

# Probabilistic Systems

## Part 1: discrete time Markov chains

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# Discrete-time Markov chains

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Discrete-time Markov chains (DTMCs)

Paths and probabilities for DTMCs

Probabilistic reachability for DTMCs

Rewards and expected reachability for DTMCs

# Discrete-time Markov chains

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## Kripke structures augmented with probabilities

### States:

- represent possible configurations of the system being modelled
- labelled by atomic propositions (properties that hold in the states)

### Transitions:

- model evolution of a system's state
- occur in **discrete time-steps**

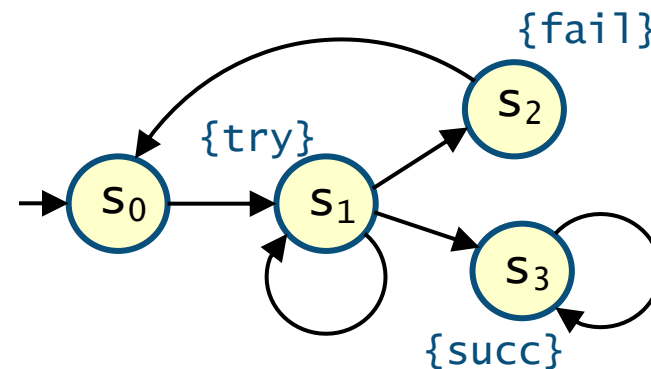
### Probabilities:

- likelihood of making transitions between states are given by **discrete probability distributions**

# Kripke structures

A Kripke structure is a tuple  $(S, s_0, T, L)$  where

- $S$  is a finite set of states
- $s_0$  is the initial state
- $T \subseteq S \times S$  is the **transition relation**  
if  $(s, s') \in T$ , then there is a transition from  $s$  to  $s'$
- $L: S \rightarrow 2^{AP}$  is the labelling function where  $AP$  is a set of atomic propositions

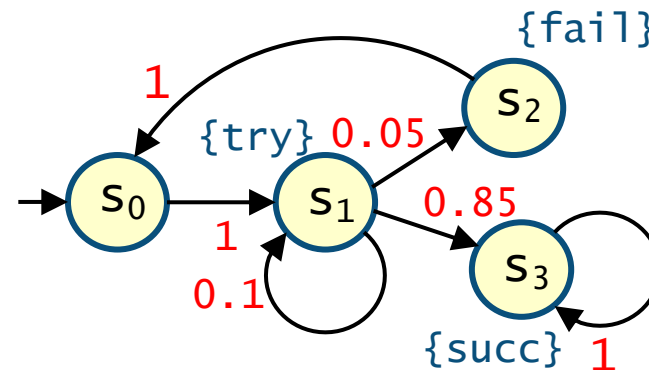


# Discrete-time Markov chains

A discrete-time Markov chain is a tuple  $(S, s_0, \mathbf{P}, L)$  where:

- $S$  is a finite set of states
- $s_0$  is the initial state
- $\mathbf{P}: S \times S \rightarrow [0, 1]$  is the **transition probability matrix** where  $\mathbf{P}(s, s')$  is the probability of making a transition from  $s$  to  $s'$ 
  - we require that  $\sum_{s' \in S} \mathbf{P}(s, s') = 1$  for all states  $s \in S$
- $L: S \rightarrow 2^{AP}$  is the labelling function where  $AP$  is a set of atomic propositions

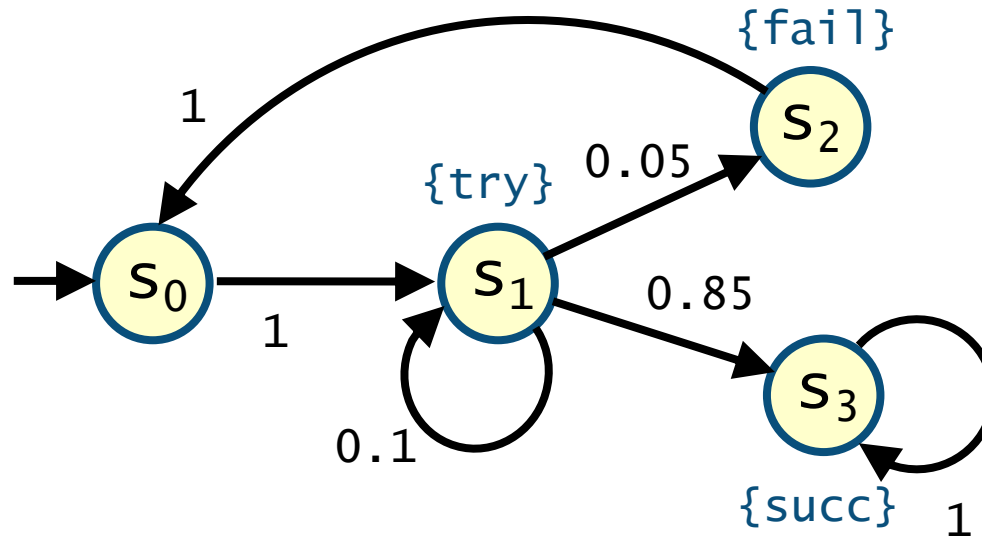
i.e. require that the total probability of making a transition from any state is **1**



# Simple DTMC example

## Modelling a very simple communication protocol

- after one step, process starts **trying** to send a message
- with probability **0.1**, channel unready so wait a step
- with probability **0.85**, send message **successfully** and stop
- with probability **0.05**, sending **fails**, then in next step it restarts



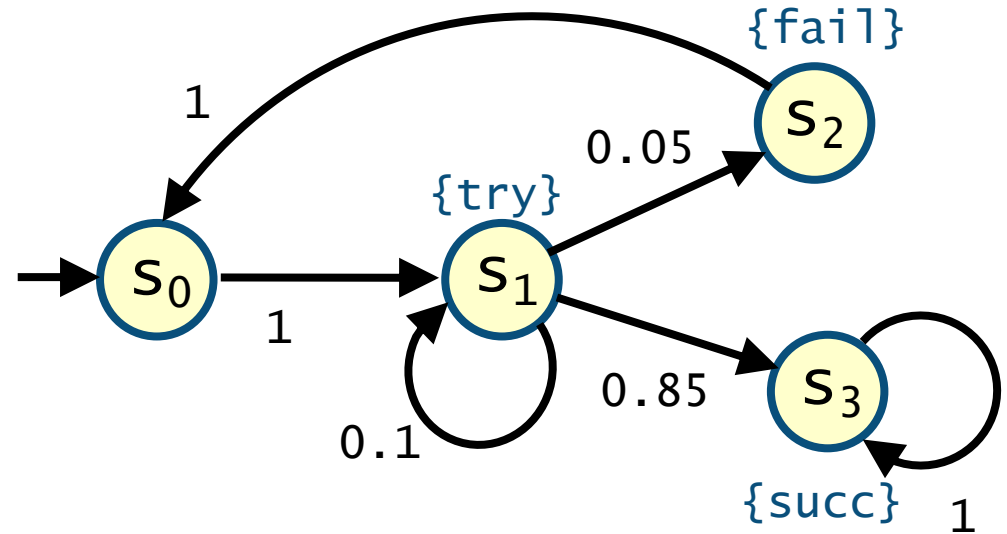
# Simple DTMC example

States:  $S = \{s_0, s_1, s_2, s_3\}$

Initial state:  $s_0$

Probability transition matrix:

$$P = \begin{bmatrix} 0 & 1 & 0 & 0 \\ 0 & 0.1 & 0.05 & 0.85 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$



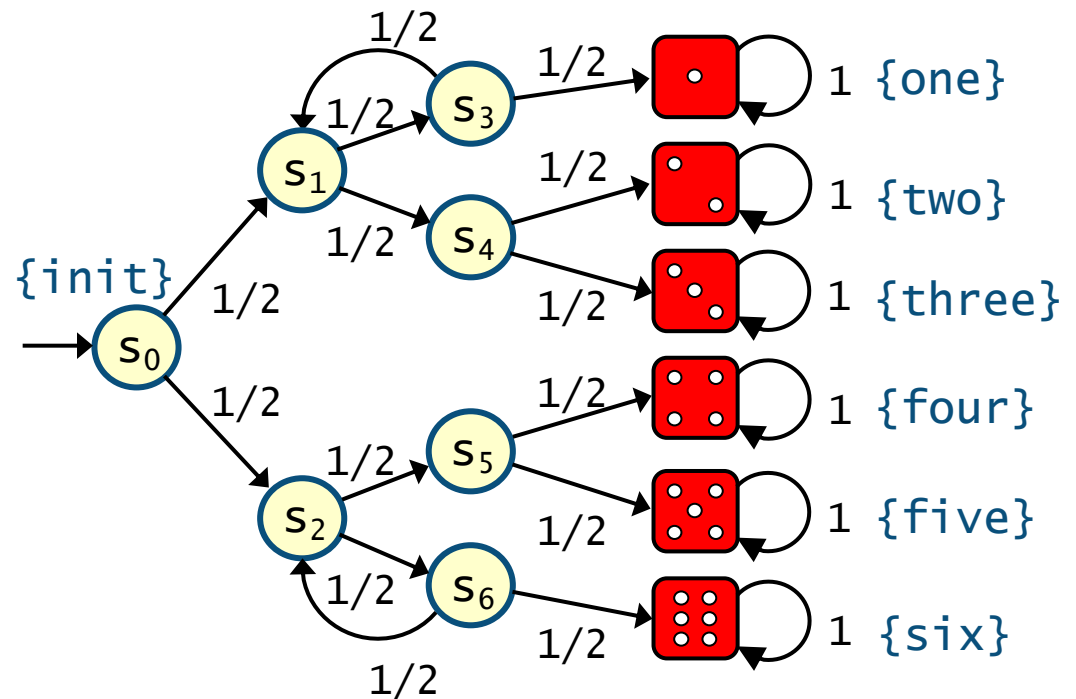
Atomic propositions:  $AP = \{\text{try}, \text{fail}, \text{succ}\}$

– labelling:  $L(s_0) = \emptyset$ ,  $L(s_1) = \{\text{try}\}$ ,  $L(s_2) = \{\text{fail}\}$  and  $L(s_3) = \{\text{succ}\}$

# DTMC example 2 – Coins and dice

## Modelling a 6-sided die using a fair coin

- algorithm due to Knuth/Yao
- start at  $s_0$ , flip a coin
- upper branch when flip **H**
- lower branch when flip **T**
- repeat until value chosen





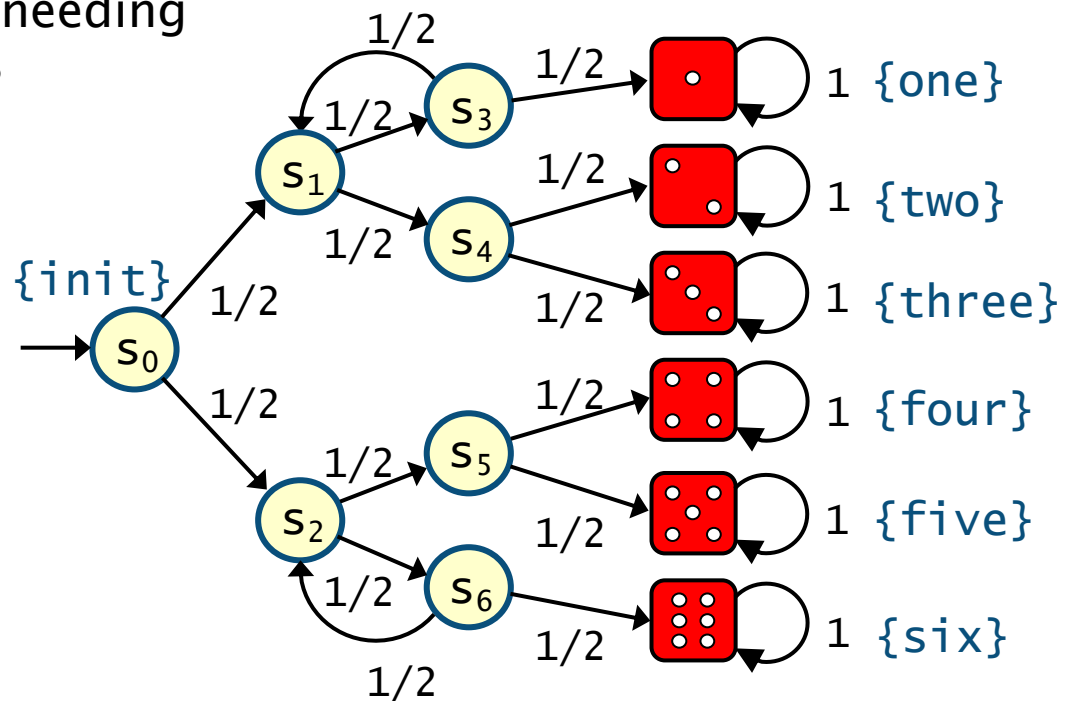
# DTMC example 2 – Coins and dice

## Is this model correct?

- e.g. probability of obtaining a 4 equals  $1/6$
- is it guaranteed to terminate?

## How efficient is it?

- what is the probability of needing more than four coin flips?
- on average, how many coin flips are needed?



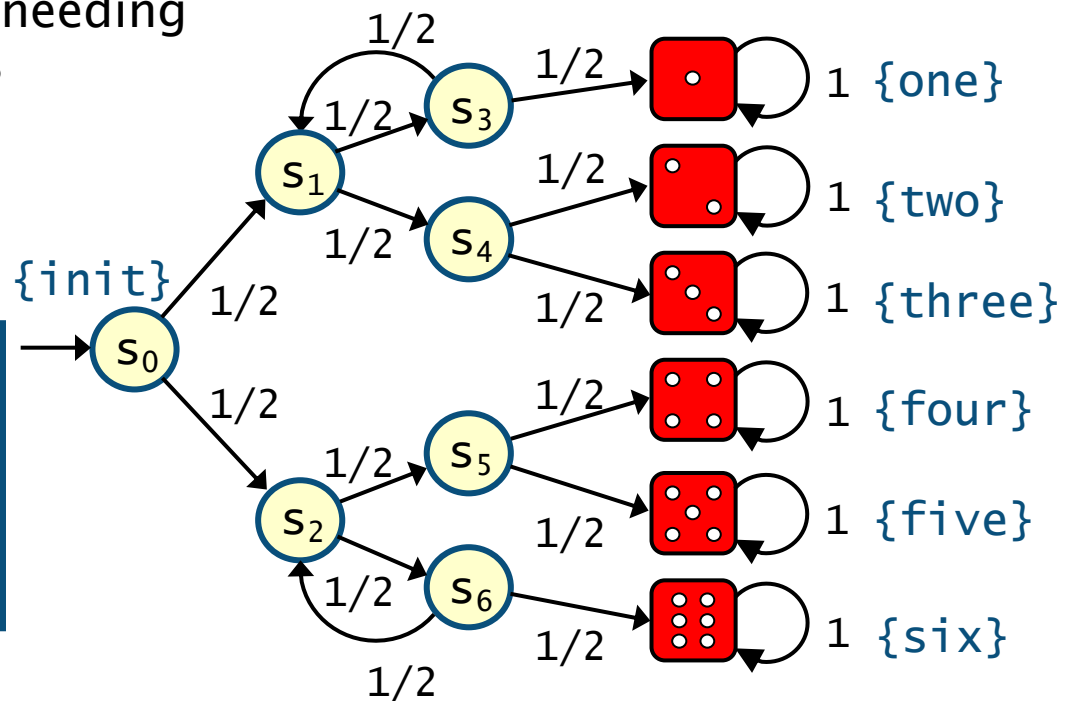
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Probabilistic model checking provides a framework for answering these kinds of questions

# Discrete time Markov chains

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Discrete-time Markov chains (DTMCs)

**Paths and probabilities for DTMCs**

Probabilistic reachability for DTMCs

Rewards and expected reachability for DTMCs

# First some probability basics

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## Need an experiment...

