

Two questions about randomized in numerical linear algebra

1. ~~Multiway random walks for reducing variance (arxiv.org/1608.04361)~~
2. An eigenvalue law for “triangle Laplacians”
3. An approximation bound for sampling zonotopes for clusters.

David F. Gleich
Purdue University

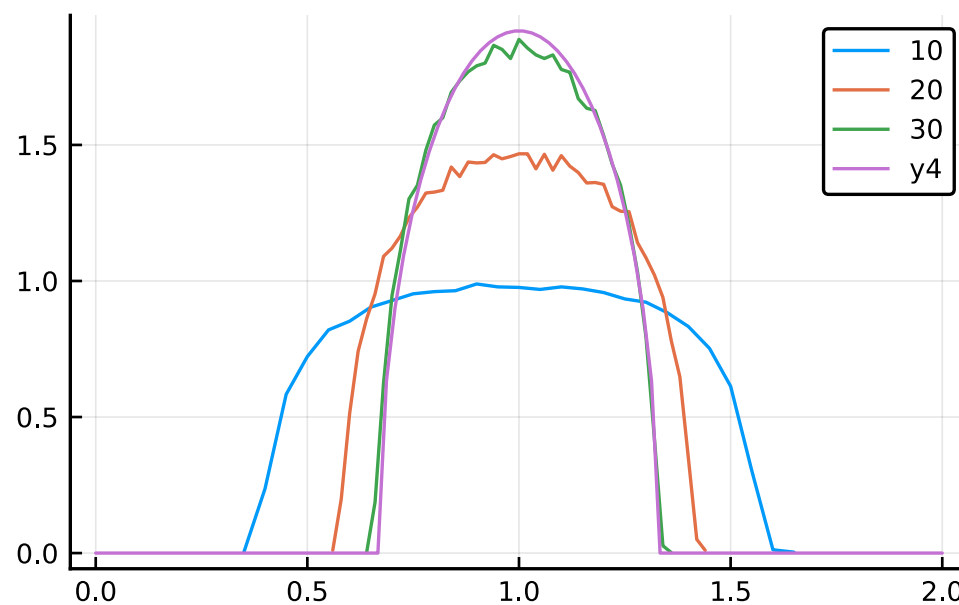
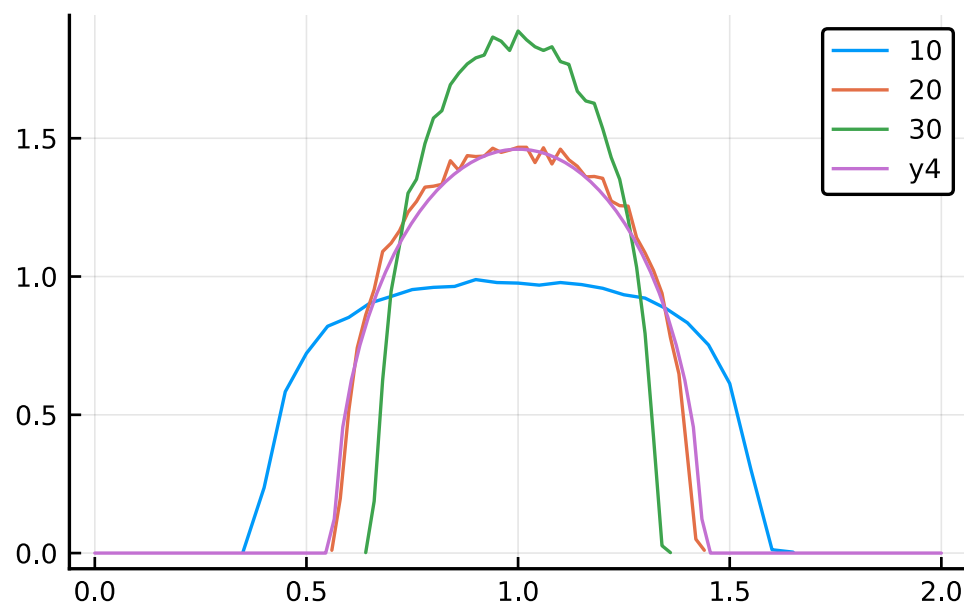
3 slides to pose a simple question.
I don't have an answer 😊

Eigenvalue laws of random ER graphs are well known

The adjacency matrix has a semi-circle law.

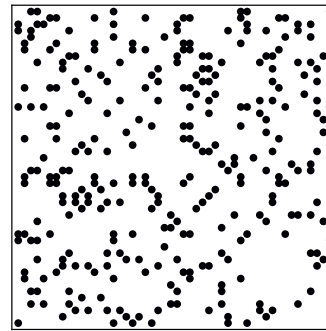
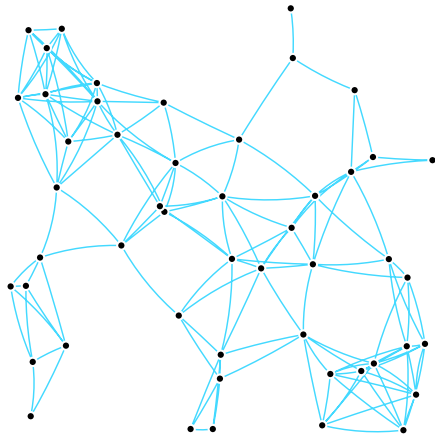
The normalized Laplacian matrix also has a semi-circle law with a specific gap

- C. Hoffman, M. Kahle, and E. Paquette. Spectral gaps of random graphs and applications to random topology. arXiv:1201.0425.



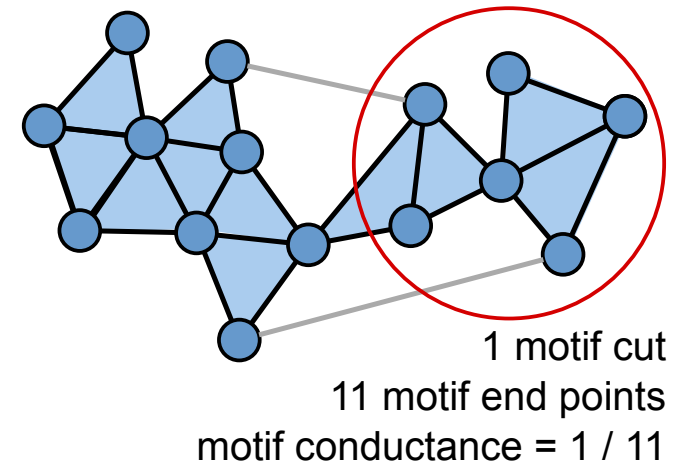
A idea based on our recent work

given a graph $G = (V, E)$ and its adjacency matrix \mathbf{A}



consider using the weighted matrix $\mathbf{W} = \mathbf{A}^2 \odot \mathbf{A}$

The matrix $\mathbf{W} = \mathbf{A}^2 \odot \mathbf{A}$ arises from our **motif and higher-order clustering framework** when using **triangles** as the motif.



The eigenvalues of the normalized Laplacian of this weighted matrix do not follow a Marchenko–Pastur law

Method.

