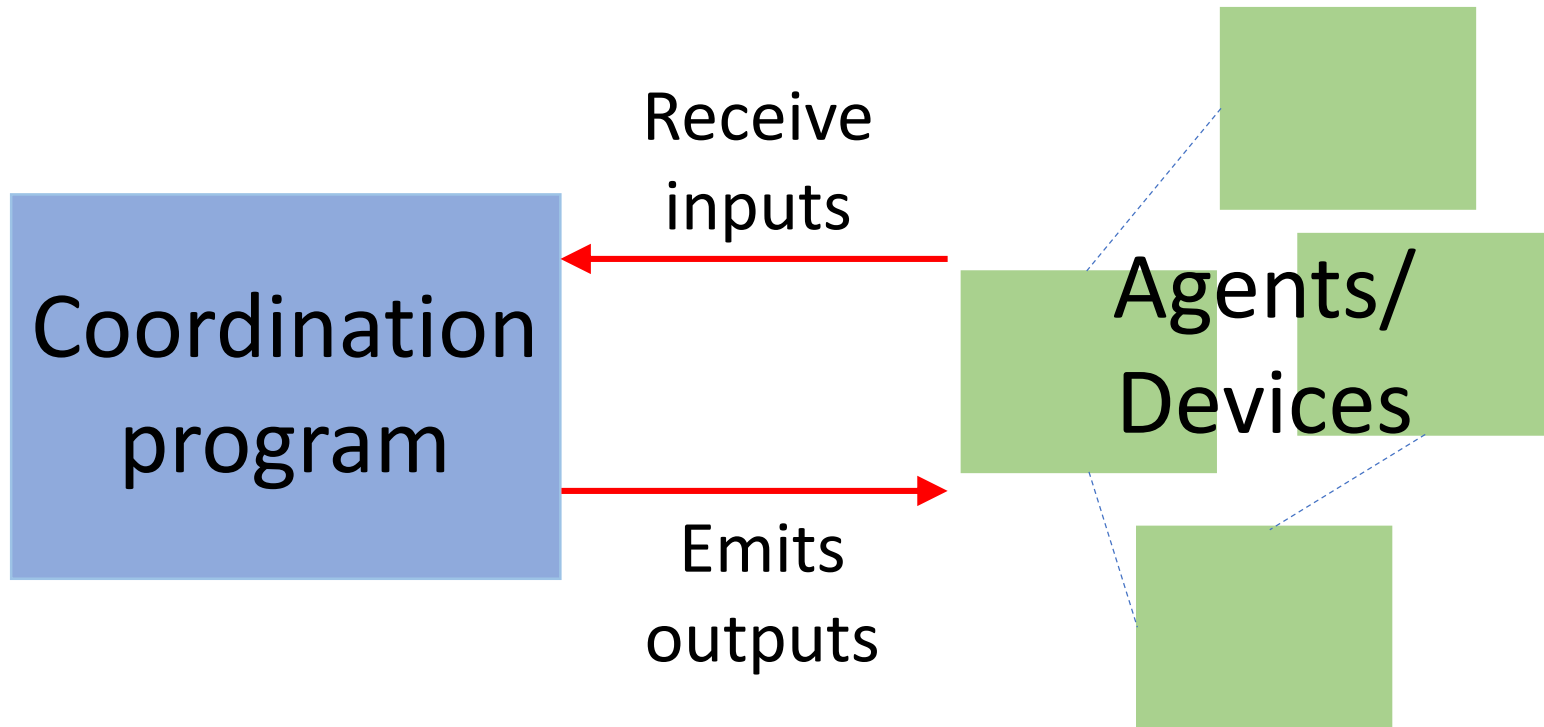






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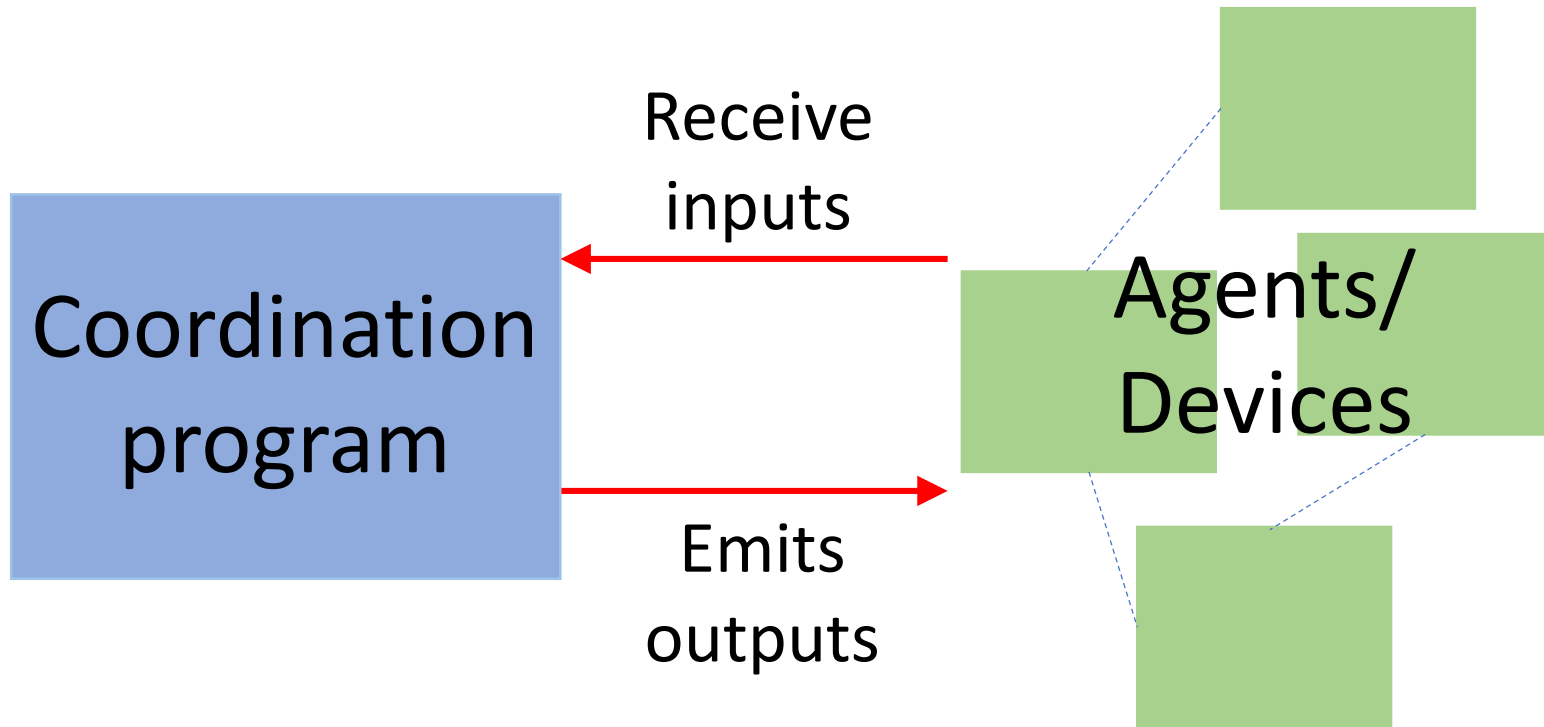