- the **sample space** is the set of possible outcomes of the experiment
- an **event** is a subset of the sample space
- the **probability** of an event is the degree of certainty an event will occur

## Example: toss two coins

- sample space:  $\{(H,H), (H,T), (T,H), (T,T)\}$
- event: “at least one H”
- probability:  $1/2 + (1/2) \cdot (1/2) = 3/4$

## Example: toss a coin infinitely often

- sample space: set of infinite sequences of H/T
- event: “H in the first 3 throws”
- probability:  $1/2 + (1/2) \cdot (1/2) + (1/2) \cdot (1/2) \cdot (1/2) = 7/8$

# Probability space $(\Omega, \Sigma, \text{Prob})$

Sample space  $\Omega$  is an arbitrary non-empty set

Event set  $\Sigma$  is family of subsets of  $\Omega$  which is

- closed under complementation
  - if  $A$  is in  $\Sigma$ , then the complement  $\Omega \setminus A$  is in  $\Sigma$
- closed under countable union
  - if  $A_i$  is in  $\Sigma$  for  $i \in \mathbb{N}$ , then the union  $\cup_i A_i$  is in  $\Sigma$
- contains the empty set ( $\emptyset$  is in  $\Sigma$ )

Elements of  $\Sigma$  are called **measurable sets** and  $\Sigma$  a  $\sigma$ -algebra on  $\Omega$

Probability measure **Prob** is a function **Prob**:  $\Sigma \rightarrow [0, 1]$  such that

- **Prob**( $\Omega$ ) = 1
- **Prob**( $\cup A_i$ ) =  $\sum_i$  **Prob**( $A_i$ )  
for any disjoint family of measurable sets  $A_1, A_2, \dots$

# Probability space – Simple example

Sample space  $\Omega = \mathbb{N} = \{ 0, 1, 2, 3, 4, \dots \}$

- the natural numbers

Event set  $\Sigma = \{ \emptyset, \text{“odd”}, \text{“even”}, \mathbb{N} \}$

- (closed under complement/countable union, contains  $\emptyset$ )
- e.g.  $\text{“odd”} \cup \text{“even”} = \mathbb{N}$  and  $\mathbb{N} \setminus \text{“odd”} = \text{“even”}$

Probability measure **Prob**

- e.g. corresponding to picking a number uniformly at random
- **Prob(“odd”)=1/2, Prob(“even”)=1/2, ...**

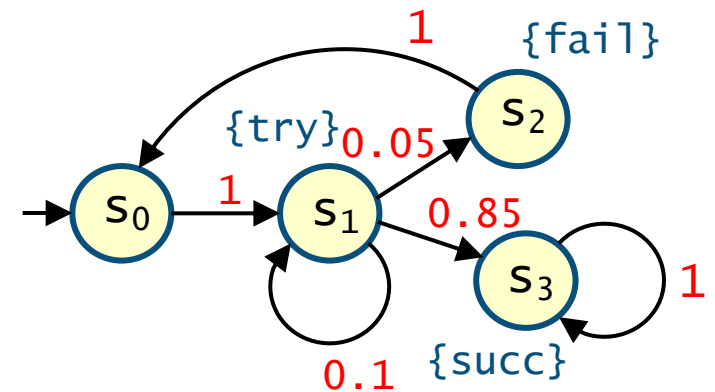
# Back to DTMCs – Paths

A **path** in a DTMC represents an **execution** of the system

- i.e. one possible behaviour

**Formally:**

- infinite sequence of states  $s_0s_1s_2s_3\dots$   
such that  $P(s_i, s_{i+1}) > 0$  for all  $i \geq 0$



# Back to DTMCs – Paths

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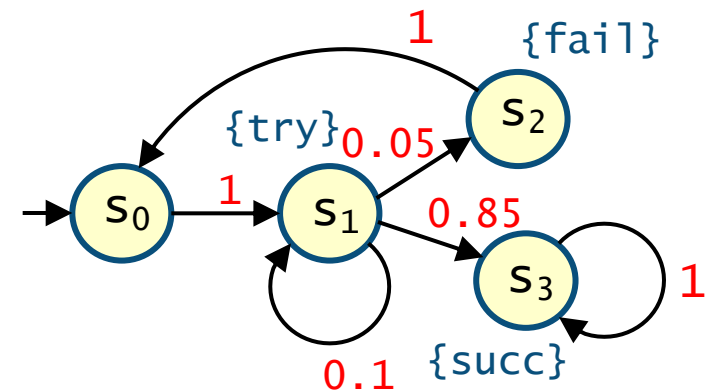
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**Example execution:**

- start, wait, fail, retry, start, succeed





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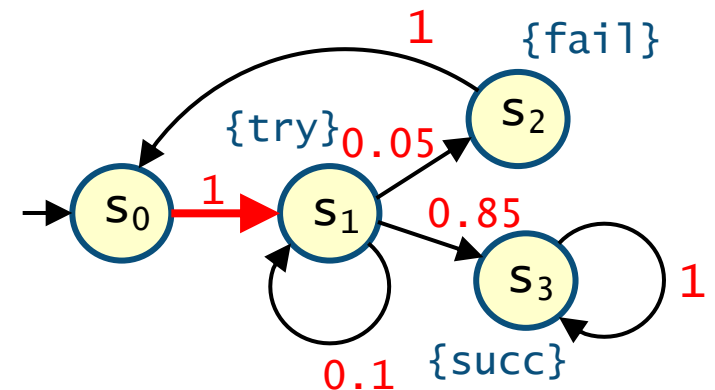
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**Example execution:**

- **start**, wait, fail, retry, start, succeed:  $s_0s_1$



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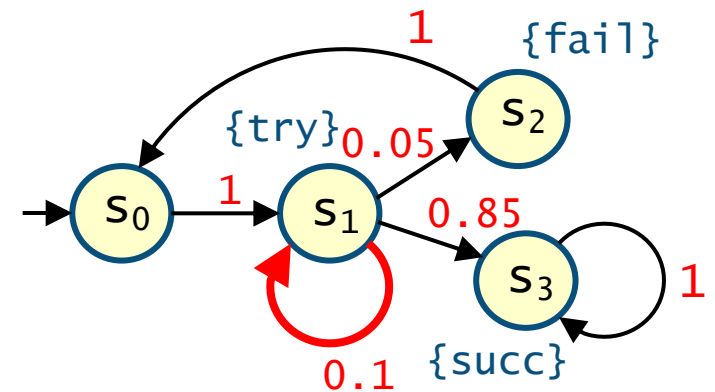
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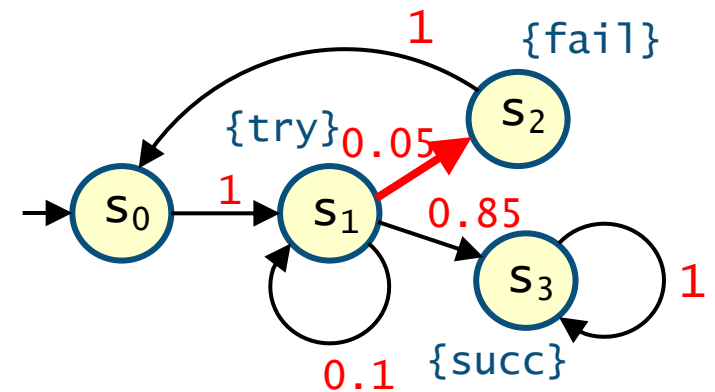
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**Example execution:**

- start, wait, **fail**, retry, start, succeed:  $s_0s_1s_1s_2$



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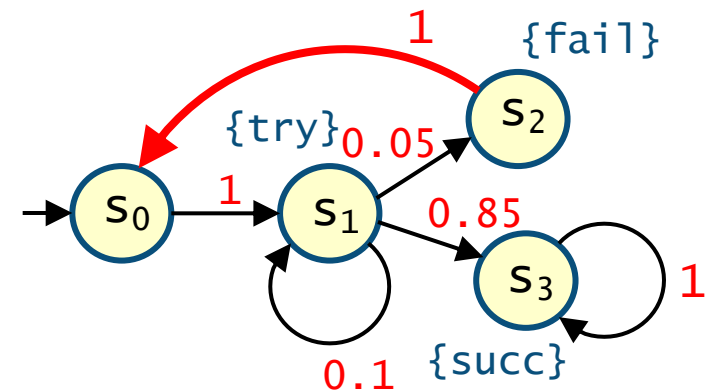
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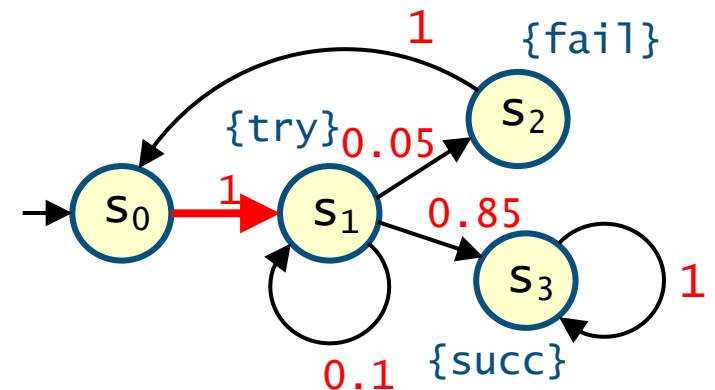
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**Example execution:**

- start, wait, fail, retry, **start**, succeed:  $s_0s_1s_1s_2s_0s_1$



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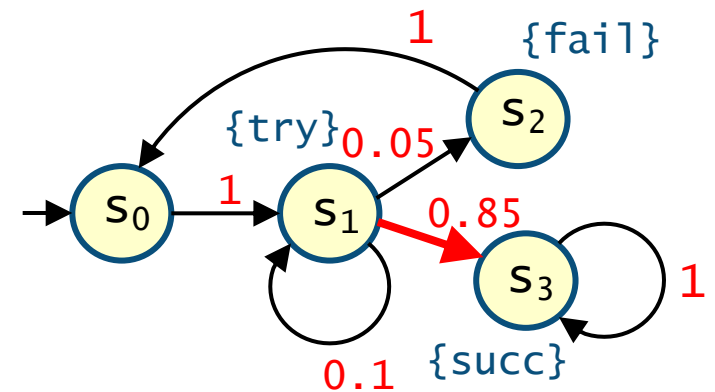
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**Example execution:**

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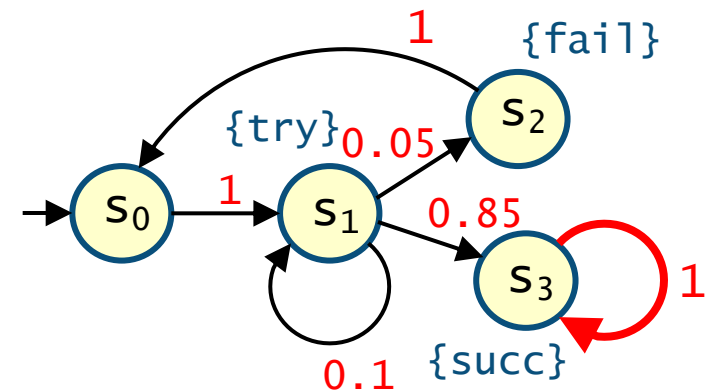
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**Formally:**

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**Example execution:**

- start, wait, fail, retry, start, **succeed**:  $s_0s_1s_1s_2s_0s_1s_3s_3\dots$



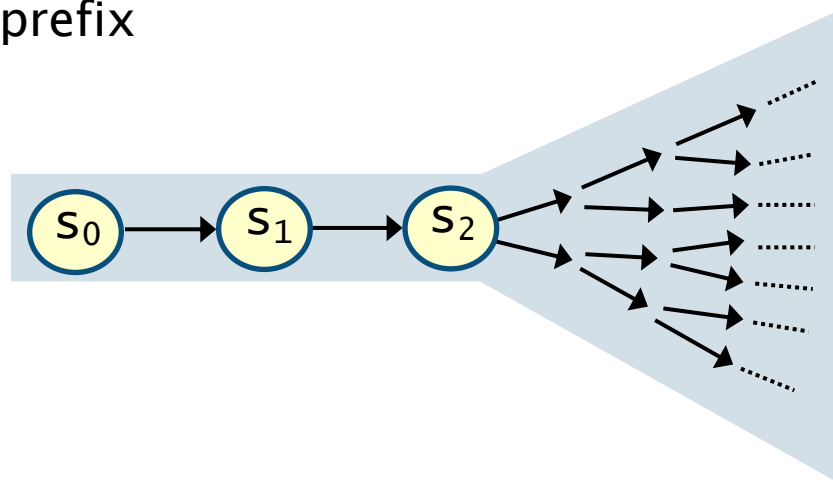
# Probability space over paths

To reason about a DTMC when starting from some state  $s$

- need to define a **probability space over paths** starting from the state  $s$

**Intuitively:**

- sample space: infinite paths starting from the state  $s$
- events: sets of infinite paths
- basic events: **cylinder sets**
- cylinder  $\text{Cyl}(\omega)$  for a finite path  $\omega$  equals the set of infinite paths that have  $\omega$  as a prefix
- e.g.  $\text{Cyl}(ss_1s_2)$





# Probability space over paths

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Probability space ( $\text{Path}_s$ ,  $\Sigma_s$ ,  $\text{Prob}_s$ )

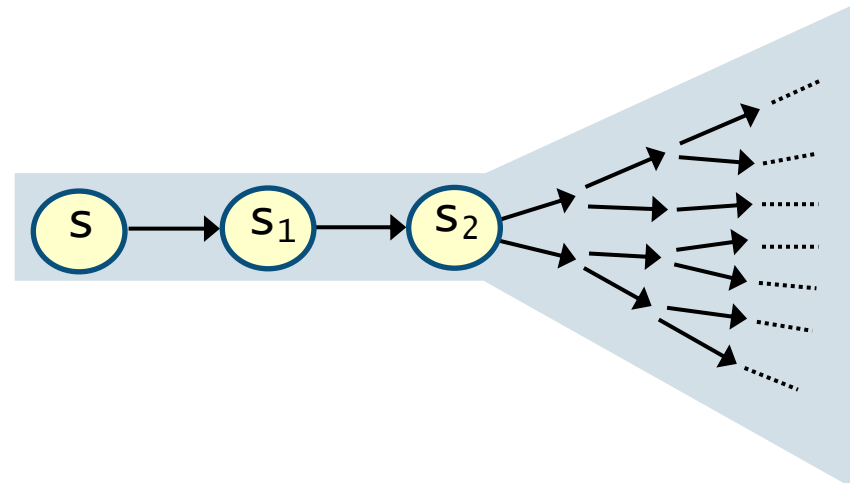
**Sample space:** all infinite paths starting from the state  $s$

# Probability space over paths

Probability space  $(\text{Path}_s, \Sigma_s, \text{Prob}_s)$

Sample space: all infinite paths starting from the state  $s$

**Event set:** least  $\sigma$ -algebra including the cylinder  $\text{Cyl}(\omega)$  of every finite path  $\omega = s s_1 s_2 \dots s_n$



# Probability space over paths

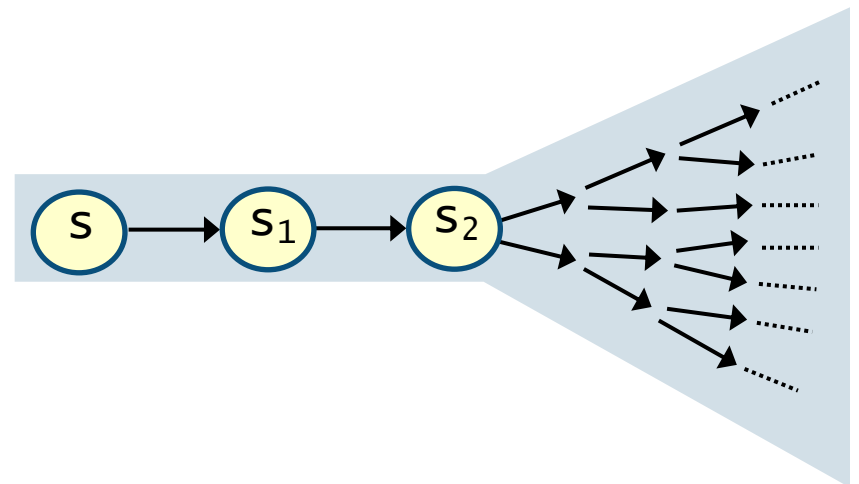
Probability space  $(\text{Path}_s, \Sigma_s, \text{Prob}_s)$