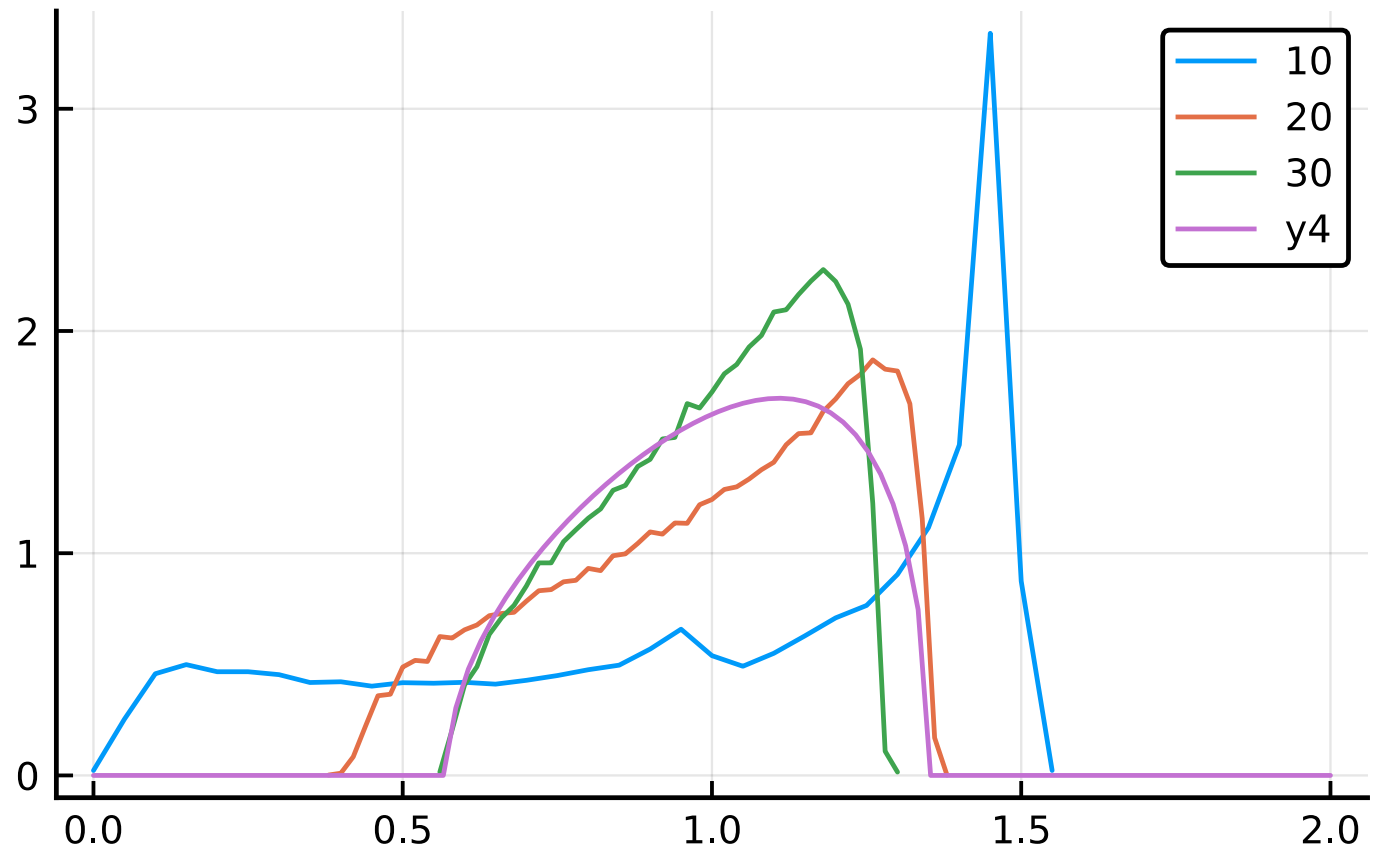
Generate a random ER graph, compute

$$W = A^2 \odot A$$

Compute eigenvalues of the normalized Laplacian of the largest component of W

Our question.

What is this distribution?



Code to reproduce <https://gist.github.com/dgleich/4d4becc858e4a7d7952af6c66c99e7b9>

Another problem where I don't know the solution.

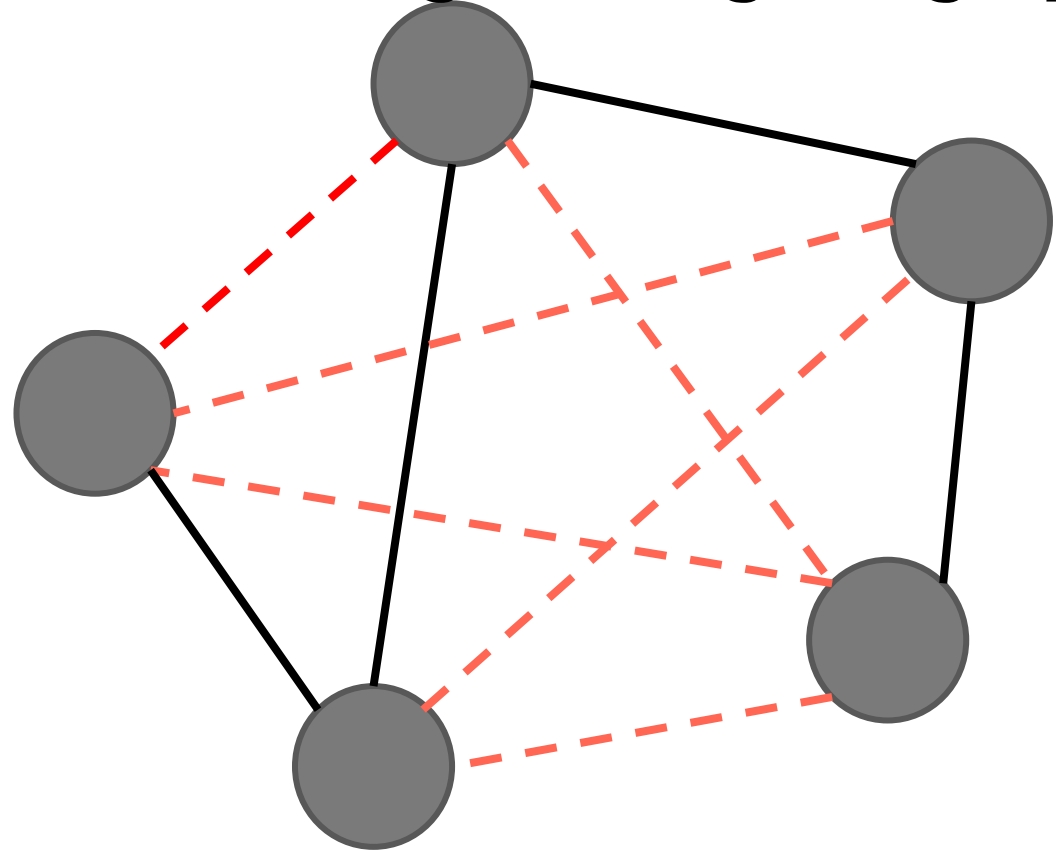
Motivation. Correlation Clustering of Low-Rank Matrices.