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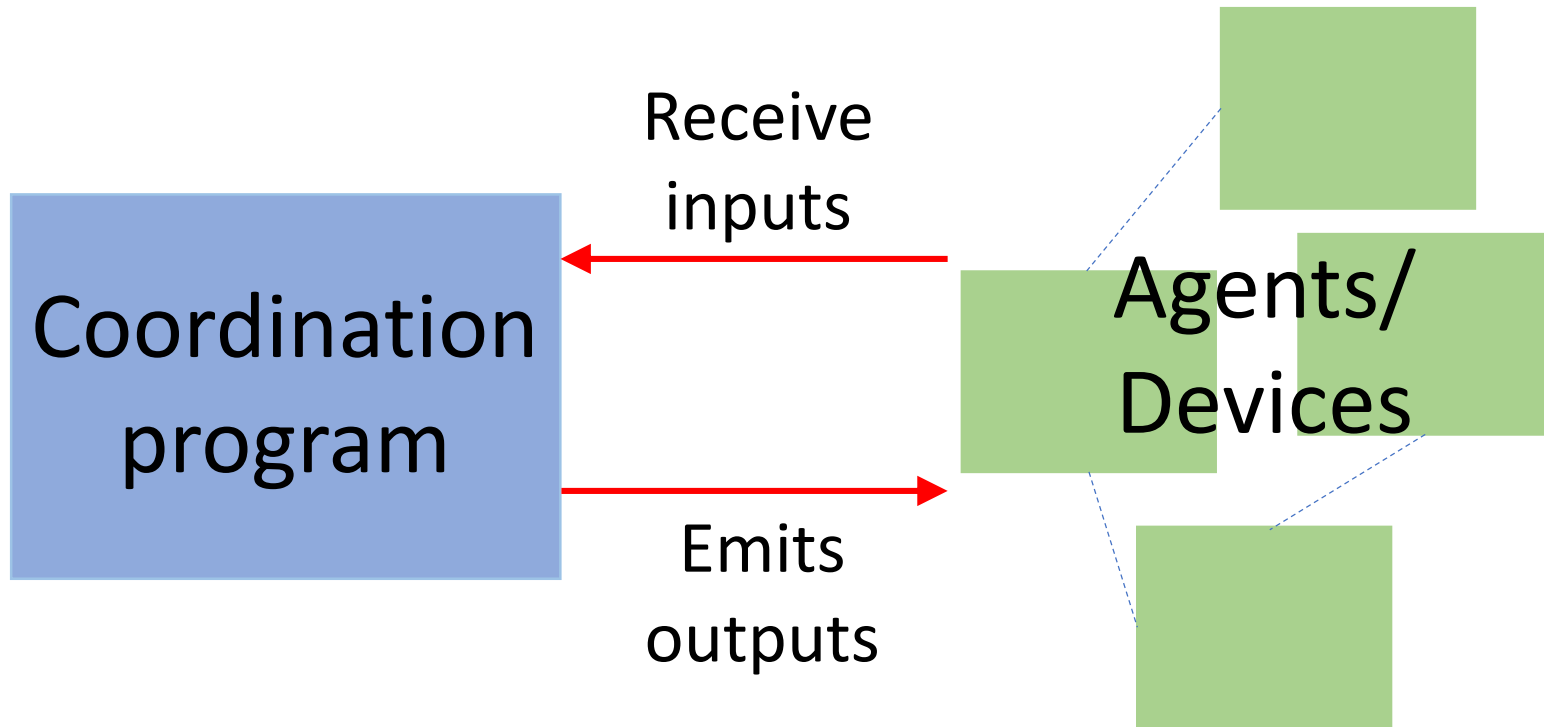
Repeated interaction with agents/devices