Sample space: all infinite paths starting from the state  $s$

**Event set:** least  $\sigma$ -algebra including the cylinder  $\text{Cyl}(\omega)$  of every finite path  $\omega = s s_1 s_2 \dots s_n$

**Probability measure:** unique extension of function  $\text{Prob}_s$  over cylinders where  $\text{Prob}_s(\text{Cyl}(\omega)) = P(s, s_1) \cdot P(s_1, s_2) \dots P(s_{n-1}, s_n)$

probability of a cylinder given by multiplying the probability of each transition of the finite path



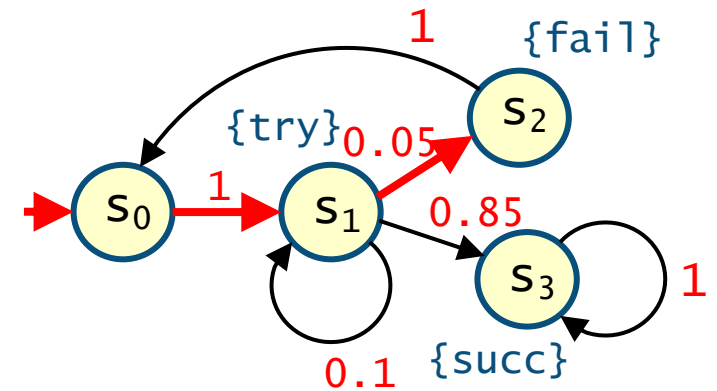
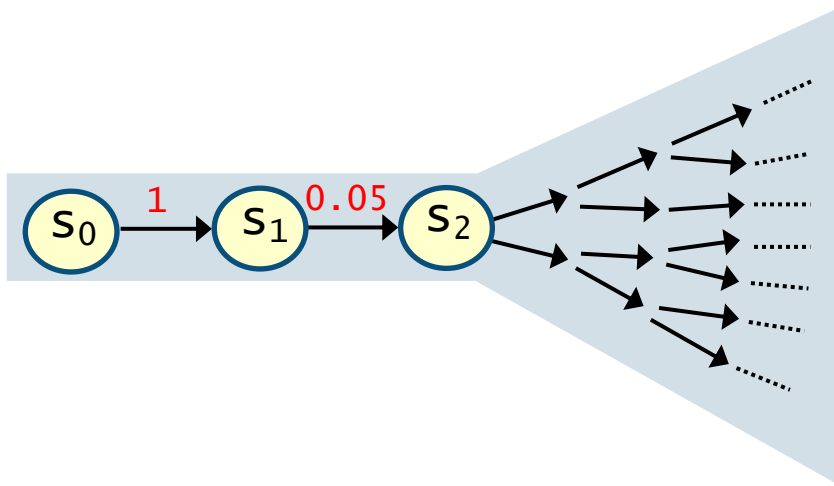
# Paths and probabilities – Example

## Paths where sending fails the first time

- all paths starting  $s_0s_1s_2$ , i.e. the cylinder  $Cyl(s_0s_1s_2)$

## Probability:

$$\mathbf{Prob}_{s_0}(Cyl(s_0s_1s_2)) = \mathbf{P}(s_0, s_1) \cdot \mathbf{P}(s_1, s_2) = 1 \cdot 0.05 = 0.05$$



# Paths and probabilities – Example

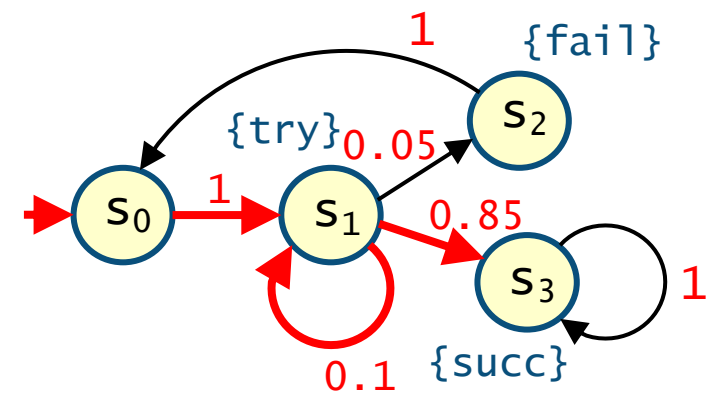
## Paths which are eventually successful with no failures

- infinite paths of the form  $s_0 (s_1)^* s_3^\omega$
- i.e. the (disjoint) union of the cylinders:

$$\text{Cyl}(s_0 s_1 s_3) \cup \text{Cyl}(s_0 s_1 s_1 s_3) \cup \text{Cyl}(s_0 s_1 s_1 s_1 s_3) \cup \dots$$

## Probability:

$$\text{Prob}_{s_0}(\text{Cyl}(s_0 s_1 s_3) \cup \text{Cyl}(s_0 s_1 s_1 s_3) \cup \text{Cyl}(s_0 s_1 s_1 s_1 s_3) \cup \dots)$$



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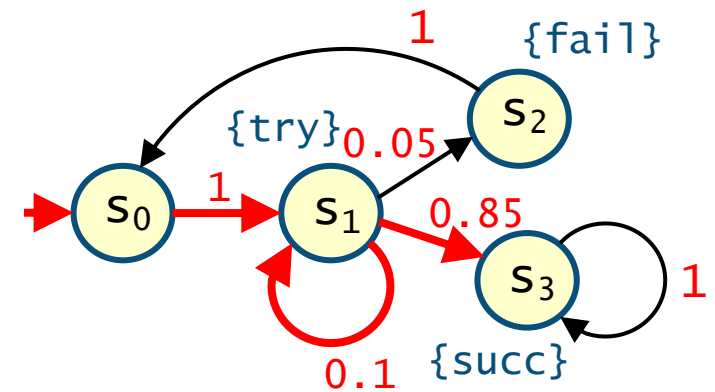
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$$\text{Prob}_{s_0}(\text{Cyl}(s_0 s_1 s_3)) + \text{Prob}_{s_0}(\text{Cyl}(s_0 s_1 s_1 s_3)) + \text{Prob}_{s_0}(\text{Cyl}(s_0 s_1 s_1 s_1 s_3)) + \dots$$

since the sets are disjoint



# Paths and probabilities – Example

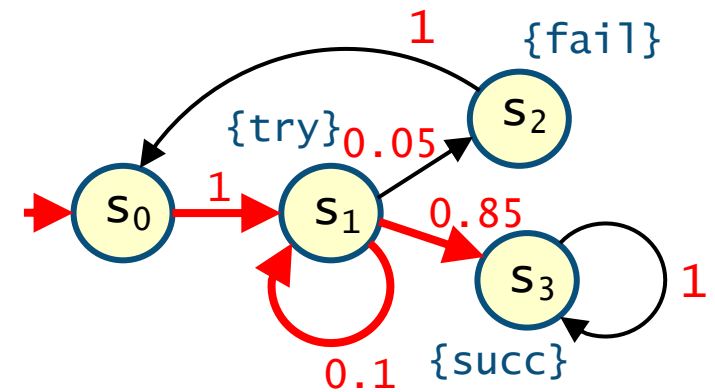
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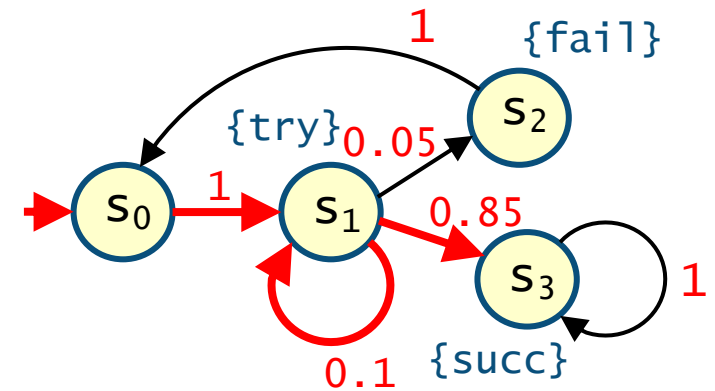
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# Paths and probabilities – Example

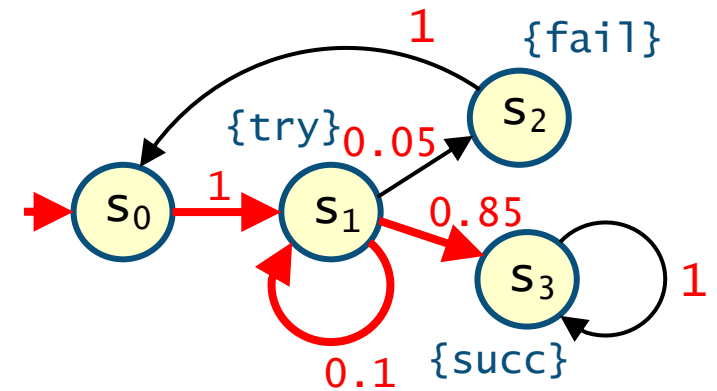
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## Probability:

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# Discrete time Markov chains

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Discrete-time Markov chains (DTMCs)

Paths and probabilities for DTMCs

**Probabilistic reachability for DTMCs**

Rewards and expected reachability for DTMCs

# Probabilistic reachability

## Fundamental property of DTMCs: **probabilistic reachability**

- probability of a path reaching some target set of states  $T$ 
  - $P(s, T)$  probability of reaching  $T$  from state  $s$
  - vector:  $P(T)$  values for all states of a DTMC
- e.g. “probability of the algorithm terminating successfully?”
- e.g. “probability that an error occurs during execution?”

## Dual of reachability: **invariance**

- probability of remaining within some class of states
- $\text{Prob}(\text{“remain in set } I\text{”}) = 1 - \text{Prob}(\text{“reach set } S \setminus I\text{”})$
- e.g. “probability that an error never occurs”

## Also other variants of reachability

- step-bounded, constrained (“until”), ...

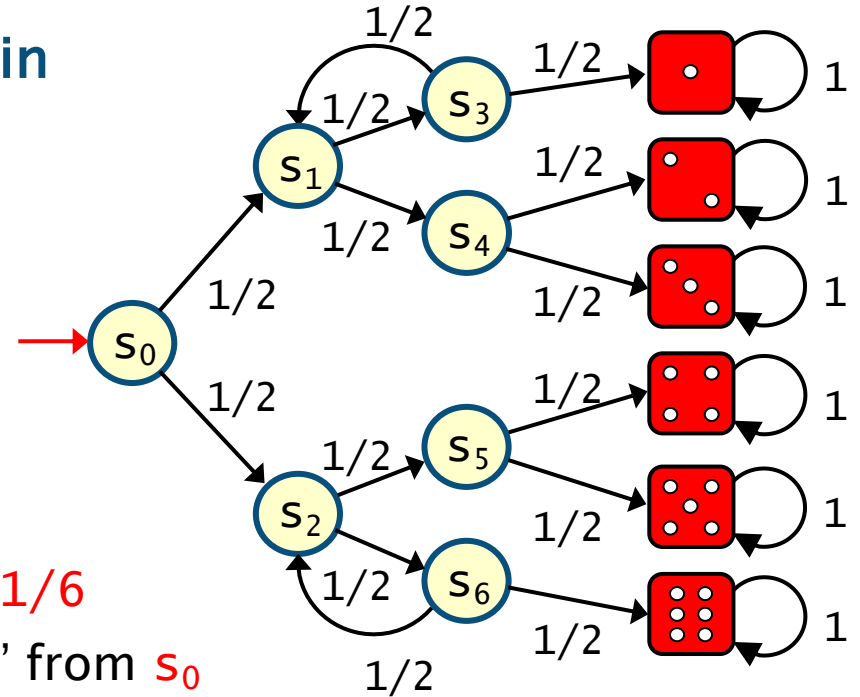
# Probabilistic reachability – Example

## Modelling a 6-sided die using a fair coin

- algorithm due to Knuth/Yao
- start at  $s_0$ , toss a coin
- upper branch when **H**
- lower branch when **T**
- repeat until value chosen

## Is this algorithm correct?

- e.g. probability of reaching “4” equals  $1/6$
- event: all possible ways of reaching “4” from  $s_0$



# Probabilistic reachability – Example

## Modelling a 6-sided die using a fair coin

- algorithm due to Knuth/Yao
- start at  $s_0$ , toss a coin
- upper branch when **H**
- lower branch when **T**
- repeat until value chosen

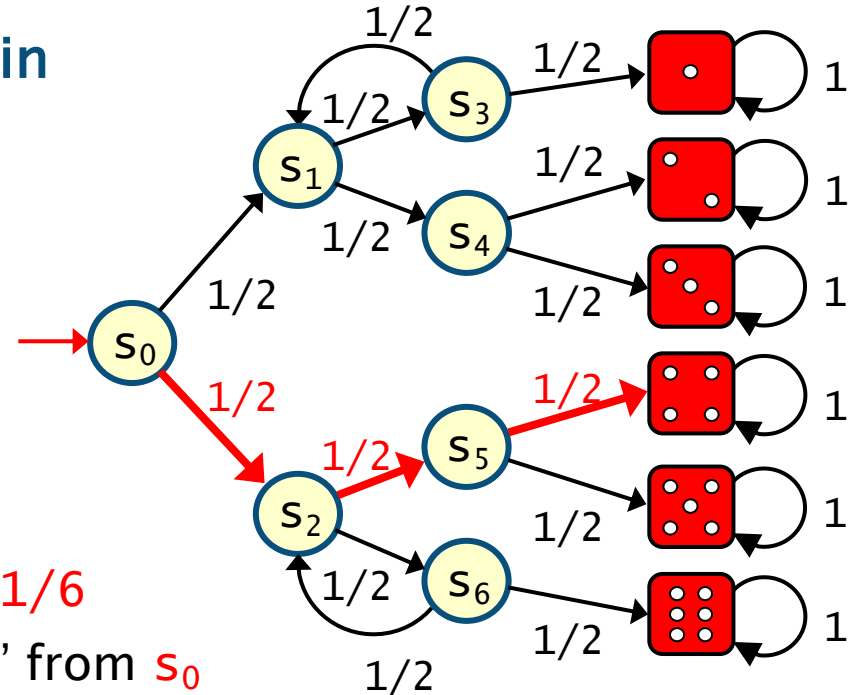
## Is this algorithm correct?