Veldt, Wirth, Gleich, 2016. arXiv:1611.07305

We have a heuristic randomized (and useful) algorithm that we don't have good analysis for 😊

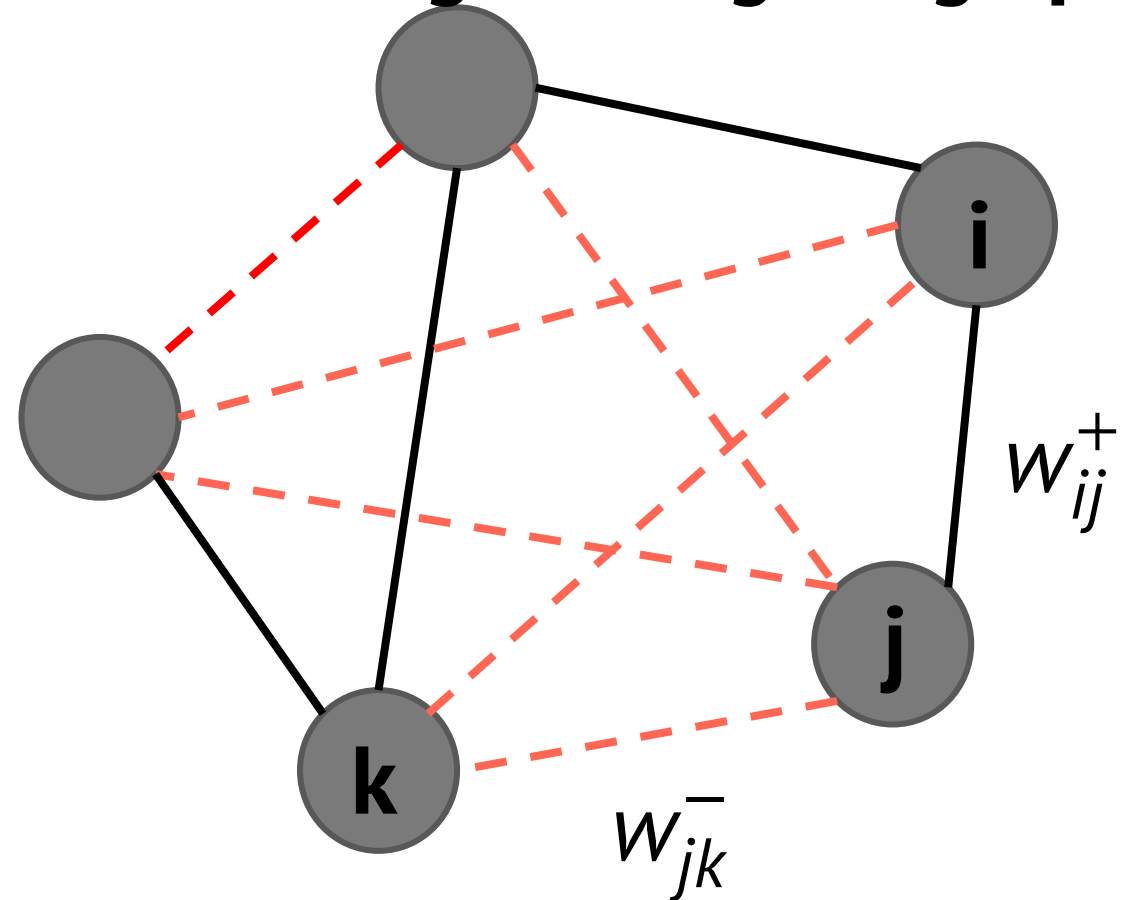
Correlation clustering involves a weighted, signed graph

Edges in a signed graph indicate
similarity (+)
or **dissimilarity** (-)



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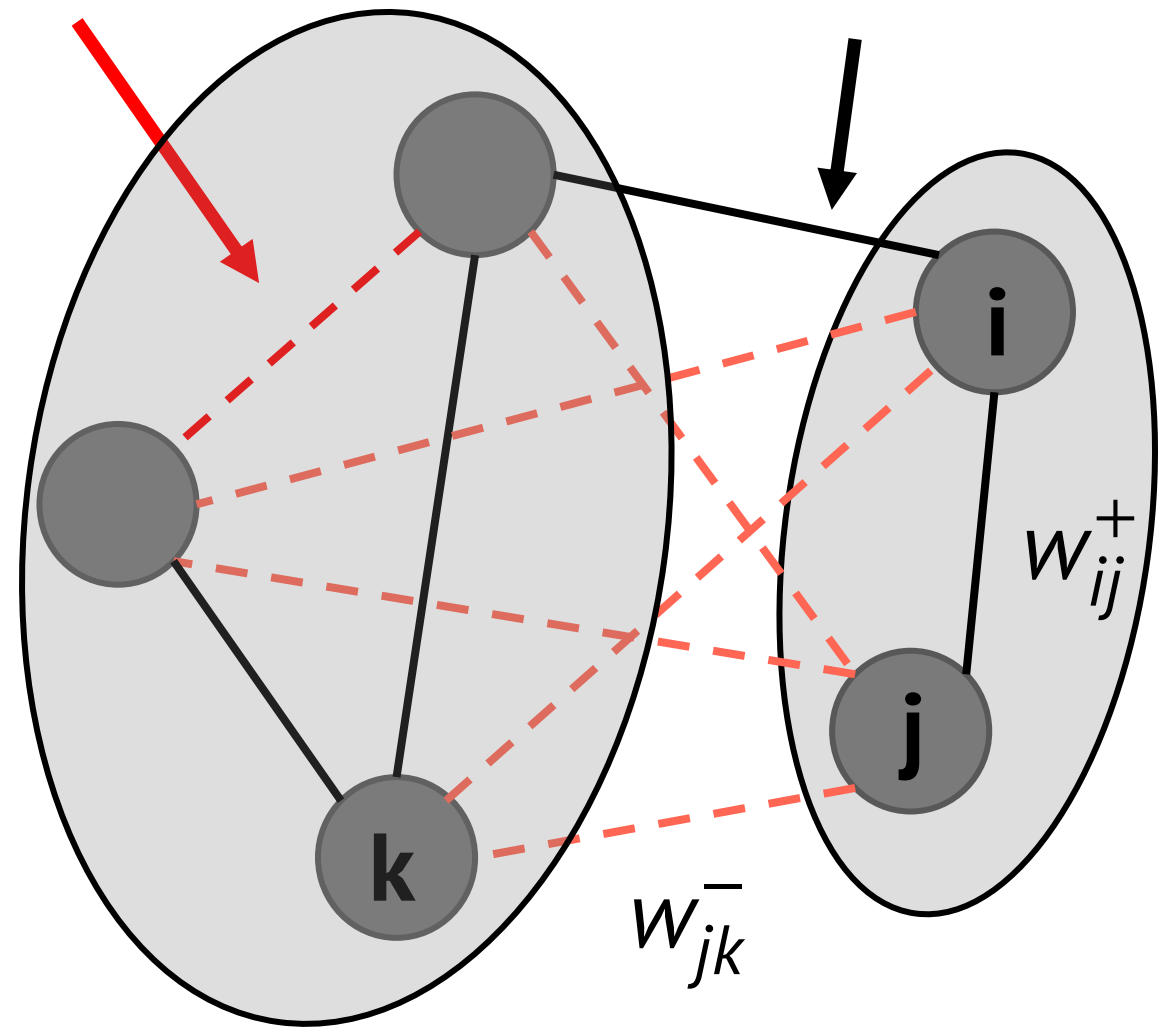


Edges can be weighted, but problems become harder.