Designing coordination programs is hard

1. Reactive system
2. Asynchrony

Each device operates **at its own clock**



Designing coordination programs is hard

1. Reactive system
2. Asynchrony
3. Partial information

Coordination program has **limited visibility**:  
Sees the interface only

**Specifying intent of a coordination program is easier**  
E.g. “Thermostat should maintain ambient temperature”

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Coordination synthesis

# Contributions

Formulate, solve, and demonstrate coordination synthesis

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Design efficient automata-based synthesis algorithm

Accounts for all three challenges – Reactive, asynchrony, partial information

Prior work accounts for at most two

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Conduct case-studies on prototype implementation



# High-level specification

Intent of coordination program

Linear Temporal Logic (LTL) [Pnueli, FOCS 1977]

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Device description

Communicating Sequential Processes (CSP) [Hoare, CACM 1978]

- **Rich structure:** asynchrony, non-determinism ...
- **Communication model:** Message passing
- **Interface:** Visible and hidden actions

# CSP processes [Hoare, , CACM 1978]

Processes:  $P, Q, R, \dots$

- Public actions ( $a_0, a_1, \dots a_n$ ) and Private actions ( $b_0, b_1, \dots b_n$ )

$$P = action_0 \rightarrow Q_0 \mid action_1 \rightarrow Q_1 \mid \dots \mid action_n \rightarrow Q_n$$

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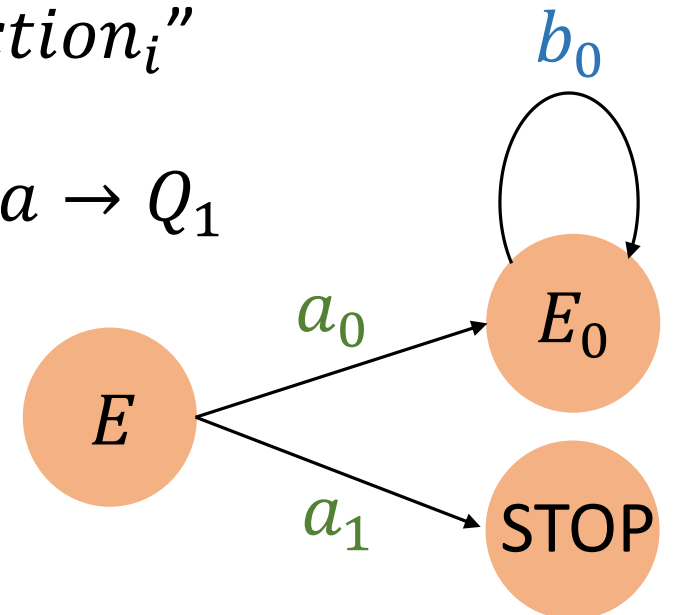
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Example:  $E = a_0 \rightarrow E_0 \mid a_1 \rightarrow STOP$

$$E_0 = b_0 \rightarrow E_0$$



# CSP interactions [Hoare, 1978]

## Synchronized public actions

- Processes evolve together on public actions

Let,  $P_0 = a \rightarrow Q_0$ ,  $P_1 = a \rightarrow Q_1$  then  $(P_0 \parallel P_1) = a \rightarrow (Q_0 \parallel Q_1)$



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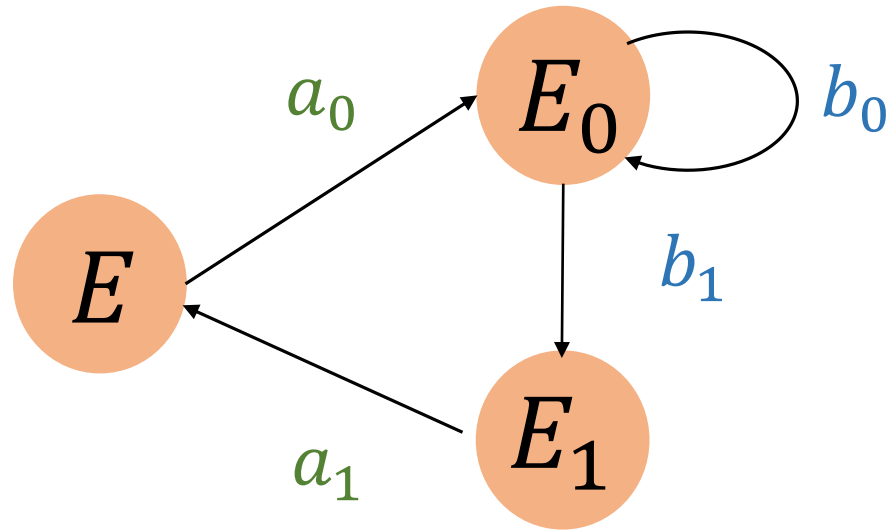
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## Internal private actions