- e.g. probability of reaching “4” equals  $1/6$
- event: all possible ways of reaching “4” from  $s_0$
- ways of reaching “4” :

**THH**,

- probability of reaching “4” :

$(1/2)^3 +$



# Probabilistic reachability – Example

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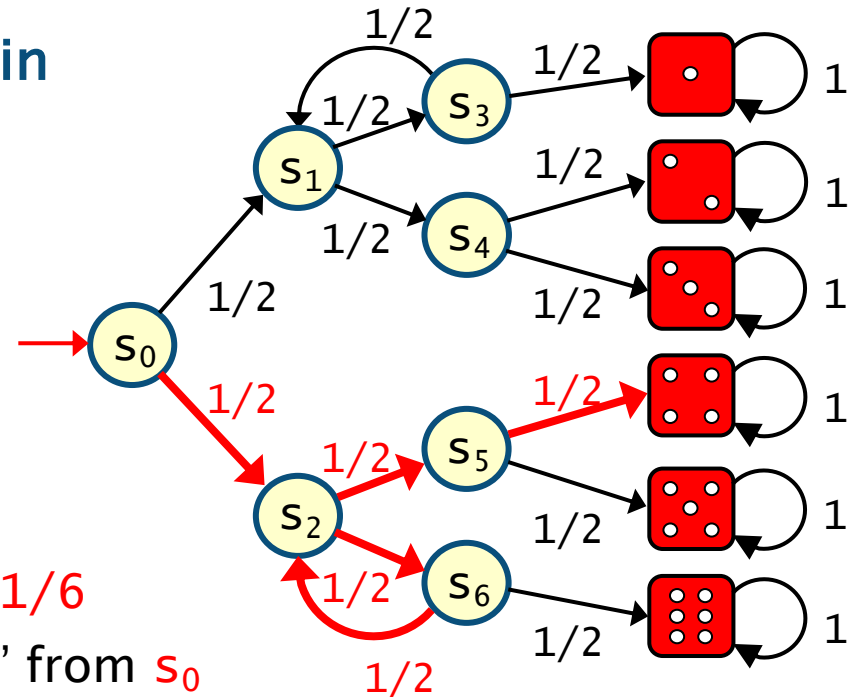
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- event: all possible ways of reaching “4” from  $s_0$
- ways of reaching “4” :

THH, **TTTHH**,

- probability of reaching “4” :

$$(1/2)^3 + (1/2)^5 +$$



# Probabilistic reachability – Example

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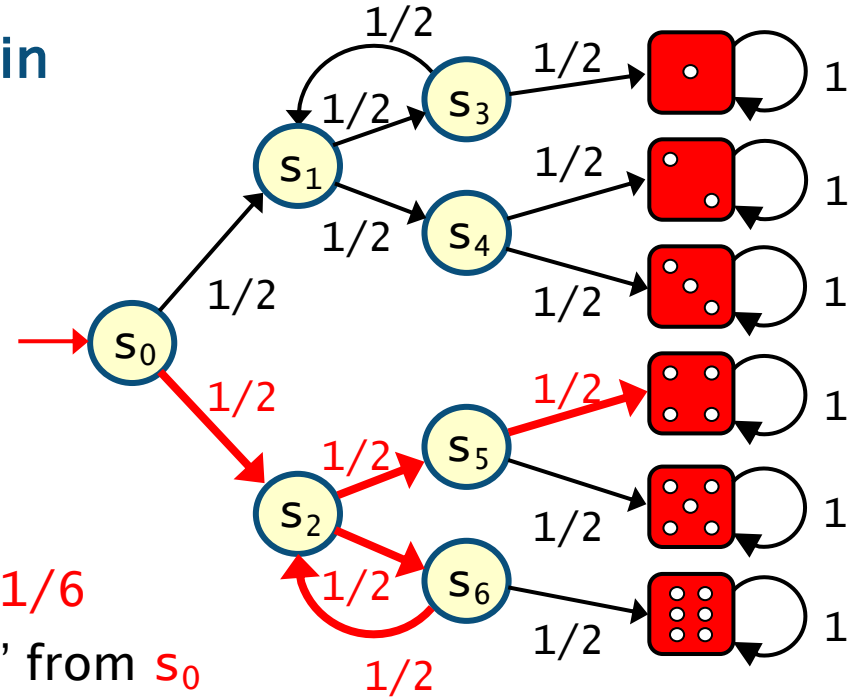
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- e.g. probability of reaching “4” equals  $1/6$
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- ways of reaching “4” :

THH, TTTHH, **TTTTTHH**, ...

- probability of reaching “4” :

$$(1/2)^3 + (1/2)^5 + (1/2)^7 + \dots$$



# Probabilistic reachability – Example

## Modelling a 6-sided die using a fair coin

- algorithm due to Knuth/Yao
- start at  $s_0$ , toss a coin
- upper branch when **H**
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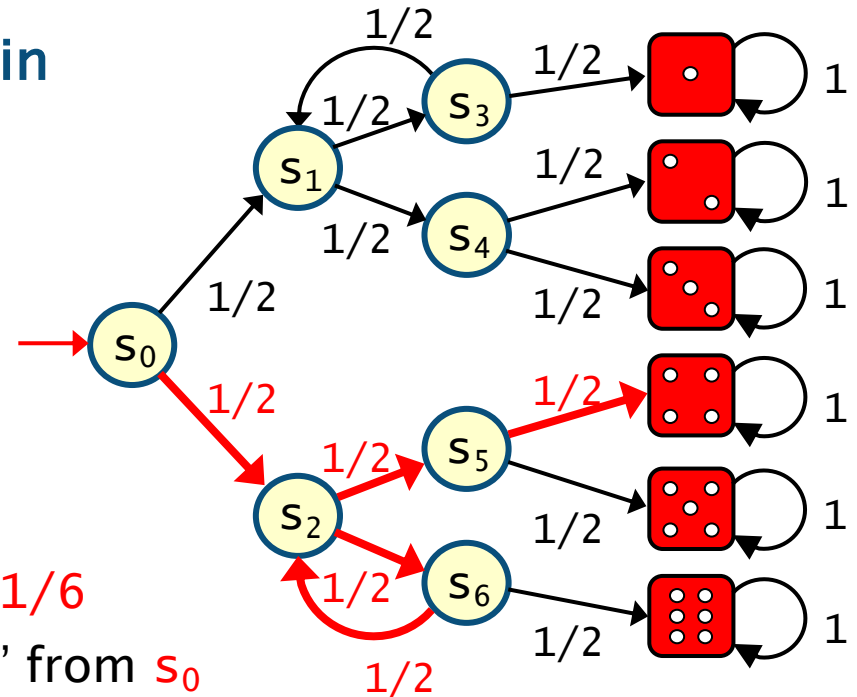
## Is this algorithm correct?

- e.g. probability of reaching “4” equals  $1/6$
- event: all possible ways of reaching “4” from  $s_0$
- ways of reaching “4” :

THH, TTTHH, TTTTTHH, ... ,  $T(TT)^nHH$  , ...

- probability of reaching “4” :

$$(1/2)^3 + (1/2)^5 + (1/2)^7 + \dots + (1/2)^{2n+3} + \dots = 1/6$$





# Probabilistic reachability – Computation

---

Computing an infinite sum not not feasible in practice

Alternative to calculate  $P(s, T)$ : derive a **linear equation system**

- calculate probabilities for all states  $s \in S$  simultaneously

Let  $x_s$  denote the probability of reaching  $T$  from state  $s$

# Probabilistic reachability – Computation

---

Computing an infinite sum not not feasible in practice

Alternative to calculate  $P(s, T)$ : derive a **linear equation system**

- calculate probabilities for all states  $s \in S$  simultaneously

Let  $x_s$  denote the probability of reaching  $T$  from state  $s$

- if  $s \in T$ , then  $x_s = 1$

if the state is in the target set then the probability of reaching the target is **1**

# Probabilistic reachability – Computation

Computing an infinite sum not not feasible in practice

Alternative to calculate  $P(s, T)$ : derive a **linear equation system**

- calculate probabilities for all states  $s \in S$  simultaneously

Let  $x_s$  denote the probability of reaching  $T$  from state  $s$

- if  $s \in T$ , then  $x_s = 1$
- if  $T$  is not reachable from  $s$ , then  $x_s = 0$   
i.e. no (finite) path from  $s$  to a state in  $T$

if one cannot reach the target, then the probability of reaching the target is  $0$

# Probabilistic reachability – Computation

Computing an infinite sum not not feasible in practice

Alternative to calculate  $P(s, T)$ : derive a **linear equation system**

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- if  $T$  is not reachable from  $s$ , then  $x_s = 0$   
i.e. no (finite) path from  $s$  to a state in  $T$
- otherwise  $x_s = \sum_{s' \in S} P(s, s') \cdot x_{s'}$

probability defined recursively using the transition probabilities:  
summation over all states  $s'$  of the probability of making a transition to  $s'$   
multiplied by the probability of reaching the target from  $s'$


# Probabilistic reachability – Computation

Can view as a **least fixed point** computation over vectors  $y \in [0, 1]^S$

– consider the function  $F : [0, 1]^S \rightarrow [0, 1]^S$  where

$$F(y)(s) = \begin{cases} 1 & \text{if } s \in T \\ \sum_{s' \in S} P(s, s') \cdot y(s') & \text{otherwise} \end{cases}$$

a vector of probabilities in each state



If we let  $x^{(0)} = \mathbf{0}$  and  $x^{(n+1)} = F(x^{(n)})$  then we have that

- $x^{(0)} \leq x^{(1)} \leq x^{(2)} \leq x^{(3)} \leq \dots$
- $P(T) = \lim_{n \rightarrow \infty} x^{(n)}$ 
  - recall  $P(T)$  is the vector of probabilities  $(P(s, T))_{s \in S}$
- $P(T)$  is the least fixed point of  $F$

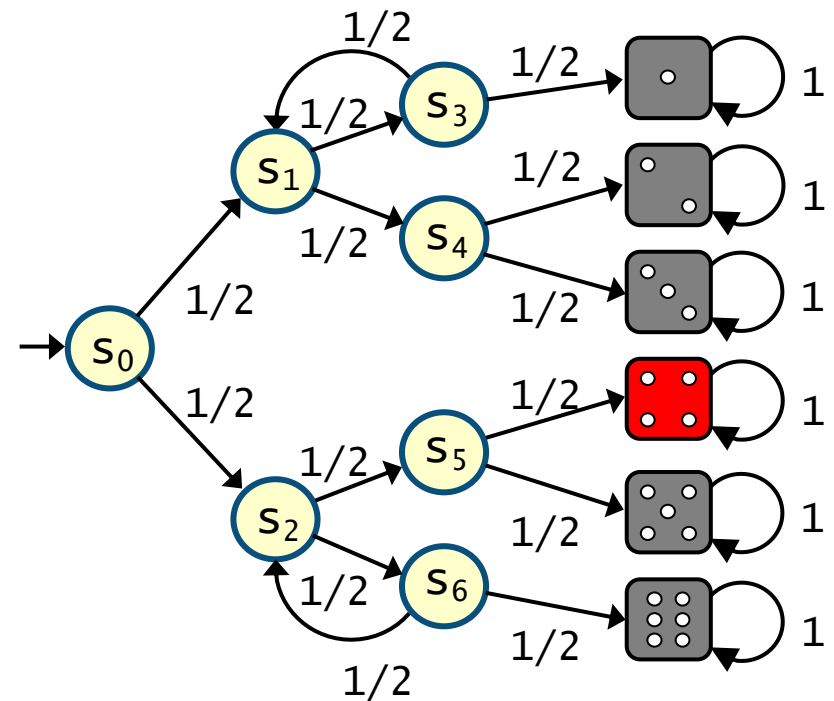
# Probabilistic reachability – Example

## Probability of reaching {4}

- i.e. tossing a four

$x_s$  denotes the probability of reaching  $T$  from  $s$

- if  $s \in T$ , then  $x_s = 1$
- if  $T$  is not reachable from  $s$ , then  $x_s = 0$
- otherwise  $x_s = \sum_{s' \in S} P(s, s') \cdot x_{s'}$



# Probabilistic reachability – Example

## Probability of reaching $\{4\}$

– i.e. tossing a four

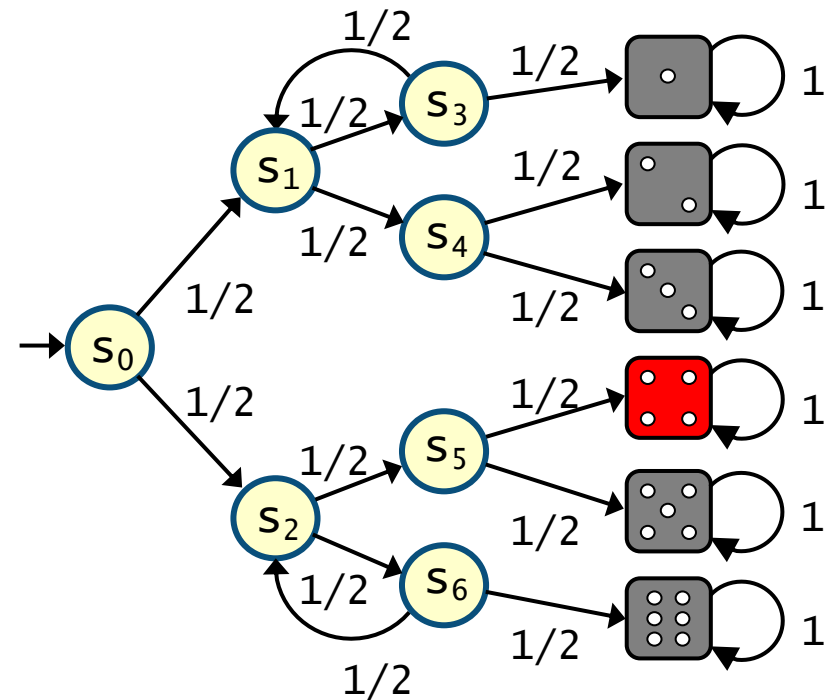
•  $x_4 = 1$

$x_s$  denotes the probability of reaching  $T$  from  $s$

– if  $s \in T$ , then  $x_s = 1$

– if  $T$  is not reachable from  $s$ , then  $x_s = 0$

– otherwise  $x_s = \sum_{s' \in S} P(s, s') \cdot x_{s'}$



# Probabilistic reachability – Example

## Probability of reaching {4}

– i.e. tossing a four

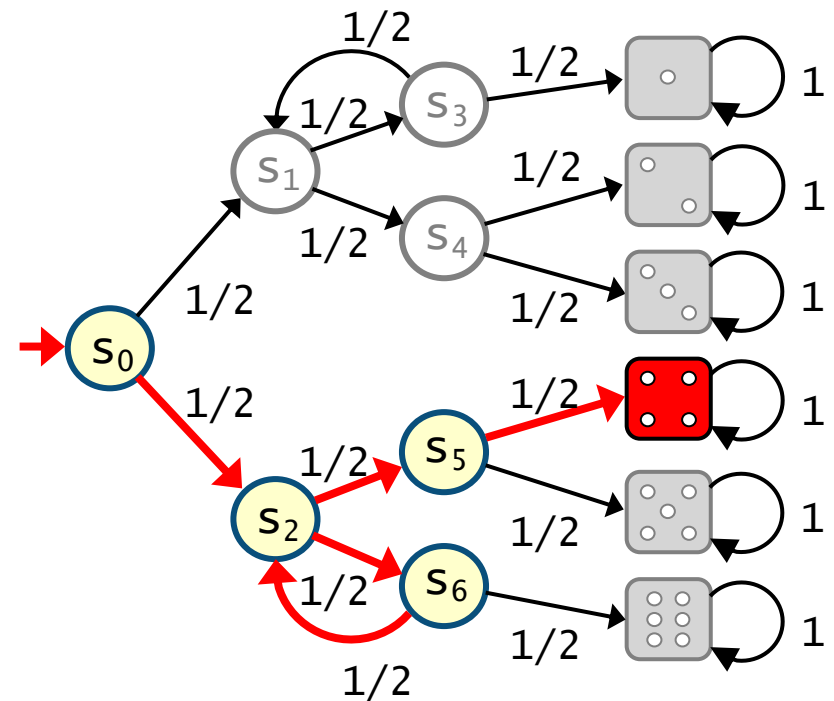
- $x_4 = 1$
- only  $s_0, s_2, s_5$  and  $s_6$  reach {4} therefore  $x_s = 0$  for all other states

$x_s$  denotes the probability of reaching  $T$  from  $s$

– if  $s \in T$ , then  $x_s = 1$

– if  $T$  is not reachable from  $s$ , then  $x_s = 0$

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# Probabilistic reachability – Example