Edges in a signed graph indicate **similarity (+)** or **dissimilarity (-)**

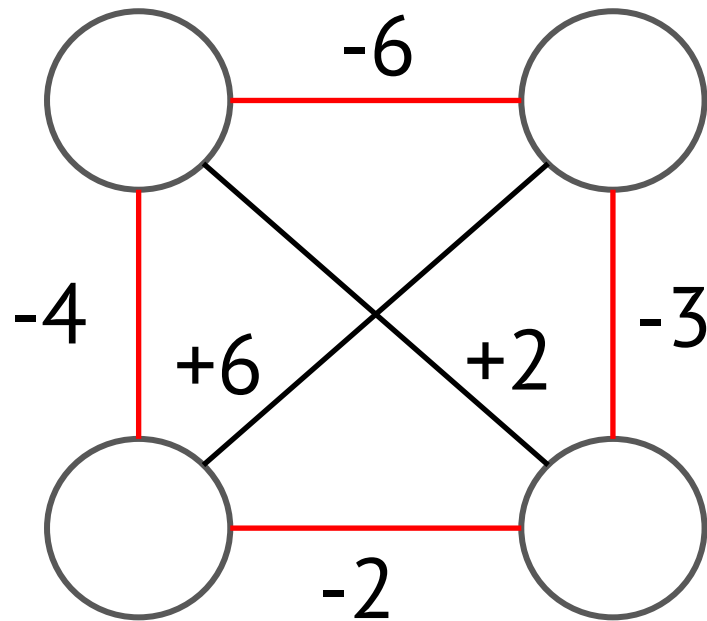
Mistake

Mistake



Objective: Minimize the weight of “mistakes”

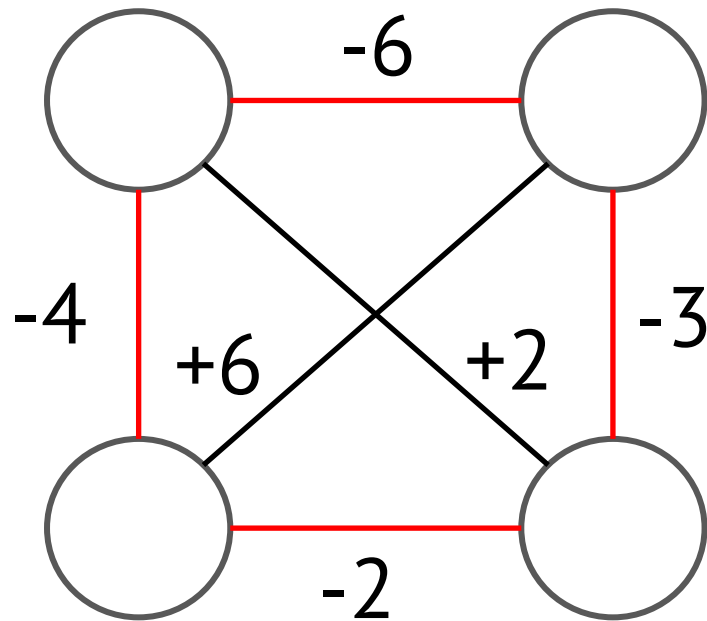
Edge weights can be stored in an adjacency matrix



$$\mathbf{A}_{ij} = w_{ij}^+ - w_{ij}^-$$

$$\mathbf{A} = \begin{bmatrix} 0 & -6 & +2 & -4 \\ -6 & 0 & -3 & +6 \\ +2 & -3 & 0 & -2 \\ -4 & +6 & -2 & 0 \end{bmatrix}$$

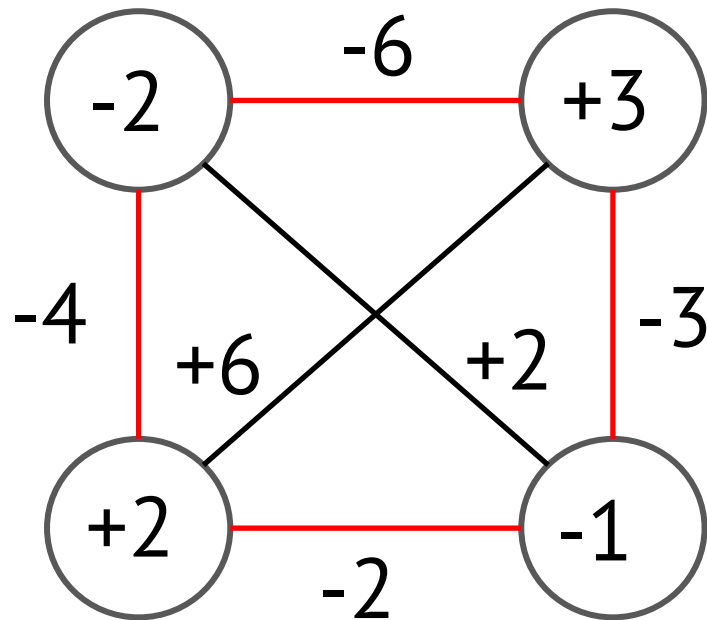
The rank-1 positive semidefinite case is very simple



$$\mathbf{A} = \mathbf{v}\mathbf{v}^T$$

$$\mathbf{A} = \begin{bmatrix} -2 \\ +3 \\ -1 \\ +2 \end{bmatrix} \begin{bmatrix} -2 & +3 & -1 & +2 \end{bmatrix}$$