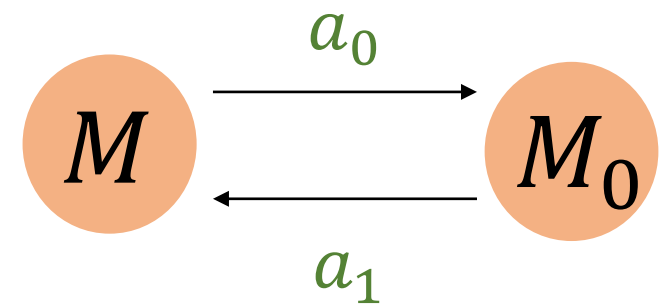
- Process with private action evolves by itself

Let,  $P_0 = b \rightarrow Q_0$ ,  $P_1 = a \rightarrow Q_1$  then  $(P_0 \parallel P_1) = b \rightarrow (Q_0 \parallel P_1)$

# Example



||



$a_0 b_0 b_0 b_0 \dots$

$a_0 b_1 a_1 \dots$

...

# Coordination synthesis

Given a CSP environment description

**$E = E1 \parallel E2 \parallel \dots \parallel E_n$**

and an LTL specification  **$S$** ,

Generate a coordinator (CSP)  **$M$**  s.t.

**$E \parallel M$**  satisfies  **$S$**

# Challenge I: Partial information

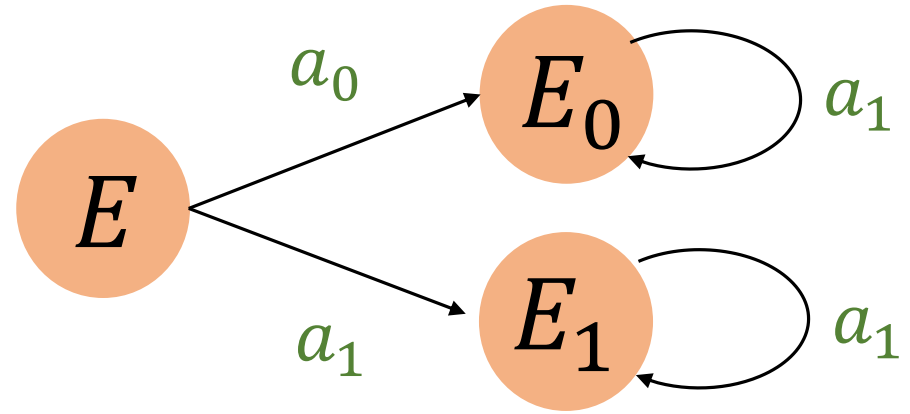
Coordinator may not know the current state of **E**

- Structural non-determinism:  
Evolution to multiple states
- Internal actions:  
Evolution due to internal actions is unknown to coordinator

# Challenge II: Deadlock freedom

Coordinator must guarantee no deadlock despite partial information

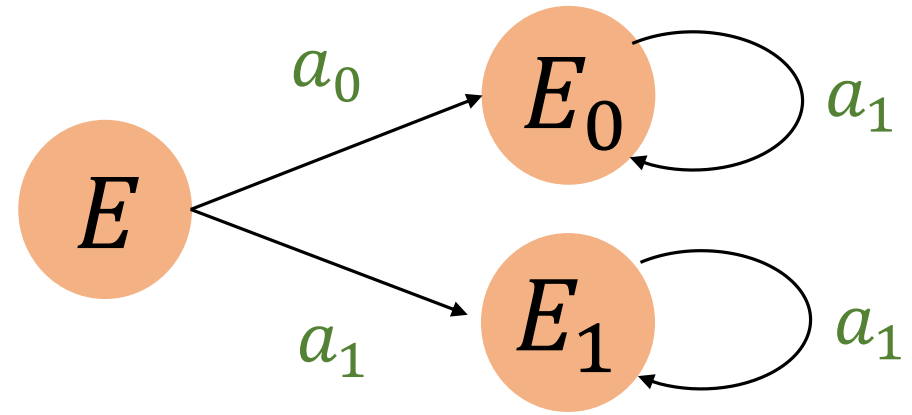
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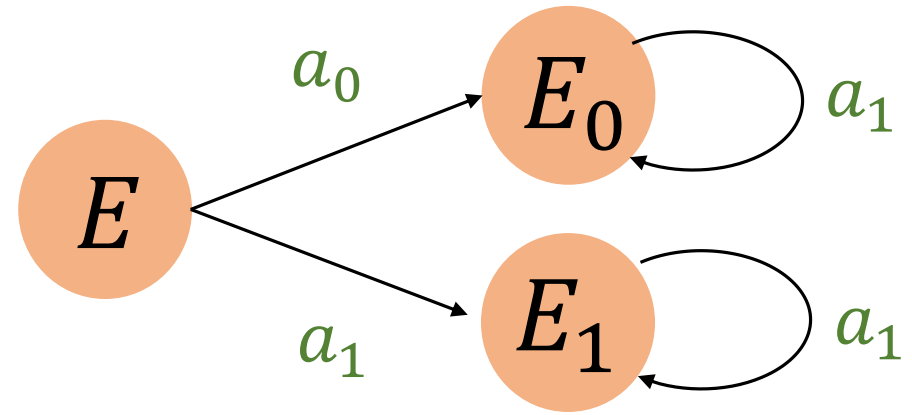