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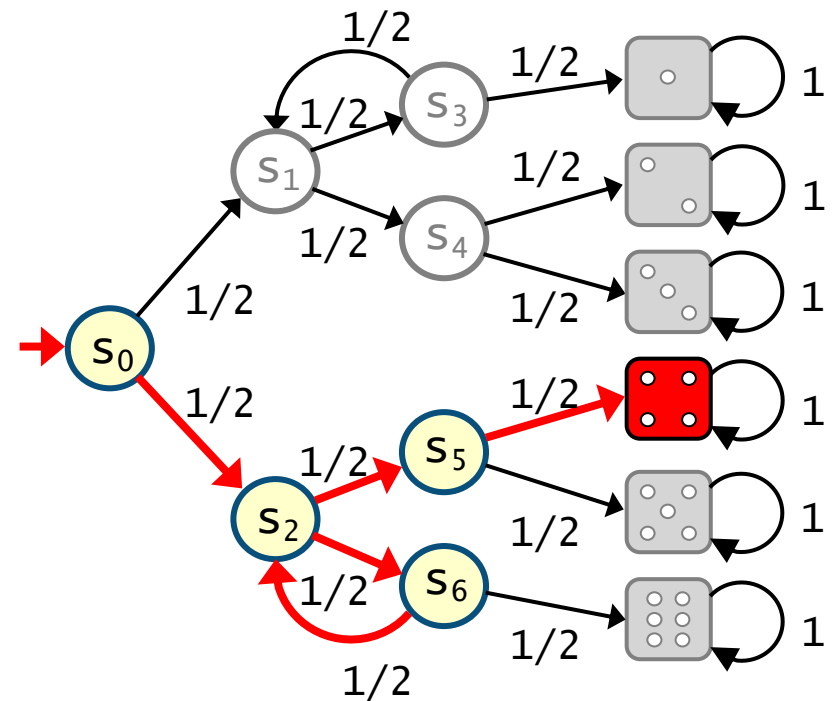
- $x_4 = 1$
- only  $s_0, s_2, s_5$  and  $s_6$  reach {4}  
therefore  $x_s = 0$  for all other states
- $x_{s_0} = 1/2 \cdot 0 + 1/2 \cdot x_{s_2}$
- $x_{s_2} = 1/2 \cdot x_{s_5} + 1/2 \cdot x_{s_6}$
- $x_{s_5} = 1/2 \cdot 1 + 1/2 \cdot 0$
- $x_{s_6} = 1/2 \cdot x_{s_2} + 1/2 \cdot 0$

$x_s$  denotes the probability of reaching T from s

– if  $s \in T$ , then  $x_s = 1$

– if T is not reachable from s, then  $x_s = 0$

– otherwise  $x_s = \sum_{s' \in S} P(s, s') \cdot x_{s'}$



# Probabilistic reachability – Example

## Probability of reaching {4}

– i.e. tossing a four

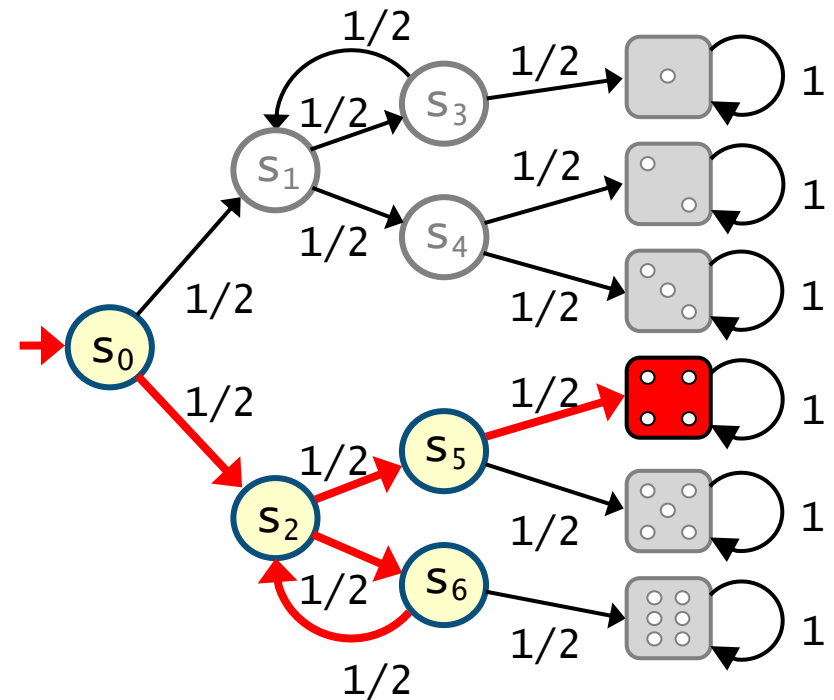
- $x_4 = 1$
- only  $s_0, s_2, s_5$  and  $s_6$  reach {4}  
therefore  $x_s = 0$  for all other states

- $x_{s_0} = 1/2 \cdot x_{s_2}$
- $x_{s_2} = 1/2 \cdot x_{s_5} + 1/2 \cdot x_{s_6}$
- $x_{s_5} = 1/2$
- $x_{s_6} = 1/2 \cdot x_{s_2}$

simplifying

$x_s$  denotes the probability of reaching  $T$  from  $s$

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- if  $T$  is not reachable from  $s$ , then  $x_s = 0$
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# Probabilistic reachability – Example

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– i.e. tossing a four

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- $x_4 = 1$
- only  $s_0, s_2, s_5$  and  $s_6$  reach {4}  
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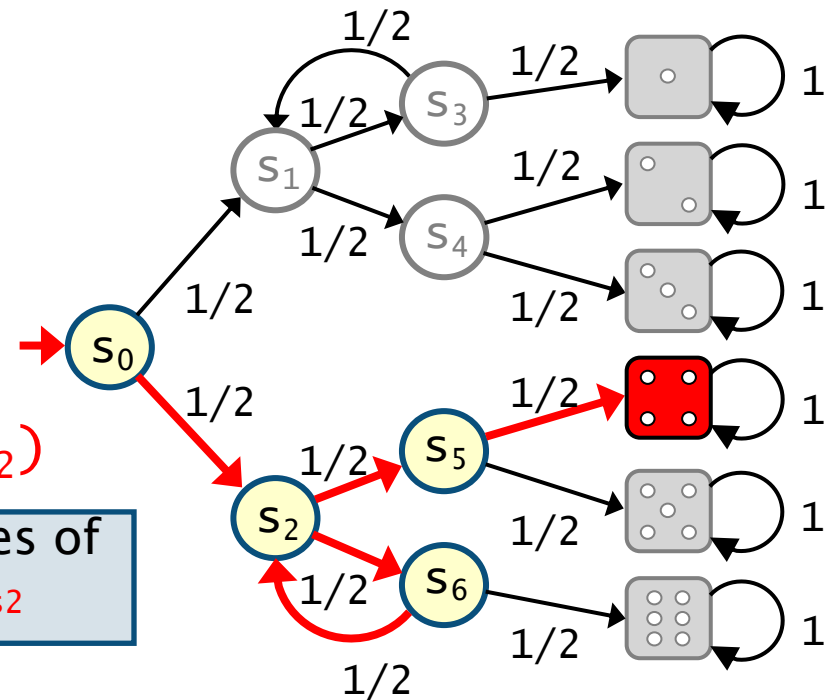
- $x_{s_0} = 1/2 \cdot x_{s_2}$

- $x_{s_2} = (1/2) \cdot (1/2) + (1/2) \cdot (1/2 \cdot x_{s_2})$

- $x_{s_5} = 1/2$

- $x_{s_6} = 1/2 \cdot x_{s_2}$

substituting the values of  $x_{s_5}$  and  $x_{s_6}$  into  $x_{s_2}$



# Probabilistic reachability – Example

## Probability of reaching {4}

– i.e. tossing a four

- $x_4 = 1$
- only  $s_0, s_2, s_5$  and  $s_6$  reach {4}  
therefore  $x_s = 0$  for all other states
- $x_{s_0} = 1/2 \cdot x_{s_2}$
- $(3/4) \cdot x_{s_2} = 1/4$
- $x_{s_5} = 1/2$
- $x_{s_6} = 1/2 \cdot x_{s_2}$

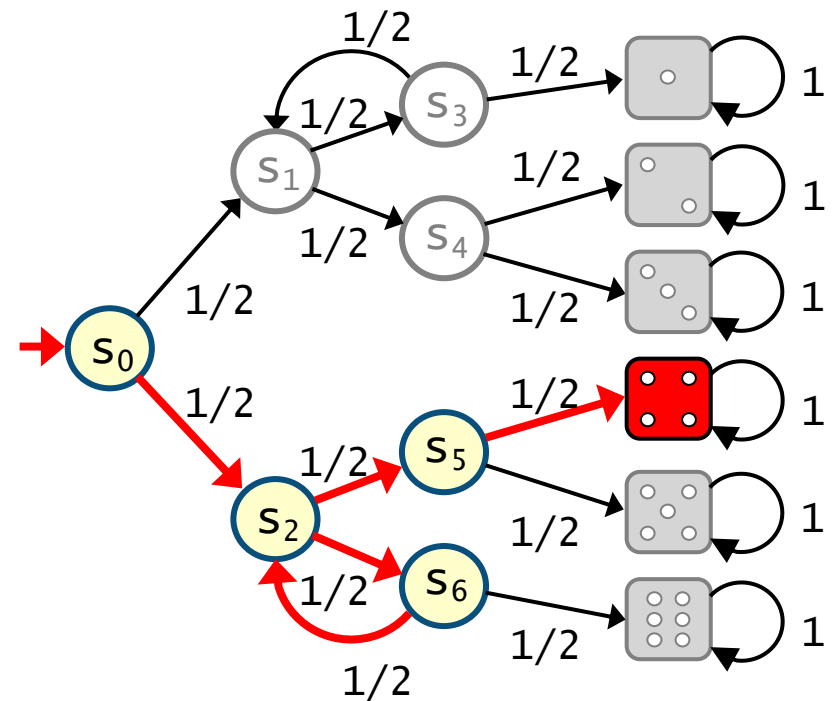
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simplifying



# Probabilistic reachability – Example

## Probability of reaching {4}

– i.e. tossing a four

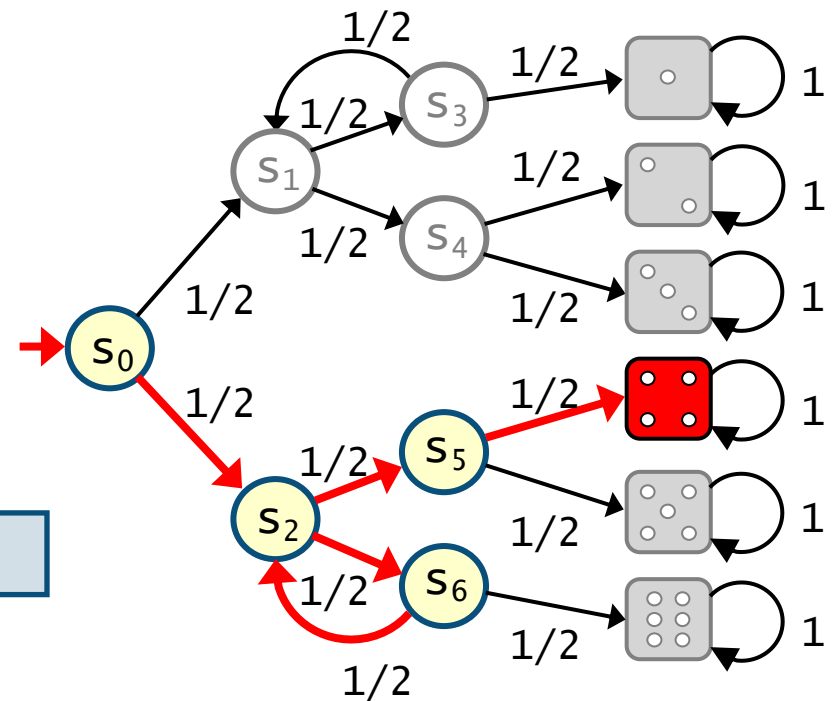
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- $x_{s_0} = 1/2 \cdot x_{s_2}$
- $x_{s_2} = (1/4) / (3/4) = 1/3$
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# Probabilistic reachability – Example

## Probability of reaching {4}

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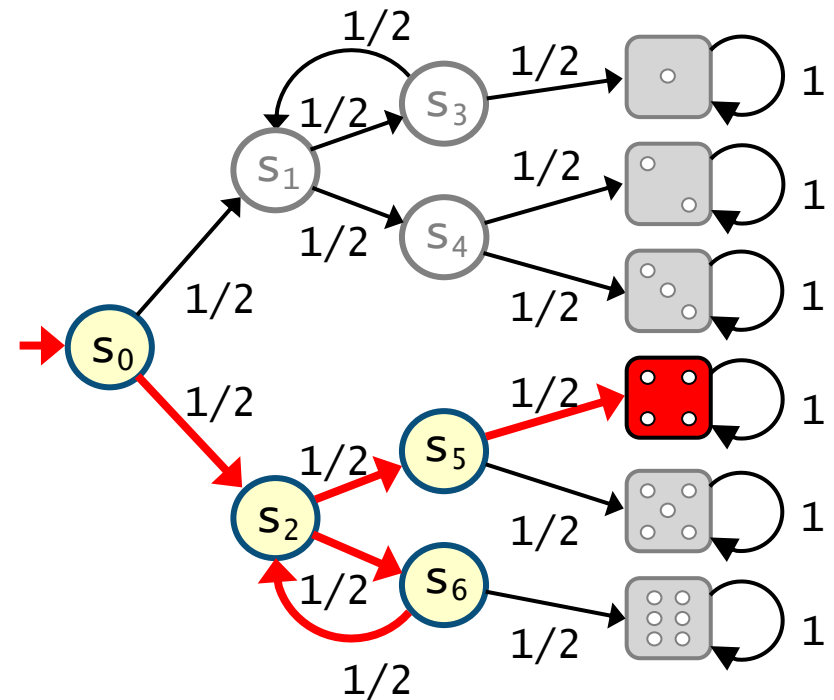
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- $x_4 = 1$
- only  $s_0, s_2, s_5$  and  $s_6$  reach {4}  
therefore  $x_s = 0$  for all other states
- $x_{s_0} = 1/2 \cdot x_{s_2} = (1/2) \cdot (1/3) = 1/6$
- $x_{s_2} = 1/3$
- $x_{s_5} = 1/2$
- $x_{s_6} = 1/2 \cdot x_{s_2} = (1/2) \cdot (1/3) = 1/6$