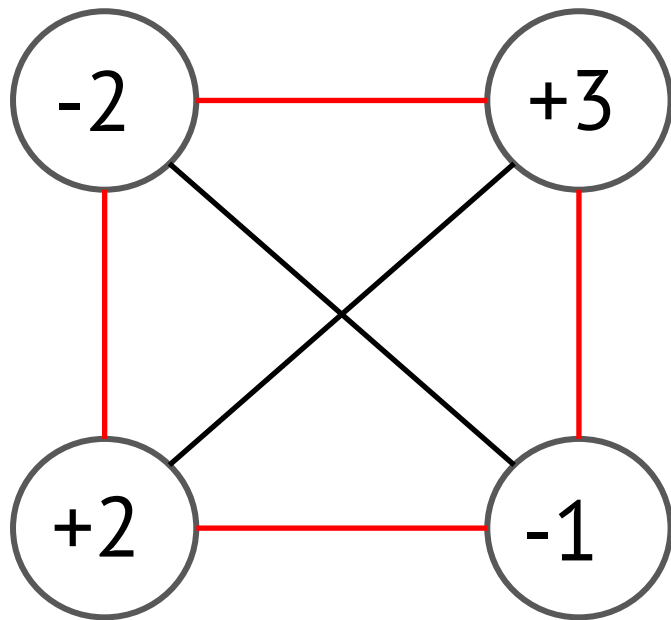
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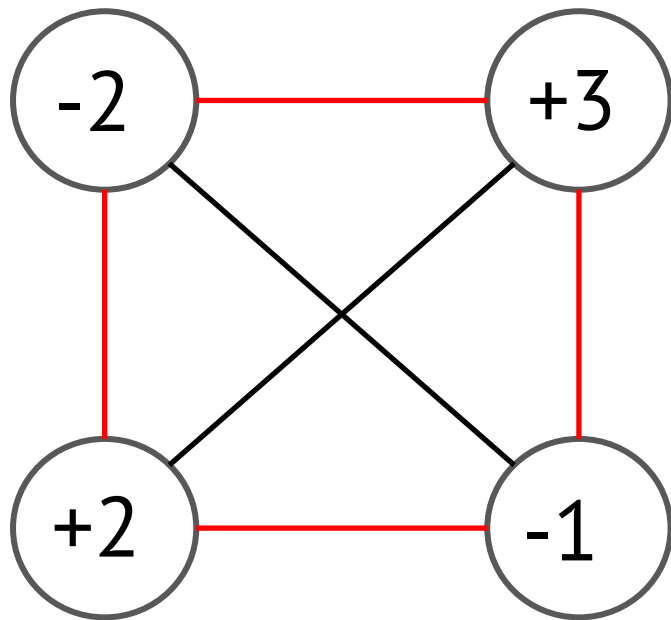
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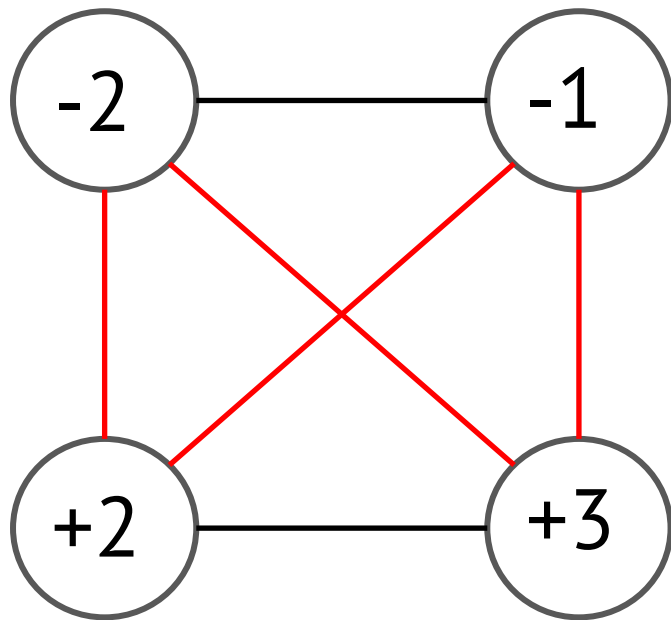


$$\mathbf{A} = \mathbf{v}\mathbf{v}^T$$

$$\begin{bmatrix} -2 \\ +3 \\ -1 \\ +2 \end{bmatrix} \longrightarrow \begin{bmatrix} -2 \\ -1 \\ +2 \\ +3 \end{bmatrix}$$

Ordering \mathbf{v} gives a perfect clustering

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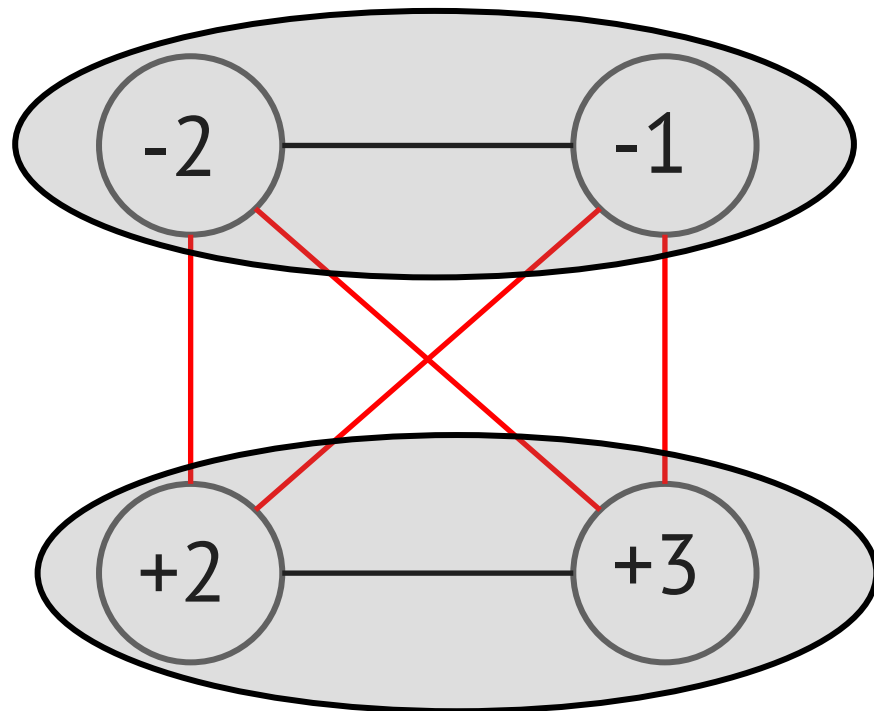


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Ordering \mathbf{v} gives a perfect clustering

What happens for other low-rank matrices?