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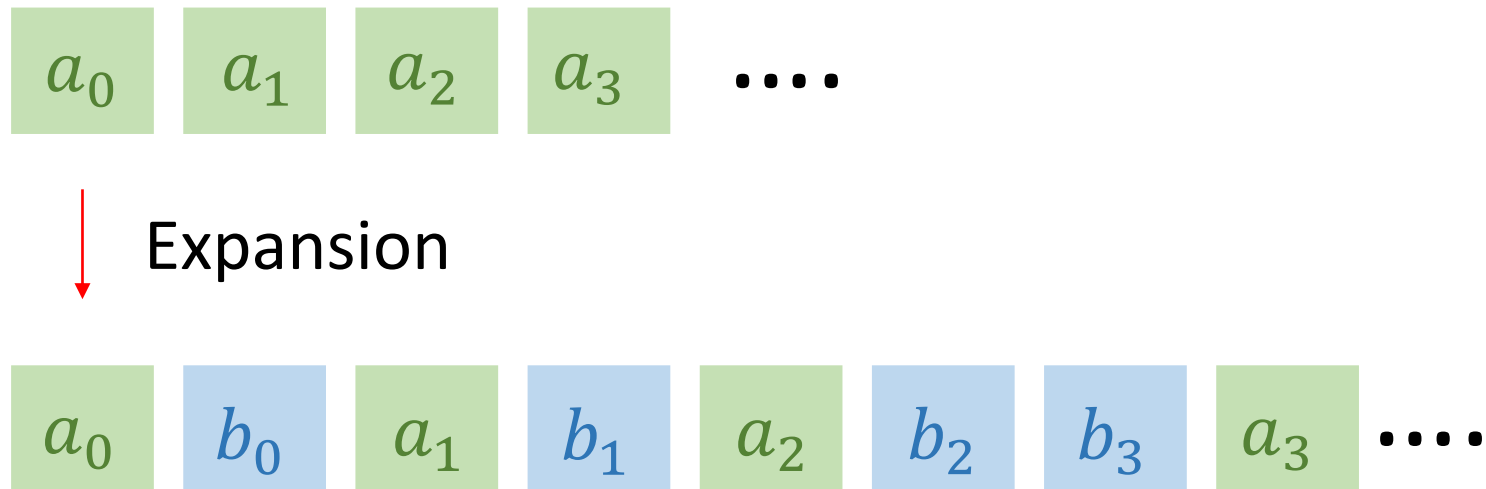
$$\begin{aligned} M &= a_0 \rightarrow M \\ E \parallel M &\text{ deadlocks} \end{aligned}$$

$$\begin{aligned} M &= a_1 \rightarrow M \\ E \parallel M &\text{ does not deadlock} \end{aligned}$$



# Challenge III: Asynchrony

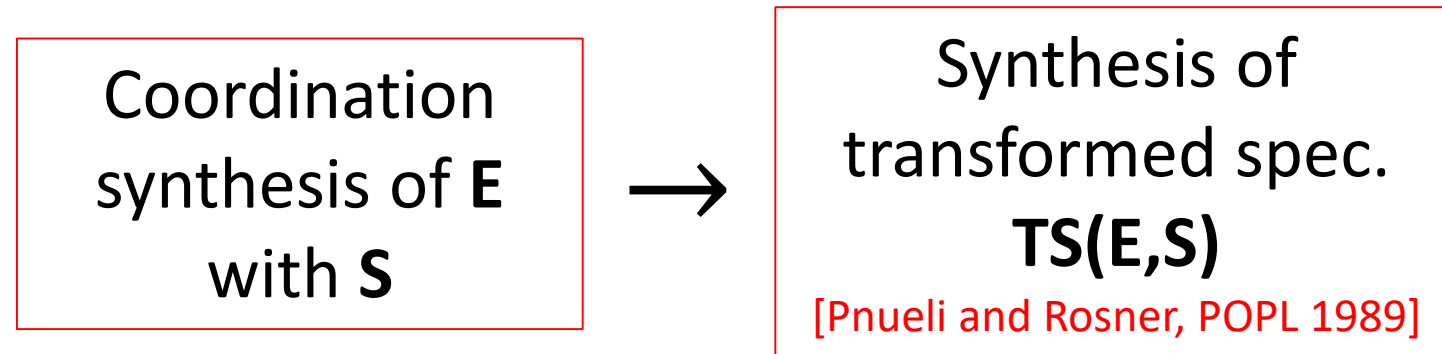
Coordinator doesn't know how many or which actions have taken place