substituting the value of  $x_{s_2}$  into the other equations

# Probabilistic reachability – Example

## Probability of reaching {4}

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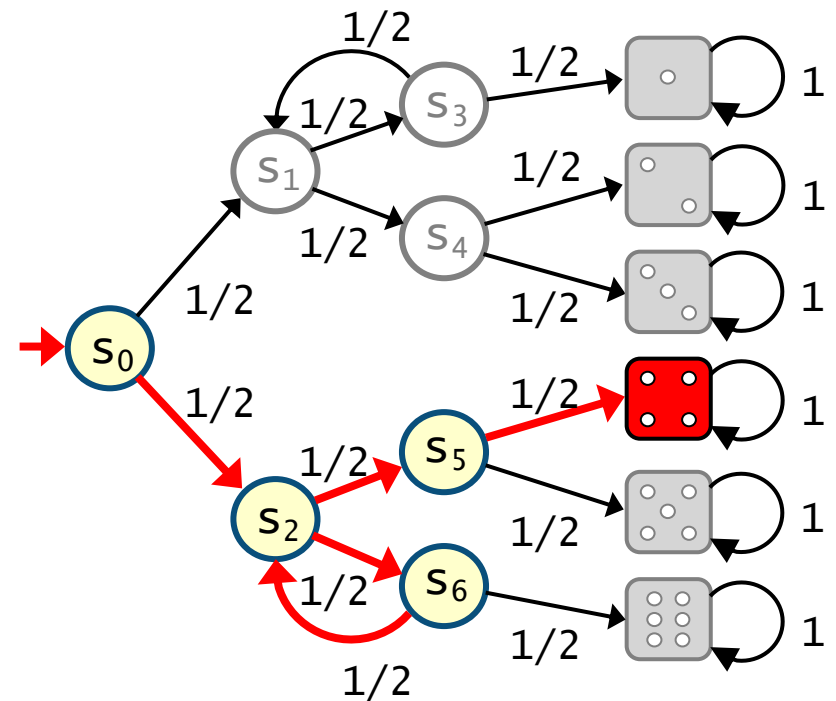
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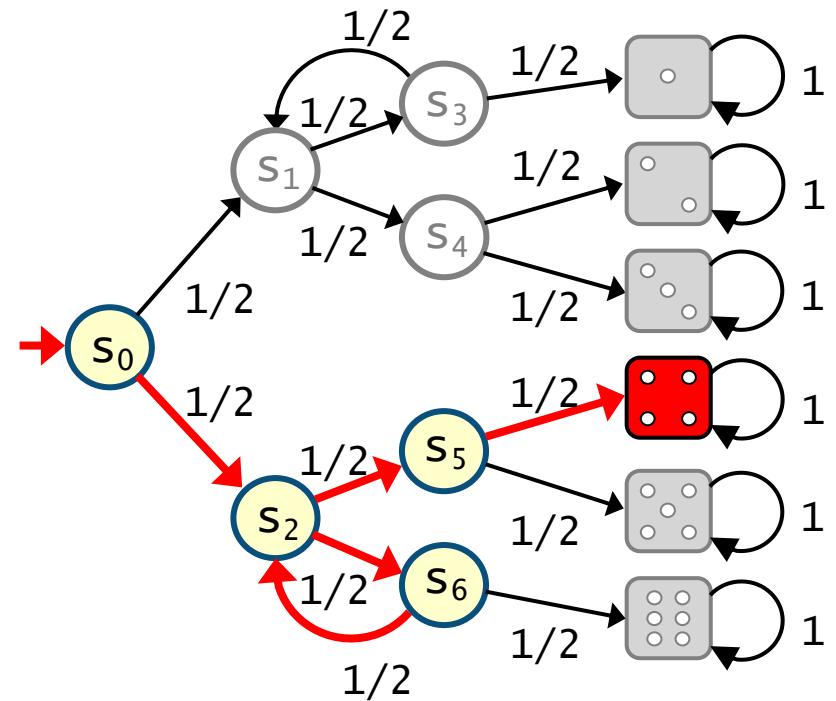
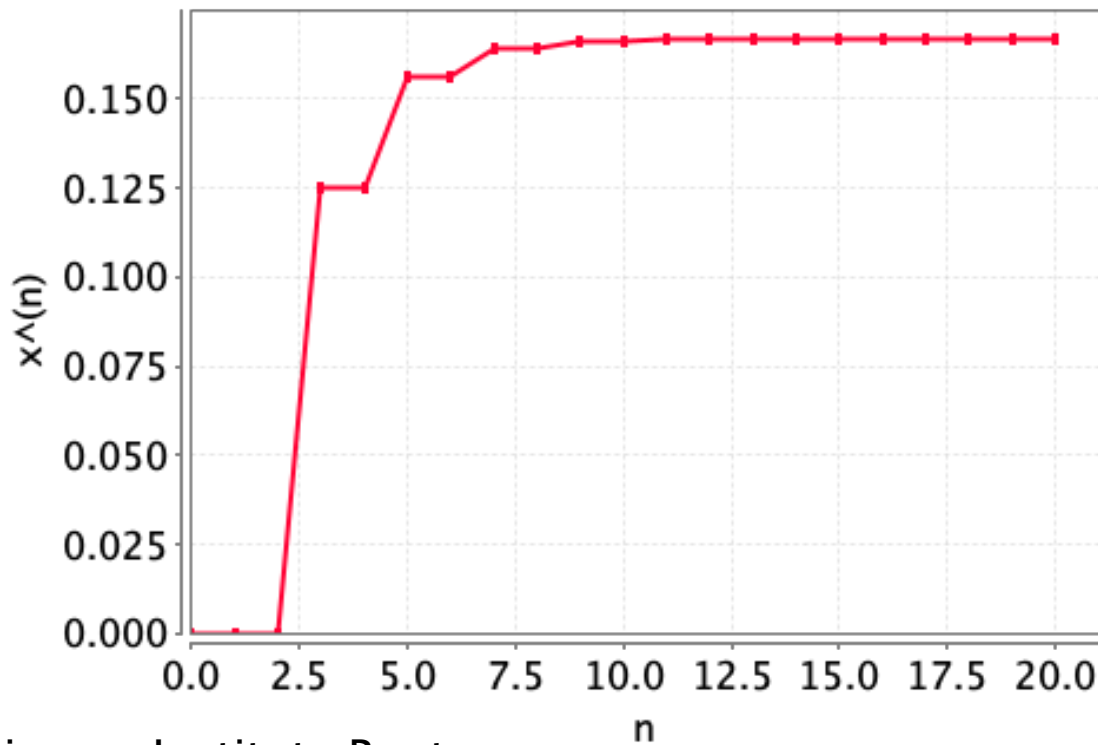


# Probabilistic reachability – Example

Probability of reaching **{4}**, as a least fixpoint

–  $x^{(0)}=0$  and  $x^{(n+1)}=F(x^{(n)})$  where

$$F(y)(s) = \begin{cases} 1 & \text{if } s \in T \\ \sum_{s' \in S} P(s, a)(s') \cdot y(s') & \text{otherwise} \end{cases}$$





# Probabilistic reachability – Complexity

---

## Computing reachability probabilities for DTMCs reduces to

- graph-based analysis (finding the states that can reach the target)
- solving a linear equation system

## Graph based analysis

- linear in the size of the DTMC (simple backwards traversal to find the states that can reach the target)

## Solving a system of linear equations

- polynomial (cubic) in the size of the DTMC (Gaussian elimination)

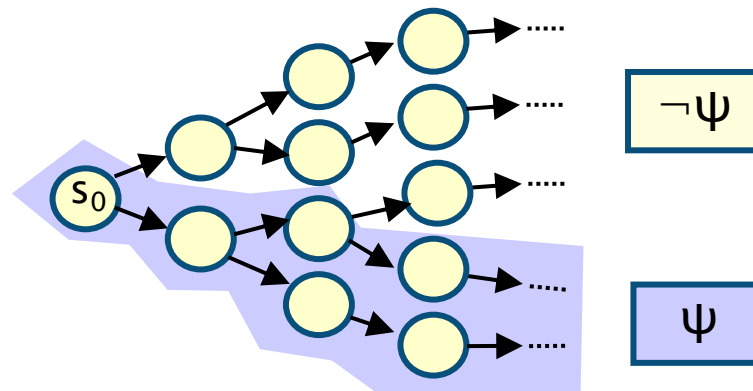
## In practice iterative methods are used for solving large linear equation systems

- power method (i.e. as a least fixed point), Gauss–Siedel

# More general probabilistic properties

For example can compute the **probability** an LTL formula  $\psi$  is true

- $P_{=?}[\psi]$  “what is the probability that  $\psi$  holds?”
- need to compute **Prob** {  $\omega \in \text{Path} \mid \omega \text{ satisfies } \psi$  }



Such sets of path are measurable (elements of the event set)

- therefore probability is well defined

# LTL model checking for DTMCs

---

Model check LTL specification  $\psi$  against a DTMC

1. **Generate a deterministic Rabin automaton (DRA) for  $\psi$** 
  - build nondeterministic Büchi automaton (NBA) for  $\psi$  [VW94]
  - convert the NBA to a DRA [Saf88]
2. **Construct product DTMC  $D \otimes A$**
3. **Identify accepting BSCCs of  $D \otimes A$** 
  - BSCC: bottom strongly connected components
    - these are sets of states such that any state can be reached from any other state and once entered one cannot leave the set
4. **Compute probability of reaching accepting BSCCs**
  - from all states of the  $D \otimes A$

# Discrete time Markov chains

---

Discrete-time Markov chains (DTMCs)

Paths and probabilities for DTMCs

Probabilistic reachability for DTMCs

**Rewards and expected reachability for DTMCs**

# Reward structures

---

We augment DTMCs with **rewards** (or, conversely, **costs**)

- real-valued quantities assigned to states and/or transitions
- these can have a wide range of possible interpretations

**Some examples:**

- elapsed time, power consumption, size of message queue, number of messages successfully delivered, net profit, ...

**Costs? or rewards?**

- mathematically, no distinction between rewards and costs
- when interpreted normally desirable to minimise costs and maximise rewards

# DTMC reward structures

For a DTMC a reward structure is a pair  $(\underline{r}, \mathbf{R})$

- $\underline{r}: S \rightarrow \mathbb{R}_{\geq 0}$  is the **state reward function** (vector over states)
- $\mathbf{R}: S \times S \rightarrow \mathbb{R}_{\geq 0}$  is the **transition reward function** (matrix over states)

$r(s)$  – the reward associated with state  $s$

$R(s, s')$  – the reward associated with the transition from  $s$  to  $s'$

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## Examples

- “**time-steps**”:  $r$  is not used and  $R$  returns  $\mathbf{1}$  for all transitions
- “**number of messages lost**”:  $r$  is not used and  $R$  maps transitions corresponding to a message loss to  $\mathbf{1}$
- “**power consumption**”:  $r$  is defined as the per-time-step energy consumption in each state and  $R$  as the energy cost of each transition

# Probability basics – Expectations

Recall a probability space is a tuple  $(\Omega, \Sigma, \text{Prob})$

- $\Omega$  is the sample space
- $\Sigma$  is the event set
- $\text{Prob}: \Sigma \rightarrow [0, 1]$  is the probability measure

Real valued **random variable  $X$**  over the probability space is a (measurable) function  $X: \Omega \rightarrow \mathbb{R}$

maps elements of the sample space to real values



# Probability basics – Expectations

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Real valued **random variable**  $X$  over the probability space is a (measurable) function  $X:\Omega\rightarrow\mathbb{R}$

Expected (“average”) value of the random variable:

$$\text{Exp}(X) = \int_{\omega\in\Omega} X(\omega) \, d\text{Prob}$$

measurability needed for integral to be well-defined

# Probability basics – Expectations

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Expected (“average”) value of the random variable:

$$\begin{aligned}\text{Exp}(X) &= \int_{\omega\in\Omega} X(\omega) \, d\text{Prob} \\ &= \sum_{\omega\in\Omega} X(\omega) \cdot \text{Prob}(\omega)\end{aligned}$$

if the probability space is discrete, e.g. finite

# Probability basics – Expectations

Example: expected (average) value of a die when tossed

## Probability space

- sample space: possible values  $\{“1”, “2”, “3”, “4”, “5”, “6”\}$
- events: subsets of the sample space
- probability measure:  $\mathbf{Prob} (“1”) = \mathbf{Prob} (“2”) = \dots = \mathbf{Pr} (“6”) = 1/6$

## Random variable $X: \Omega \rightarrow \mathbb{R}$

- the value of the die:  $X (“1”) = 1, X (“2”) = 2, \dots, X (“6”) = 6$

## Expected value of the random variable

- i.e. the expected (average) value of the die when thrown
- $E(X) = \mathbf{Prob} (“1”) \cdot X (“1”) + \mathbf{Prob} (“2”) \cdot X (“2”) + \dots + \mathbf{Prob} (“6”) \cdot X (“6”)$   
 $= 1/6 \cdot 1 + 1/6 \cdot 2 + \dots + 1/6 \cdot 6 = 21/6 = 3\frac{1}{2}$

# Expected reward properties

## Probability space for DTMCs

- sample space is the set infinite paths  $\text{Path}$
- therefore random variables of the form  $X:\text{Path}\rightarrow\mathbb{R}$

function from  
infinite paths to  
real values

Consider any infinite path  $\omega=s_0s_1s_2\dots$

## Cumulative (reachability)

- “reward cumulated before reaching a target set  $T$ ”
- random variable  $X$  where

$$X(\omega) \text{ equals } r(s_0) + \dots + r(s_{k-1}) + R(s_0, s_1) + \dots + R(s_{k-1}, s_k) \\ \text{if } k = \min\{ j \mid s_j \in T \} \text{ exists}$$

find the first time that a state in  $T$  is reached along the path

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summation of rewards up until  $T$  is reached for the first time

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- sample space is the set infinite paths  $\text{Path}$
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function from  
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Consider any infinite path  $\omega = s_0 s_1 s_2 \dots$

## Cumulative (reachability)

- “reward cumulated before reaching a target set  $T$ ”
- random variable  $X$  where

$X(\omega)$  equals  $r(s_0) + \dots + r(s_{k-1}) + R(s_0, s_1) + \dots + R(s_{k-1}, s_k)$   
if  $k = \min\{ j \mid s_j \in T \}$  exists and **infinity** otherwise

summation of rewards up until  $T$  is reached for the first time  
however, if  $T$  is never reached the cumulated reward is infinity

# Computing the expected rewards

---

Computing expected cumulated reward before reaching a state in  $T$

- graph-based analysis and solving a system of linear equations
- compute the expectations for all states

Let  $y_s$  denote the value of  $E(X)$  when starting from state  $s$

# Computing the expected rewards

## Computing expected cumulated reward before reaching a state in $T$

- graph-based analysis and solving a system of linear equations
- compute the expectations for all states

Let  $y_s$  denote the value of  $E(X)$  when starting from state  $s$

- if  $s$  is in  $T$ , then  $y_s = 0$

we have reached a state in  $T$  so no rewards to cumulate



# Computing the expected rewards

## Computing expected cumulated reward before reaching a state in $T$

- graph-based analysis and solving a system of linear equations
- compute the expectations for all states

Let  $y_s$  denote the value of  $E(X)$  when starting from state  $s$

- if  $s$  is in  $T$ , then  $y_s = 0$
- if  $s$  does not reach  $T$  with probability  $1$ , then  $y_s = \infty$

follows from the fact that if no state in  $T$  is reached we set the cumulated reward to infinity for the path

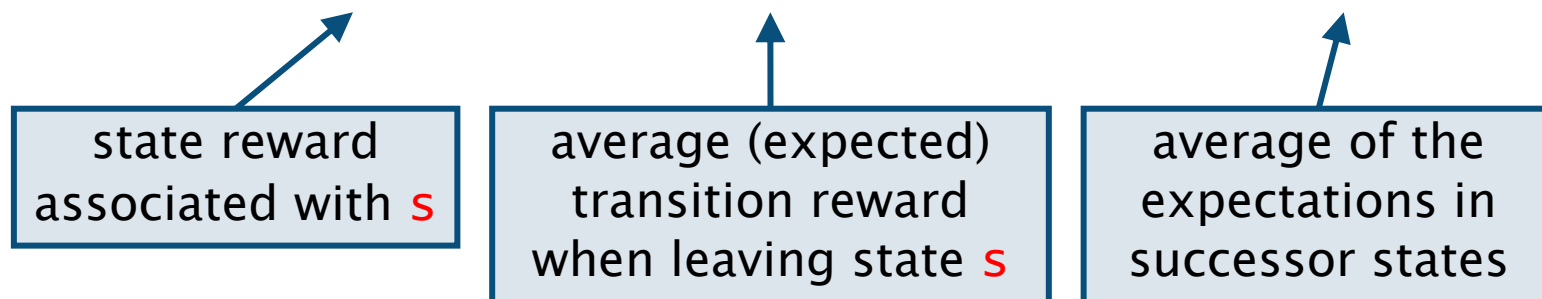
# Computing the expected rewards

## Computing expected cumulated reward before reaching a state in $T$

- graph-based analysis and solving a system of linear equations
- compute the expectations for all states

Let  $y_s$  denote the value of  $E(X)$  when starting from state  $s$

- if  $s$  is in  $T$ , then  $y_s = 0$
- if  $s$  does not reach a state in  $T$  with probability 1, then  $y_s = \infty$
- otherwise  $y_s = r(s) + \sum_{s' \in S} P(s, s') \cdot R(s, s') + \sum_{s' \in S} P(s, s') \cdot y_{s'}$