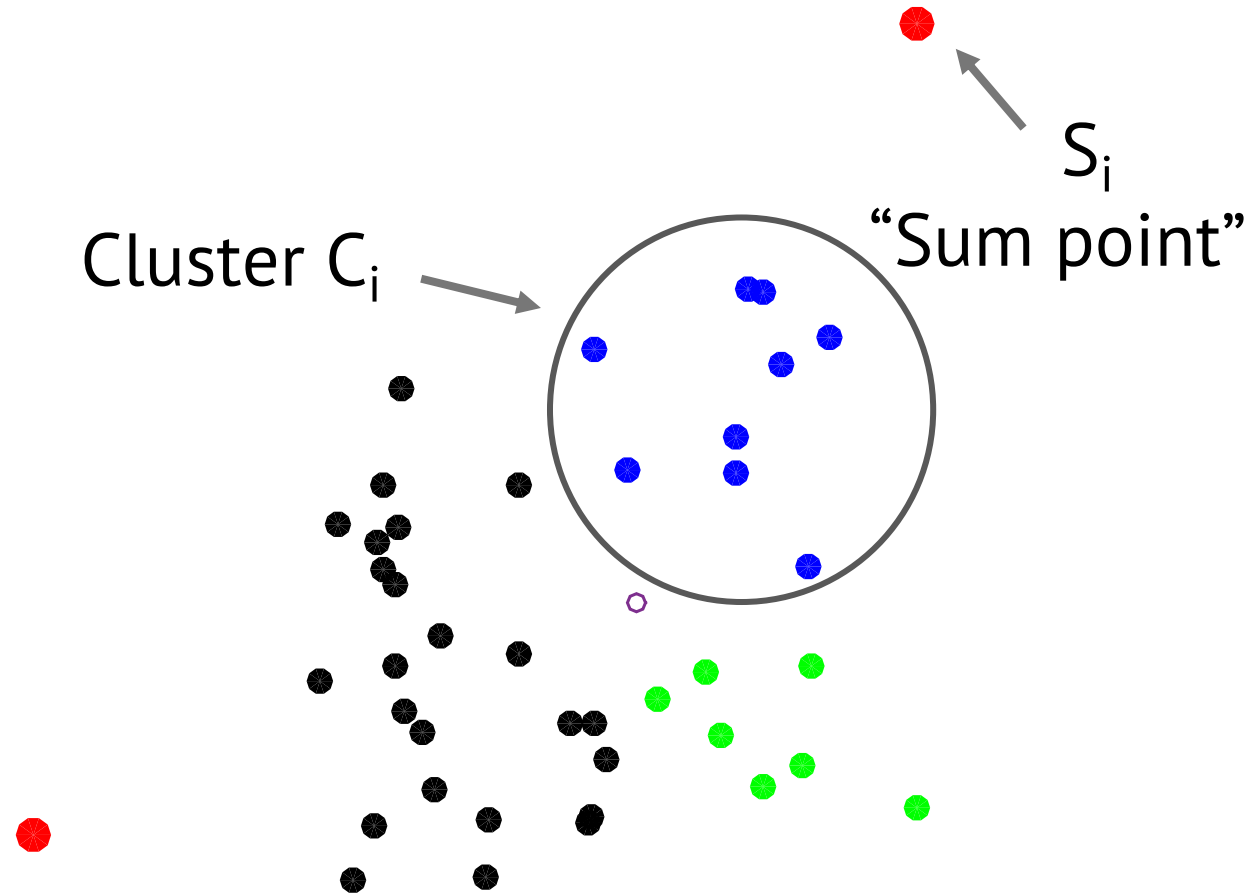
A simple solution for rank-1 positive semidefinite correlation clustering always exists.

Veldt, Gleich, Wirth arXiv:1611.07305

- A single negative eigenvalue makes the problem NP-hard ☹️
- A rank- r pos def matrix can be solved in “polynomial” time $O(n^{r^2-1})$.
Rank 3 = $O(n^8)$ time
- Algorithm is based on equiv between low-rank CC and vector partitioning [Onn & Schulman 2001] based on the vertices of a signing zonotope.
- In Stinson, Gleich, Constantine (arXiv:1602.06620) we proposed a randomized algorithm to find vertices of a zonotope.

Vector partitioning formulation

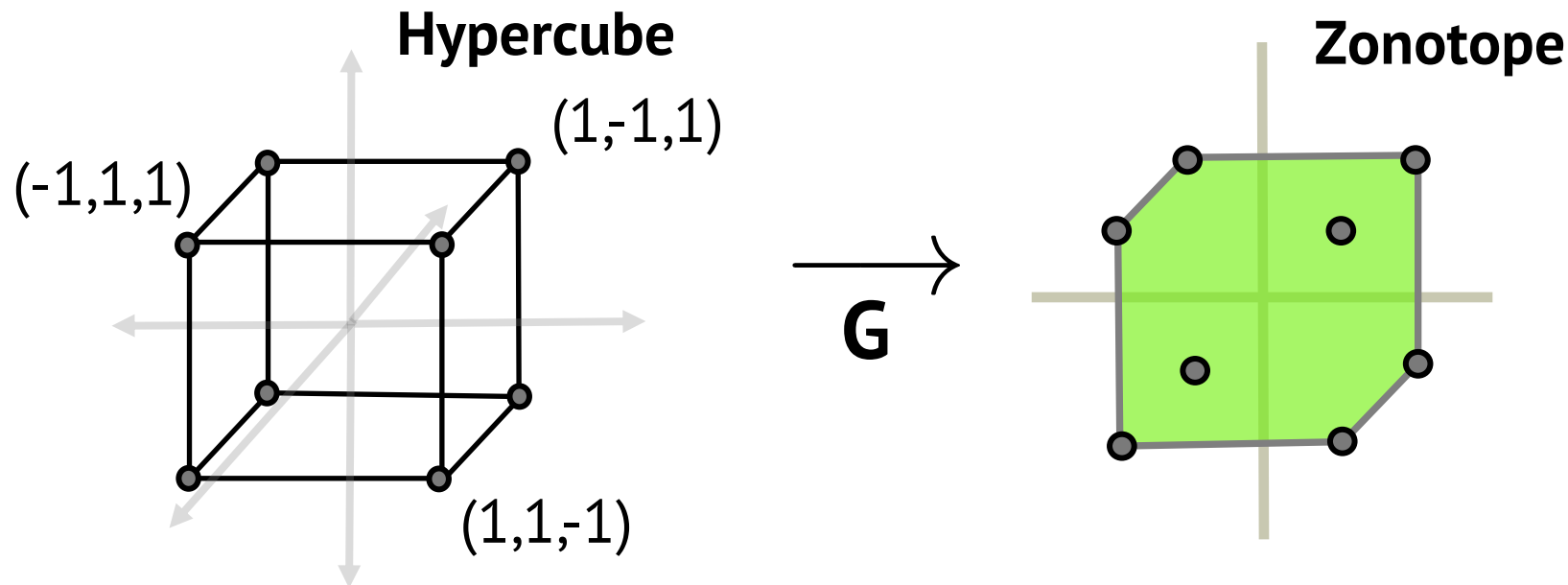
$$\max \sum_{k=1}^{\text{\# clusters}} \|S_k\|^2$$



A zonotope is the linear projection of a hypercube into a lower dimension

$$\mathcal{Z}(\mathbf{G}) = \text{conv}\{\mathbf{G}\mathbf{x} \mid \mathbf{x} \in \{\pm 1\}^n\}$$