Every valid expansion satisfies **S**

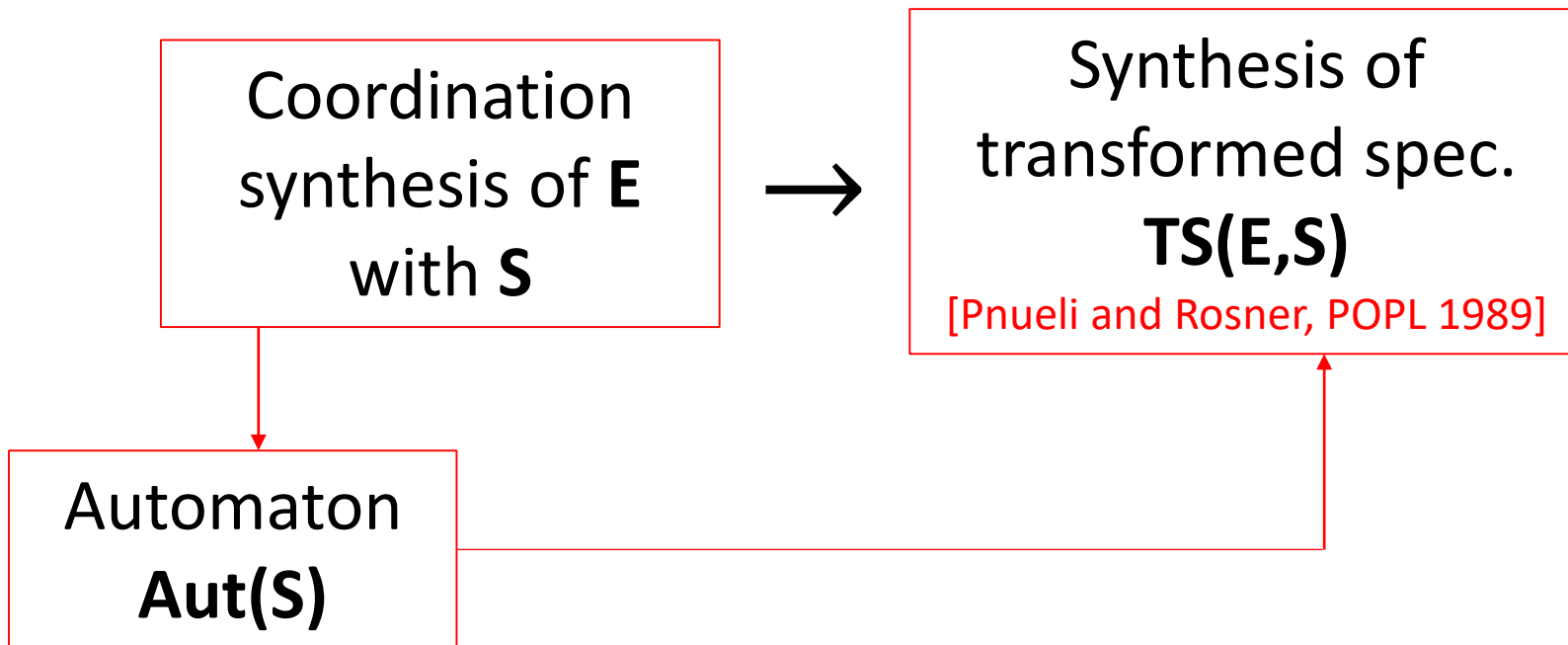
# Core technique

Coordination synthesis with  $E, S$  reduces to Synch. synthesis with  $TS(E, S)$