# Expected rewards – Example

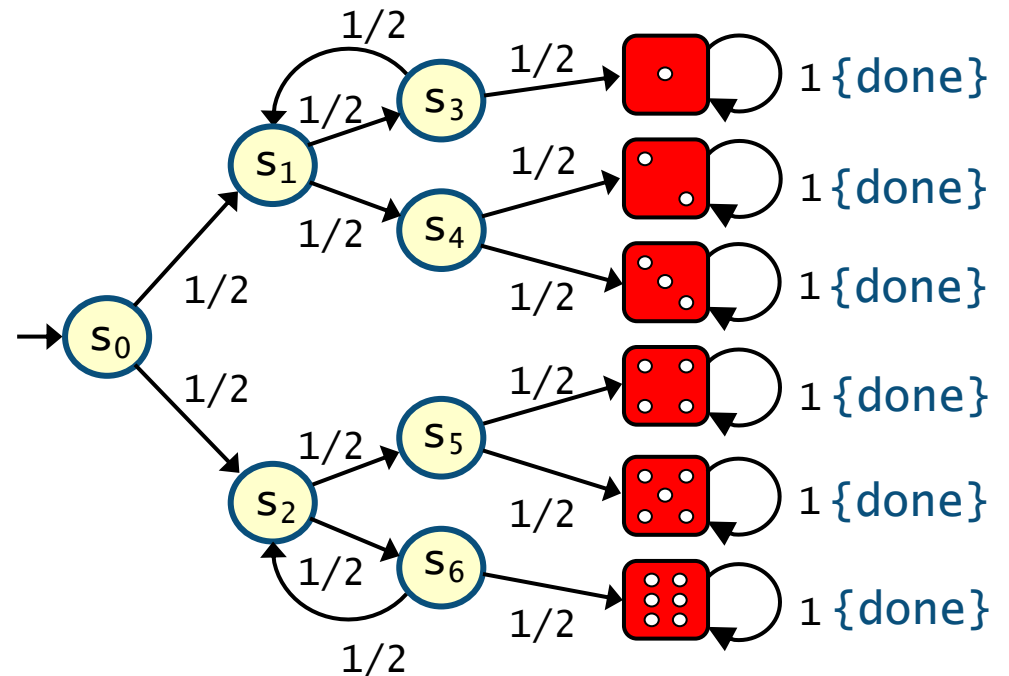
Expected  
number of  
coin flips

if  $s$  is in  $T$ , then  $y_s = 0$

if  $s$  does not reach  $T$  with probability 1, then  $y_s = \infty$

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- reward structure  $R(s, s')=1$  for all states labelled **done**

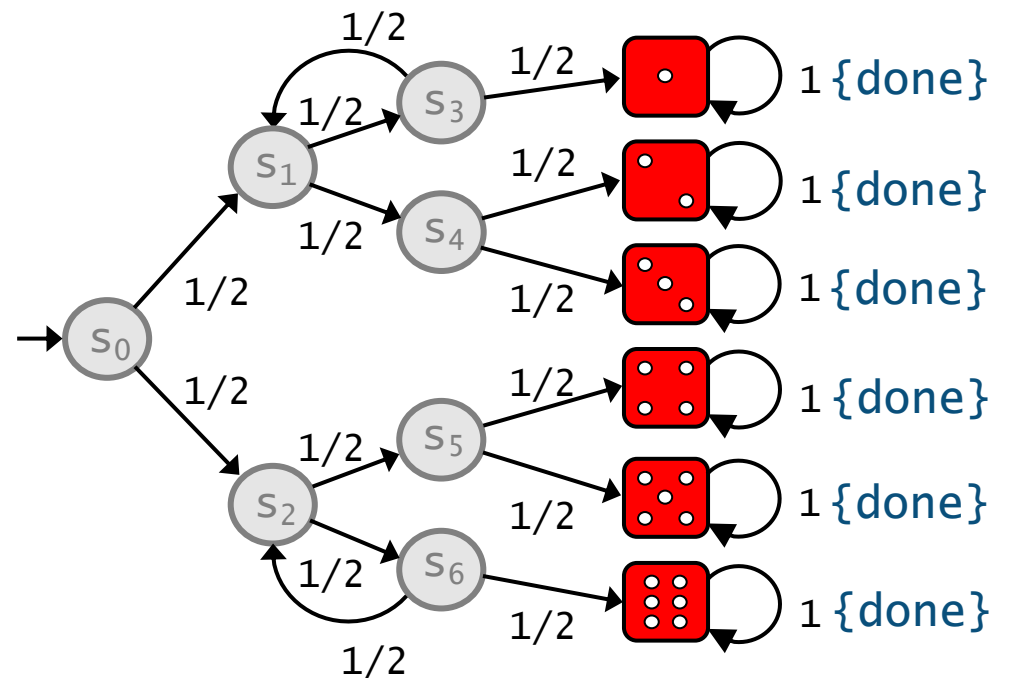


# Expected rewards – Example

Expected number of coin flips

if  $s$  is in  $T$ , then  $y_s = 0$   
if  $s$  does not reach  $T$  with probability 1, then  $y_s = \infty$   
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- reward structure  $R(s, s')=1$  for all states labelled **done**
- $y_s = 0$  for all states labelled **done**



# Expected rewards – Example

## Expected number of coin flips

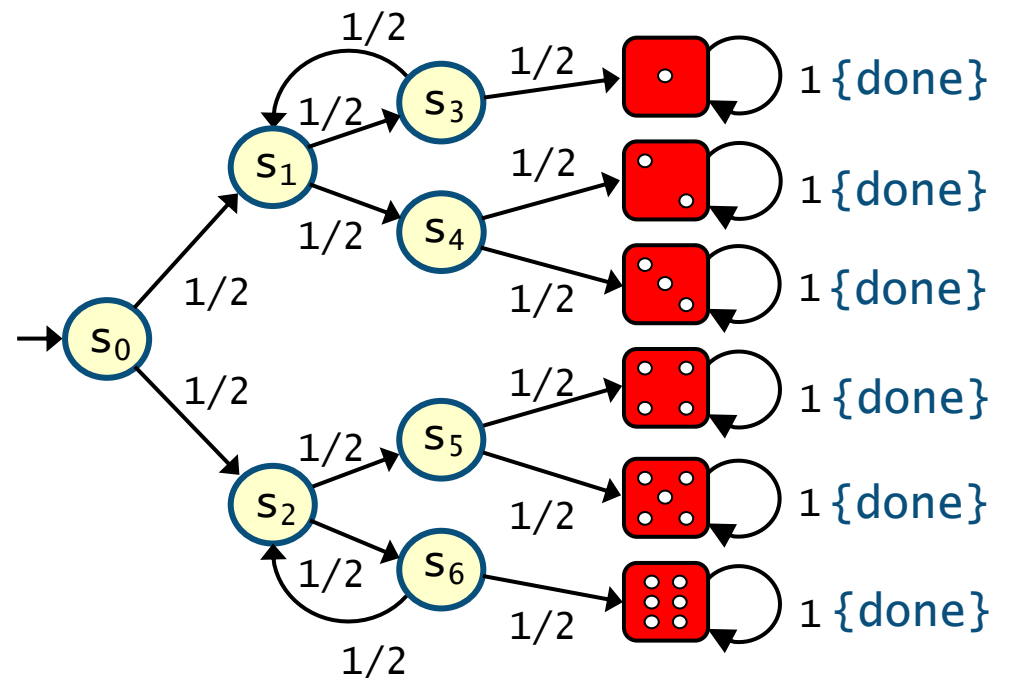
if  $s$  is in  $T$ , then  $y_s = 0$

if  $s$  does not reach  $T$  with probability 1, then  $y_s = \infty$

otherwise  $y_s = r(s) + \sum_{s' \in S} P(s, s') \cdot R(s, s') + \sum_{s' \in S} P(s, s') \cdot y_{s'}$

– reward structure  $R(s, s')=1$  for all states labelled **done**

- $y_s = 0$  for all states labelled **done**
- all states reach done with probability 1 therefore no state has value infinity



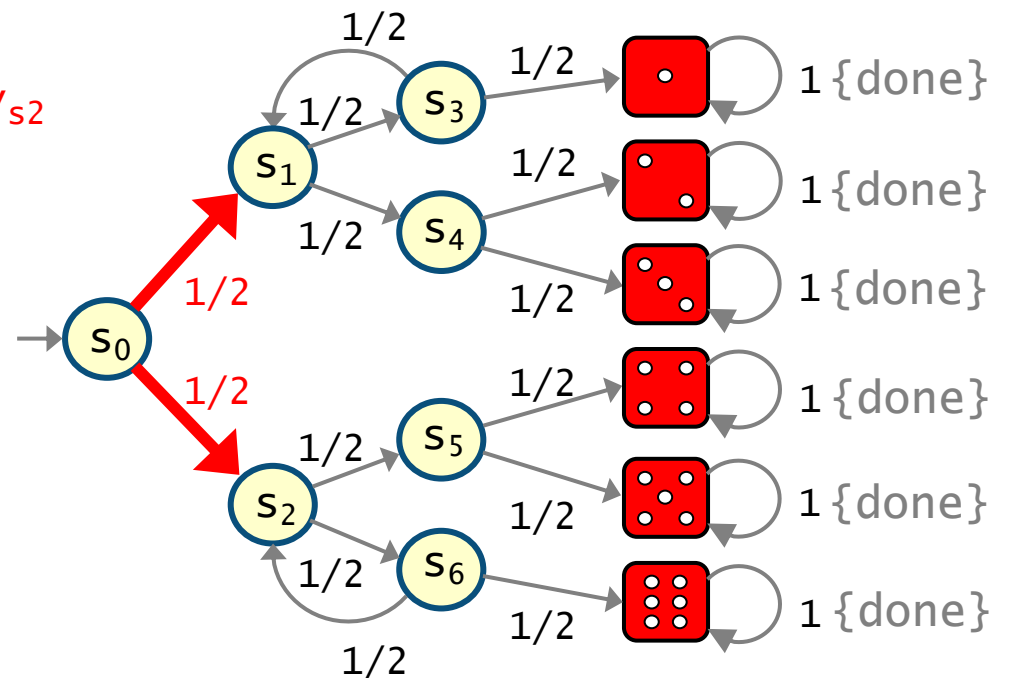
# Expected rewards – Example

## Expected number of coin flips

if  $s$  is in  $T$ , then  $y_s = 0$   
if  $s$  does not reach  $T$  with probability 1, then  $y_s = \infty$   
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– reward structure  $R(s, s')=1$  for all states labelled **done**

- $y_s = 0$  for all states labelled **done**
- $y_{s_0} = 1/2 \cdot 1 + 1/2 \cdot 1 + 1/2 \cdot y_{s_1} + 1/2 \cdot y_{s_2}$



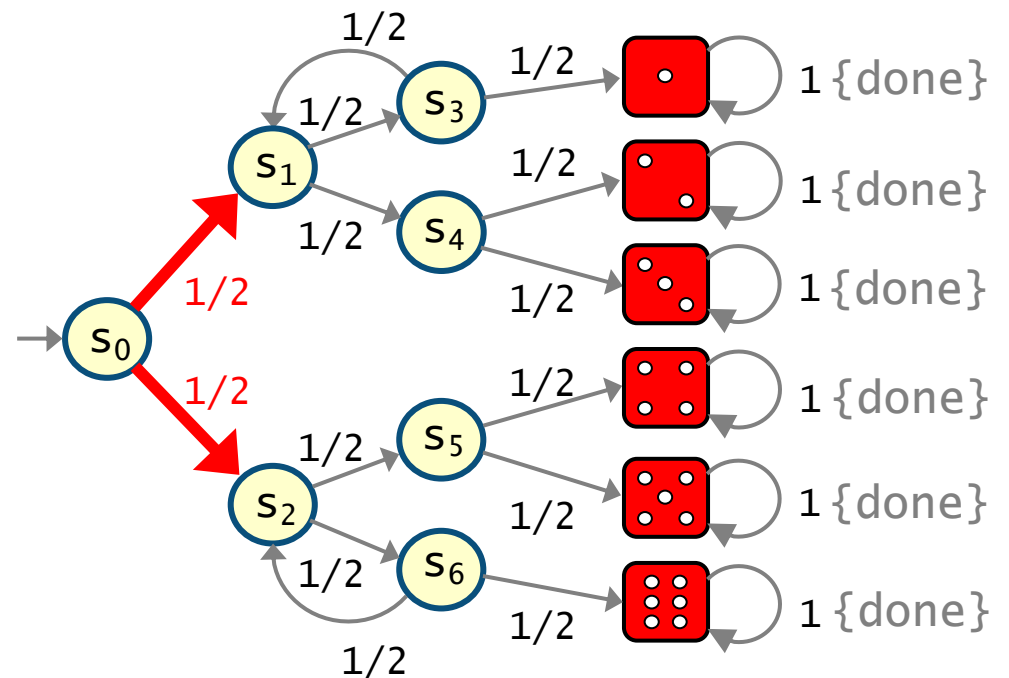
# Expected rewards – Example

## Expected number of coin flips

if  $s$  is in  $T$ , then  $y_s = 0$   
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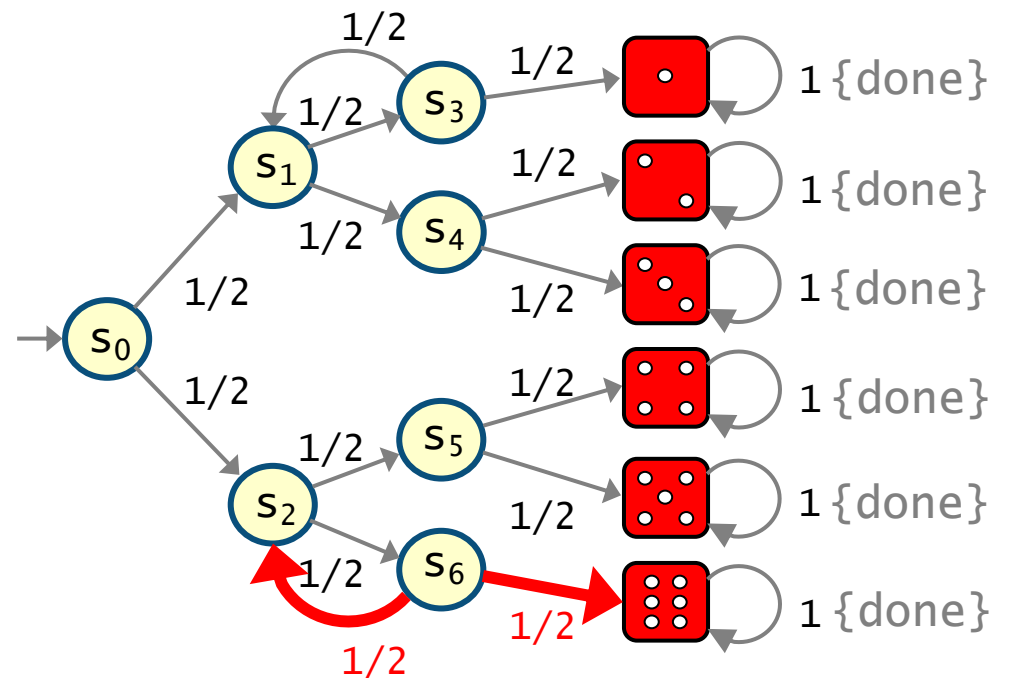
# Expected rewards – Example

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if  $s$  is in  $T$ , then  $y_s = 0$   
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– reward structure  $R(s, s')=1$  for all states labelled **done**

- $y_s = 0$  for all states labelled **done**
- $y_{s_0} = 1 + 1/2 \cdot y_{s_1} + 1/2 \cdot y_{s_2}$
- $y_{s_1} = 1 + 1/2 \cdot y_{s_3} + 1/2 \cdot y_{s_4}$
- $y_{s_2} = 1 + 1/2 \cdot y_{s_5} + 1/2 \cdot y_{s_6}$
- $y_{s_3} = 1 + 1/2 \cdot y_{s_1} + 1/2 \cdot 0$
- $y_{s_4} = 1 + 1/2 \cdot 0 + 1/2 \cdot 0$
- $y_{s_5} = 1 + 1/2 \cdot 0 + 1/2 \cdot 0$
- $y_{s_6} = 1/2 + 1/2 + 1/2 \cdot 0 + 1/2 \cdot y_{s_2}$