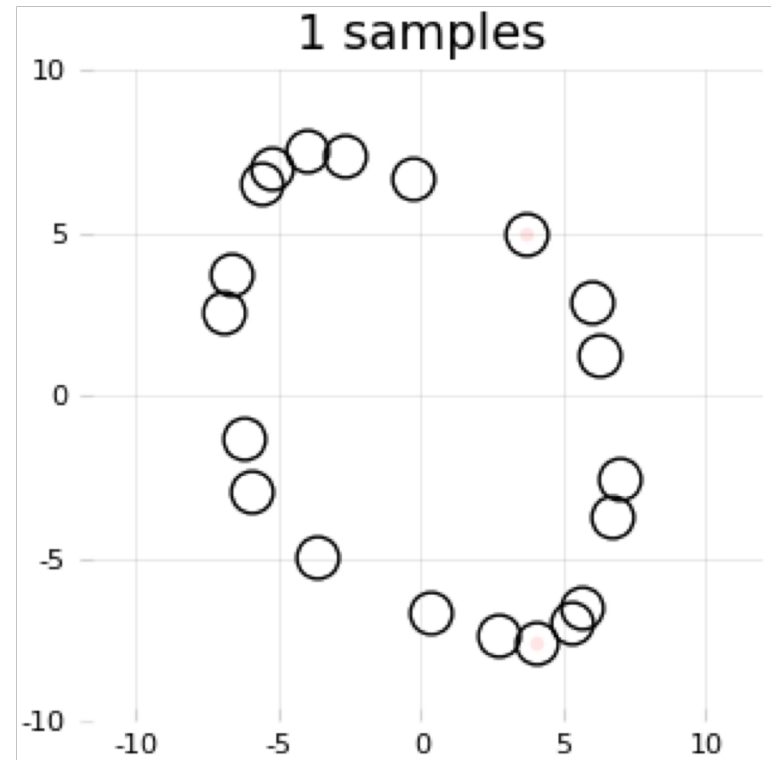
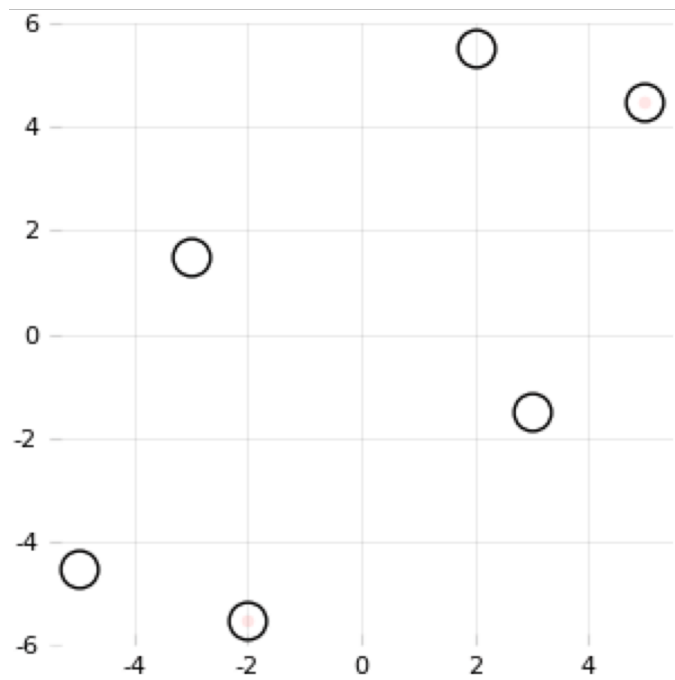
$\mathbf{G} \in \mathbb{R}^{D \times N}$ “generator” matrix



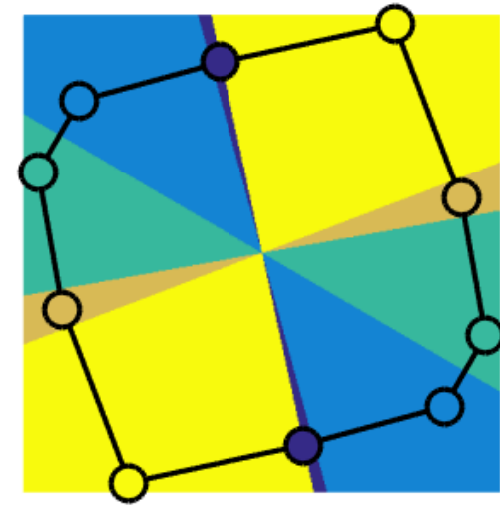
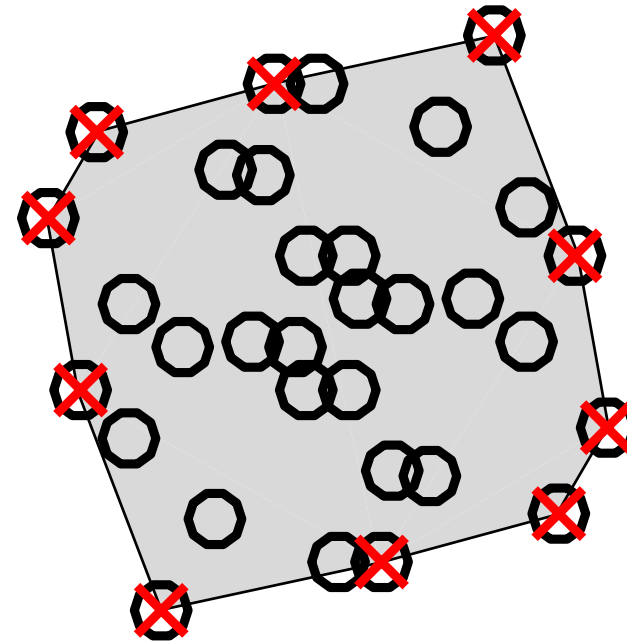
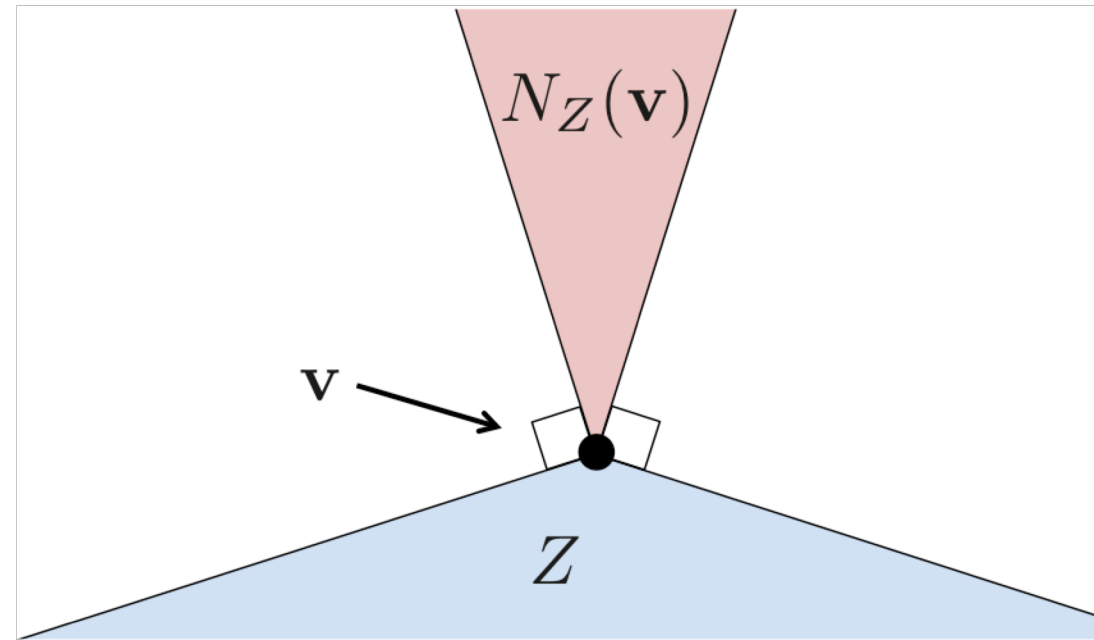
A really simple randomized algorithm to get a vertex of the zonotope