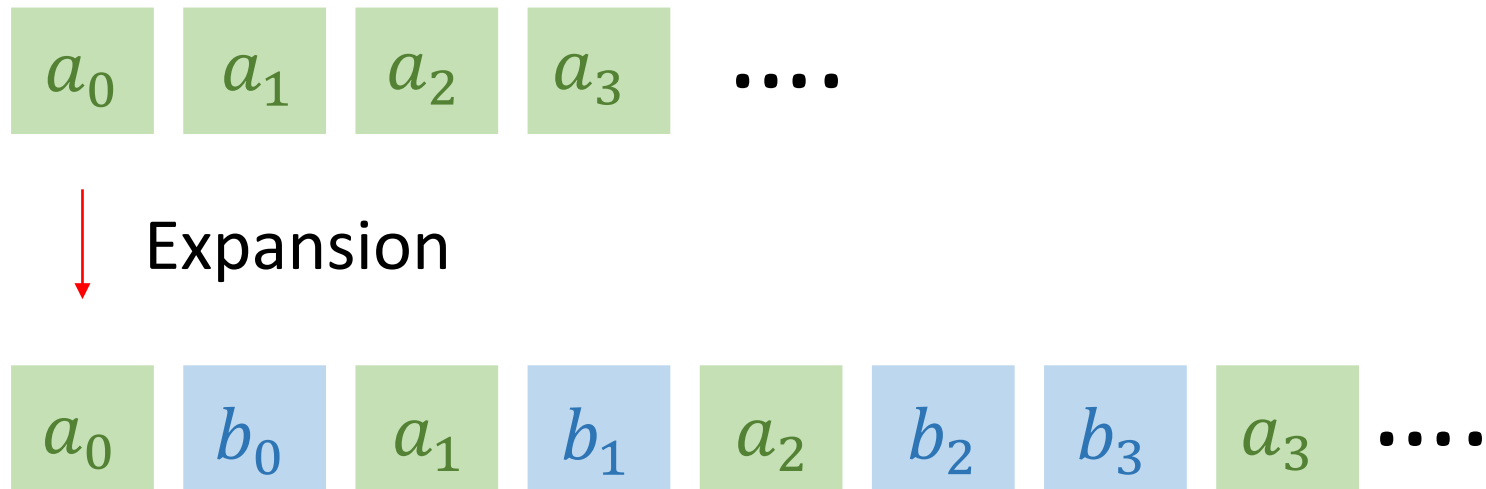
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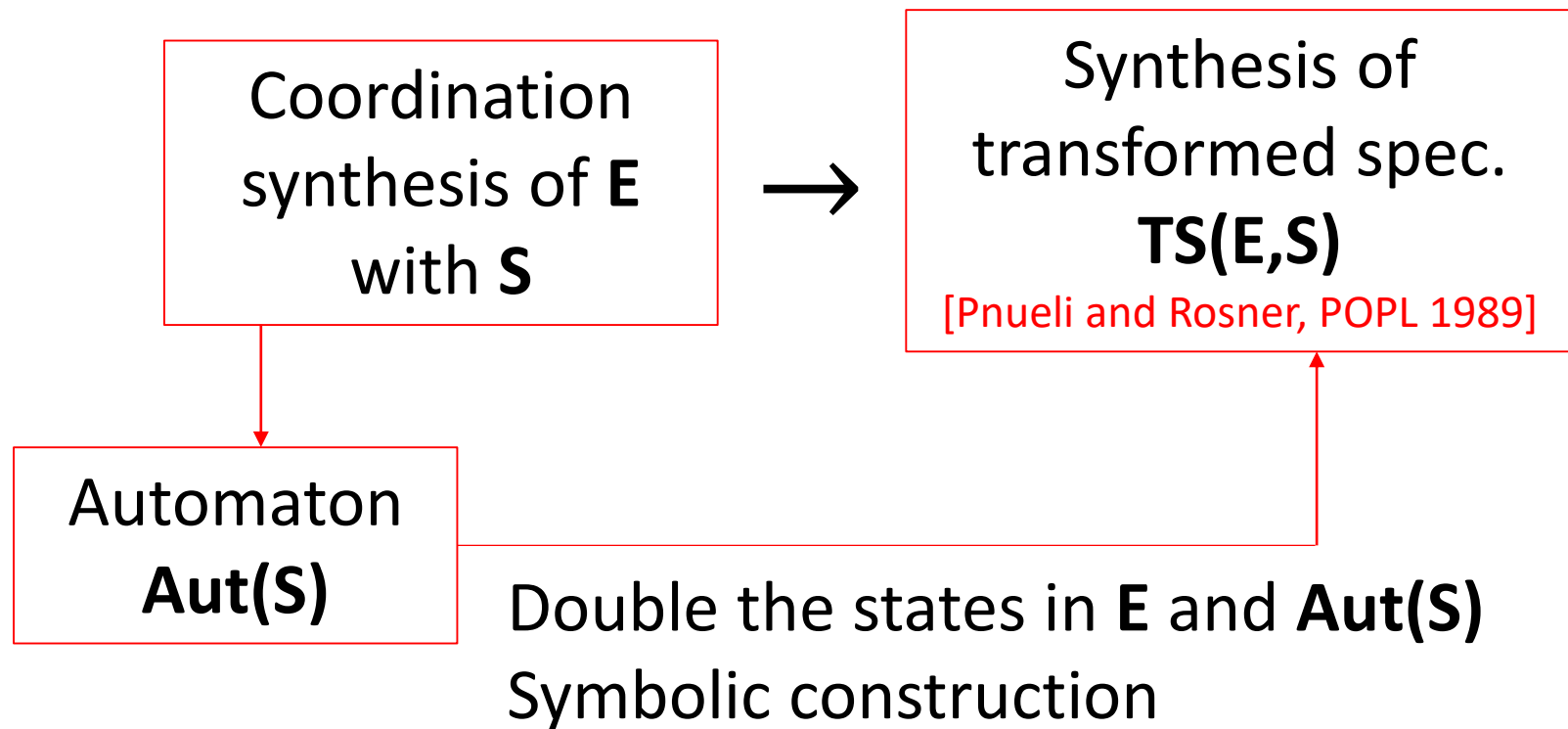
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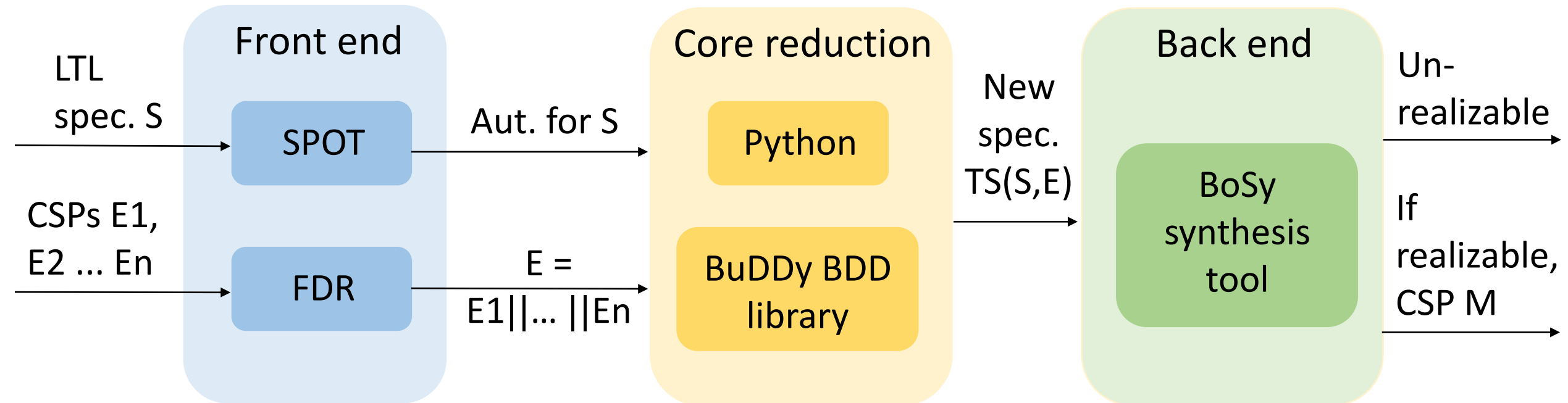