# Expected rewards – Example

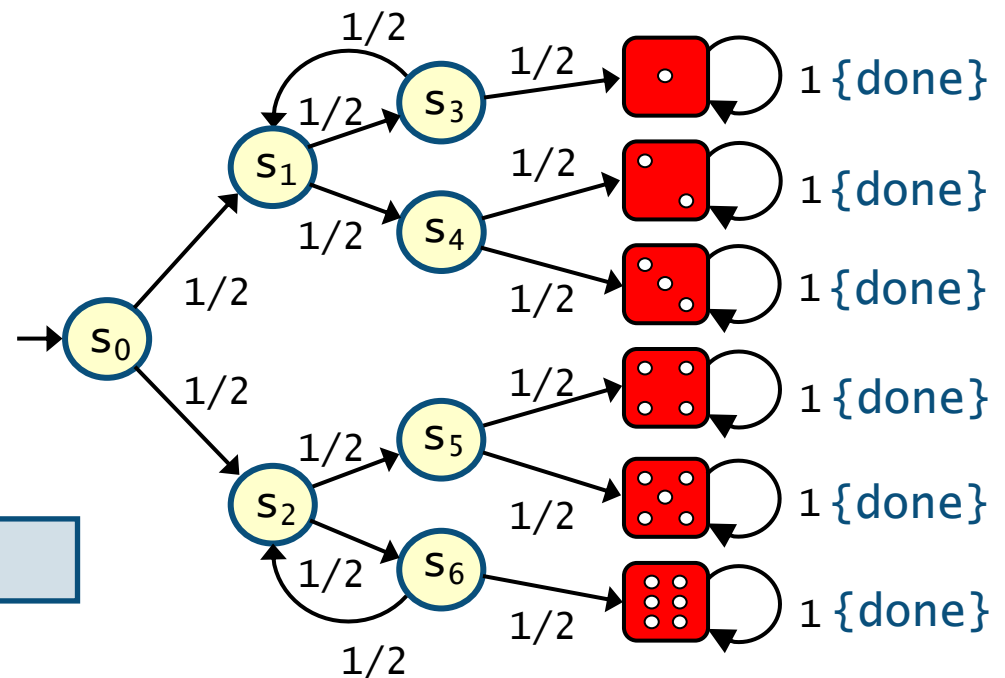
Expected number of coin flips

if  $s$  is in  $T$ , then  $y_s = 0$   
 if  $s$  does not reach  $T$  with probability  $1$ , then  $y_s = \infty$   
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– reward structure  $R(s, s')=1$  for all states labelled **done**

- $y_s = 0$  for all states labelled **done**
- $y_{s_0} = 1 + 1/2 \cdot y_{s_1} + 1/2 \cdot y_{s_2}$
- $y_{s_1} = 1 + 1/2 \cdot y_{s_3} + 1/2 \cdot y_{s_4}$
- $y_{s_2} = 1 + 1/2 \cdot y_{s_5} + 1/2 \cdot y_{s_6}$
- $y_{s_3} = 1 + 1/2 \cdot y_{s_1}$
- $y_{s_4} = 1$
- $y_{s_5} = 1$
- $y_{s_6} = 1 + 1/2 \cdot y_{s_2}$

simplifying



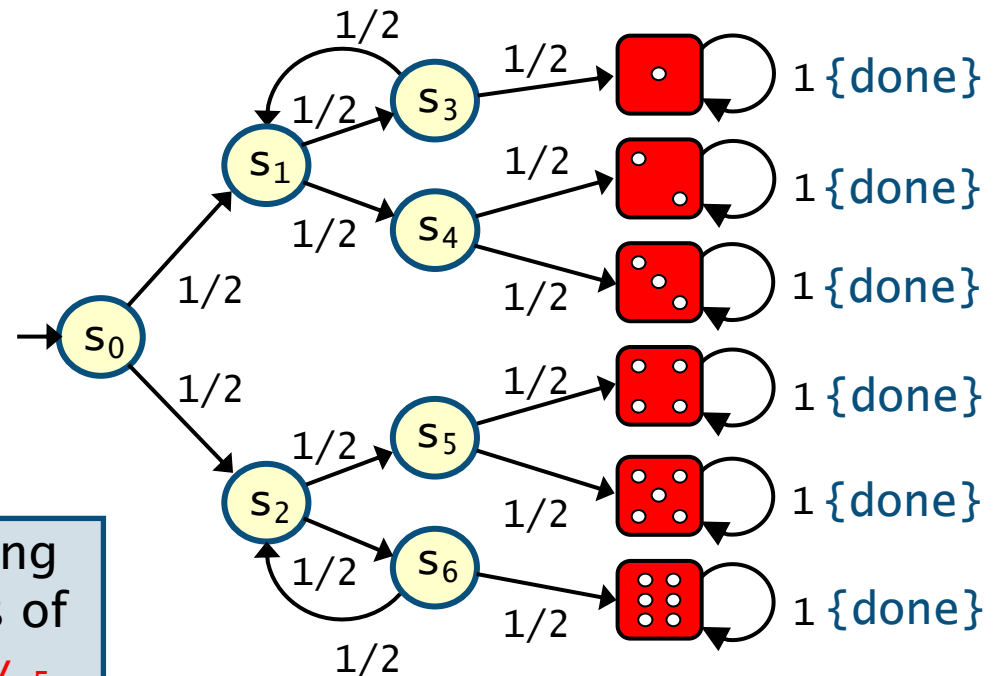
# Expected rewards – Example

## Expected number of coin flips

if  $s$  is in  $T$ , then  $y_s = 0$   
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– reward structure  $R(s, s')=1$  for all states labelled **done**

- $y_s = 0$  for all states labelled **done**
- $y_{s_0} = 1 + 1/2 \cdot y_{s_1} + 1/2 \cdot y_{s_2}$
- $y_{s_1} = 1 + 1/2 \cdot (1 + 1/2 \cdot y_{s_1}) + 1/2$
- $y_{s_2} = 1 + 1/2 + 1/2 \cdot (1 + 1/2 \cdot y_{s_2})$
- $y_{s_3} = 1 + 1/2 \cdot y_{s_1}$
- $y_{s_4} = 1$
- $y_{s_5} = 1$
- $y_{s_6} = 1 + 1/2 \cdot y_{s_2}$



substituting  
the values of  
 $y_{s_3}, y_{s_4}, y_{s_5}$   
and  $y_{s_6}$

# Expected rewards – Example

## Expected number of coin flips

if  $s$  is in  $T$ , then  $y_s = 0$   
 if  $s$  does not reach  $T$  with probability  $1$ , then  $y_s = \infty$   
 otherwise  $y_s = r(s) + \sum_{s' \in S} P(s, s') \cdot R(s, s') + \sum_{s' \in S} P(s, s') \cdot y_{s'}$

– reward structure  $R(s, s')=1$  for all states labelled **done**

•  $y_s = 0$  for all states labelled **done**

•  $y_{s_0} = 1 + 1/2 \cdot y_{s_1} + 1/2 \cdot y_{s_2}$

•  $y_{s_1} = 8/3$

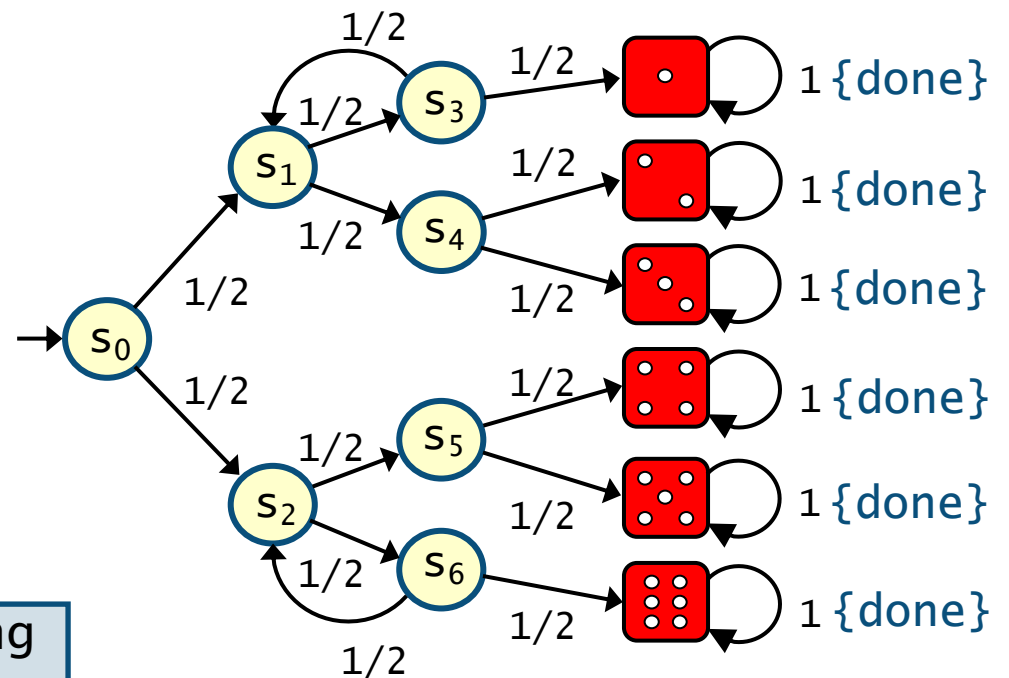
•  $y_{s_2} = 8/3$

•  $y_{s_3} = 1 + 1/2 \cdot y_{s_1}$

•  $y_{s_4} = 1$

•  $y_{s_5} = 1$

•  $y_{s_6} = 1 + 1/2 \cdot y_{s_2}$



simplifying

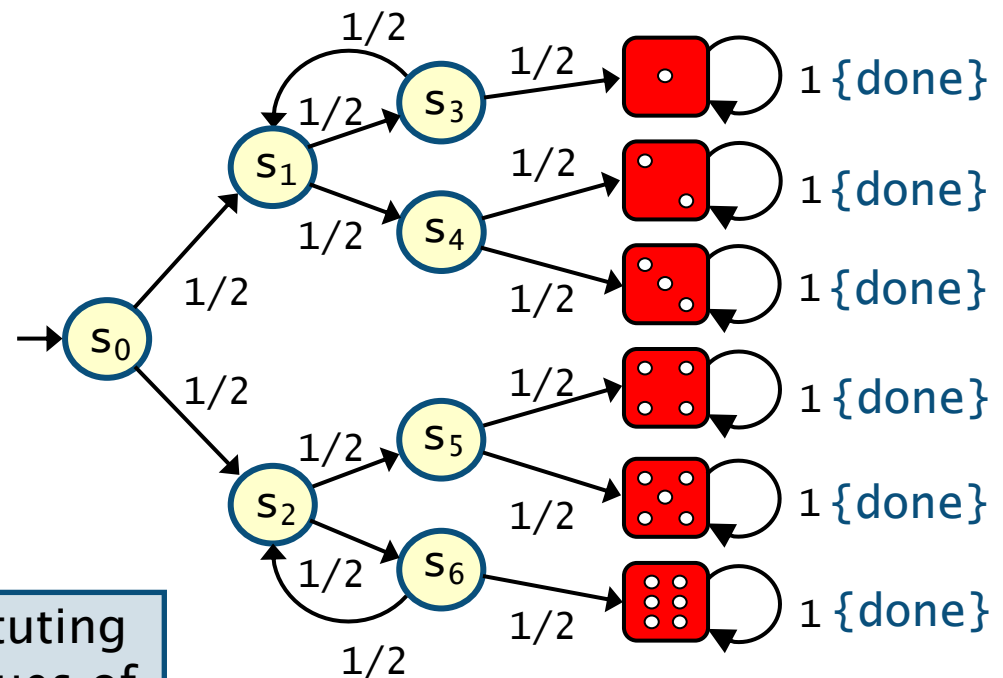
# Expected rewards – Example

## Expected number of coin flips

if  $s$  is in  $T$ , then  $y_s = 0$   
if  $s$  does not reach  $T$  with probability 1, then  $y_s = \infty$   
otherwise  $y_s = r(s) + \sum_{s' \in S} P(s, s') \cdot R(s, s') + \sum_{s' \in S} P(s, s') \cdot y_{s'}$

– reward structure  $R(s, s')=1$  for all states labelled **done**

- $y_s = 0$  for all states labelled **done**
- $y_{s_0} = 1 + 1/2 \cdot (8/3) + 1/2 \cdot (8/3)$
- $y_{s_1} = 8/3$
- $y_{s_2} = 8/3$
- $y_{s_3} = 1 + 1/2 \cdot (8/3)$
- $y_{s_4} = 1$
- $y_{s_5} = 1$
- $y_{s_6} = 1 + 1/2 \cdot (8/3)$



substituting  
the values of  
 $y_{s_1}$  and  $y_{s_2}$

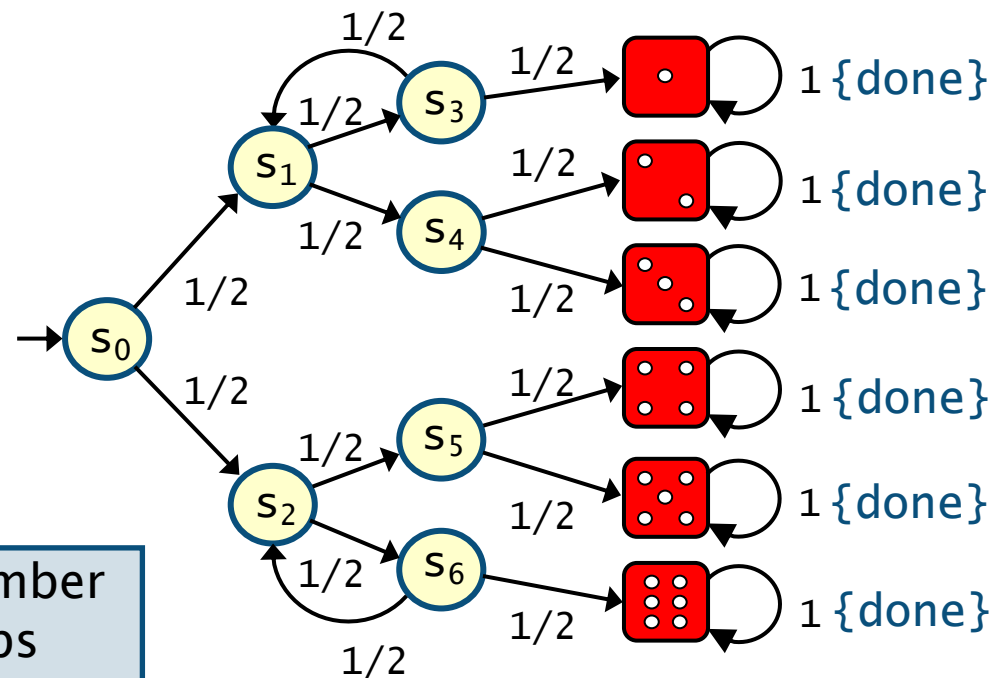
# Expected rewards – Example

## Expected number of coin flips

if  $s$  is in  $T$ , then  $y_s = 0$   
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– reward structure  $R(s, s')=1$  for all states labelled **done**

- $y_s = 0$  for all states labelled **done**
- $y_{s_0} = 11/3$
- $y_{s_1} = 8/3$
- $y_{s_2} = 8/3$
- $y_{s_3} = 7/3$
- $y_{s_4} = 1$
- $y_{s_5} = 1$
- $y_{s_6} = 7/3$



expected number of coin flips equals  $11/3$

# Expected reachability – Complexity

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## Computing expected reachability values for DTMCs reduces to

- graph-based analysis (find states that reach the target with probability 1)
- solving a linear equation system

## Graph based analysis

- linear in the size of the DTMC (simple backwards traversal to find the states that can reach the target)

## Solving a system of linear equations

- polynomial (cubic) in the size of the DTMC (Gaussian elimination)
- again in practice use iterative methods
- as for probabilistic reachability can express as a least fixed point

# Additional Reward Properties

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## Instantaneous

- “the expected value of the state reward at time–step  $k$ ”
- e.g. “the expected queue size after exactly 90 seconds”

## Cumulative (time–bounded)

- “the expected reward cumulated up to time–step  $k$ ”
- e.g. “the expected power consumption over one hour”

Also long run average and multi–objective properties

# In the next video

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## Markov decision processes

- extend DTMCs to allow the modelling of non-deterministic behaviour