# Implementation



## Prototype **CoSy** (**Co**ordination **Sy**nthesis)

- Core reduction in Python
- BDD-based symbolic reduction



# Thermostat

How will coordination synthesis help?

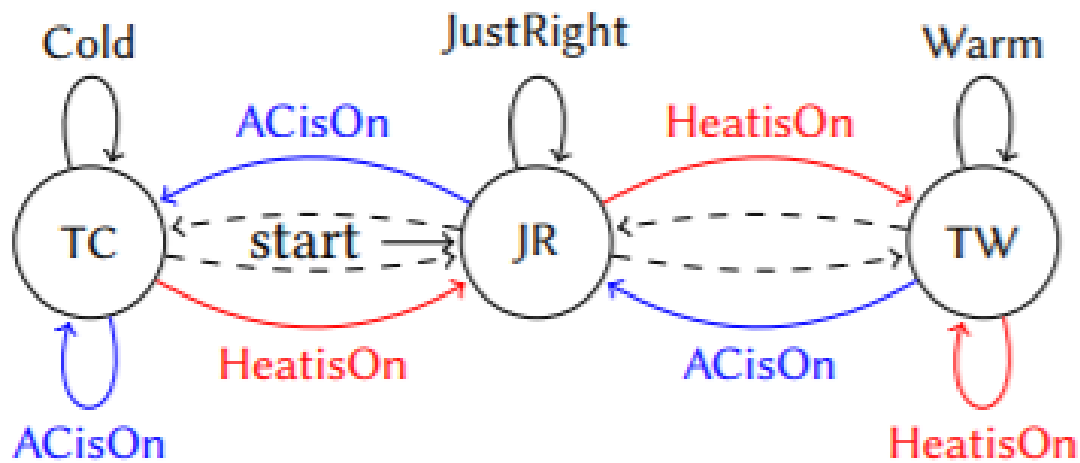


Fig. 6. Room temperature Sensor process (JR: Just right, TW: Too warm, TC: Too cold).

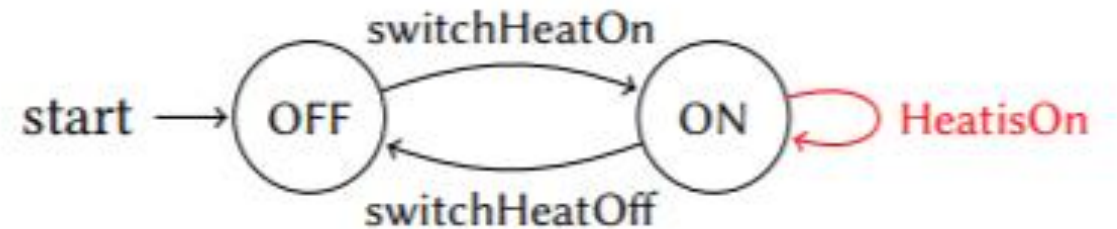


Fig. 7. Heater process.

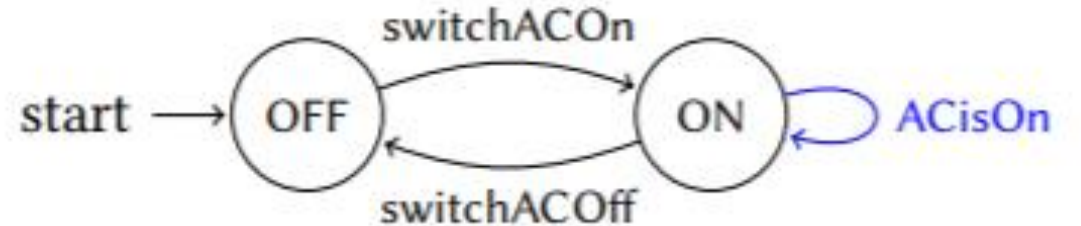


Fig. 8. AC process.

# Thermostat case study

## Phase I

Maintain ambient temperature

Thermostat  
**doesn't interact**  
with all devices

+

## Phase II

Thermostat must  
interact with all  
devices (Fairness)

AC and Heater are  
**switched on at**  
**the same time**

+

## Phase III

Energy efficient

**SATISFIED!**



# Coordination synthesis, in a nutshell

- Modelling
  - CSP environments – Interface, non-determinism, private/public actions
  - LTL specifications – Expressive
- Algorithm
  - Efficient automata-based reduction to synchronous synthesis
  - Prototype + Case studies demonstrate utility
- Complexity analysis
  - PSPACE-hard in size of  $E$
  - Algorithm is exponential in  $E$ , number of devices