To generate a non-random vertex of a zonotope, compute

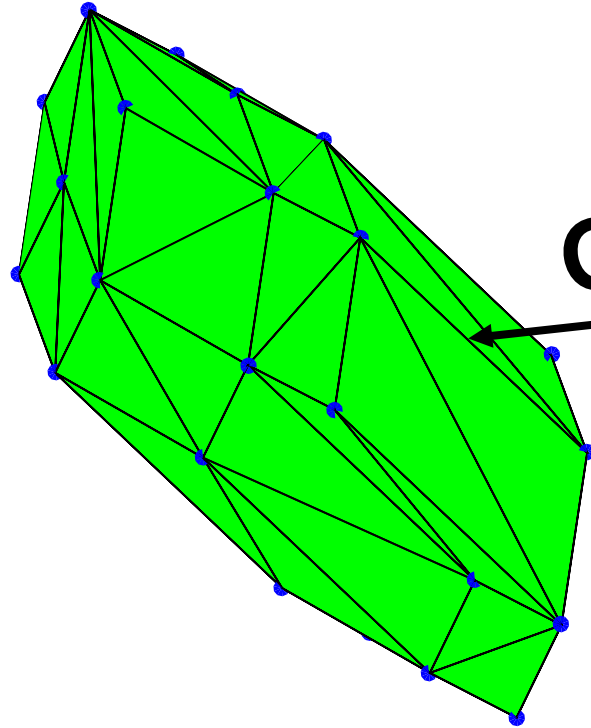
$$\mathbf{G} \text{sign}(\mathbf{G}^T \mathbf{x}) \text{ where } x_i \in N(0, 1)$$



Prob. that a vertex is generated depends on normal cone



Vertices of the signing zonotope correspond to clusterings

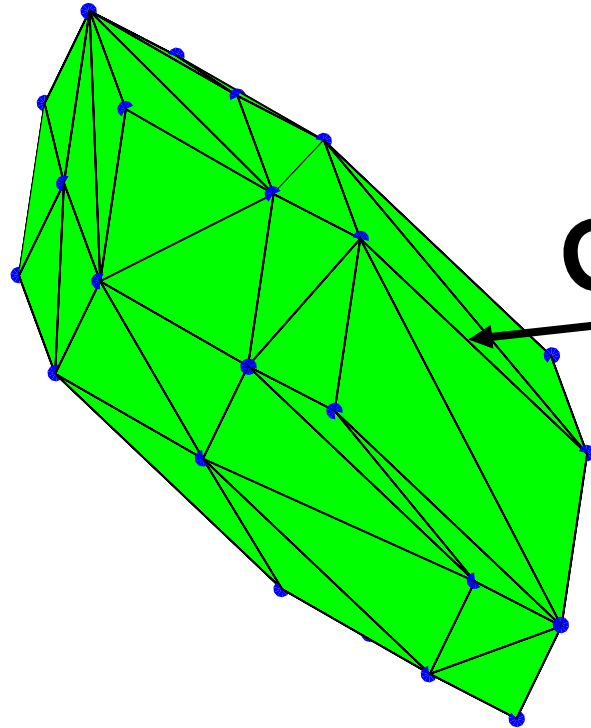


1. Begin with “signing” vector

$$\sigma = (\sigma_{r,s}^i) \in \{\pm 1\}^{n \binom{d+1}{2}}$$

2. Map signings into a zonotope

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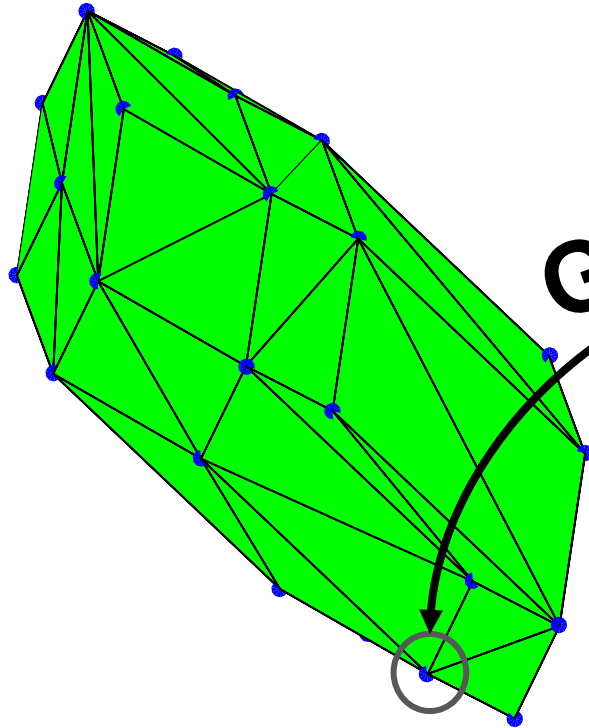
$$\mathbf{G}\sigma \quad \sigma = (\sigma_{r,s}^i) \in \{\pm 1\}^{n \binom{d+1}{2}}$$

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Columns of \mathbf{G} are vectorized outer products:

$$\mathbf{v}_i \cdot (\mathbf{e}_r - \mathbf{e}_s)^\top$$

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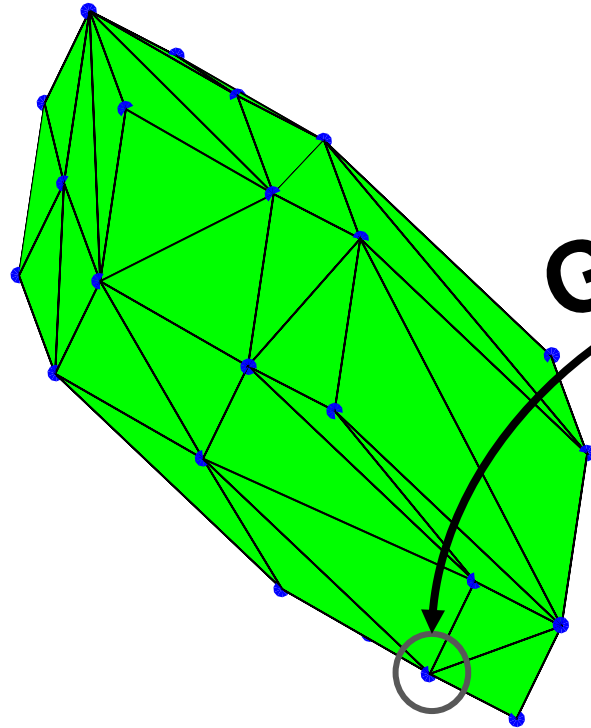
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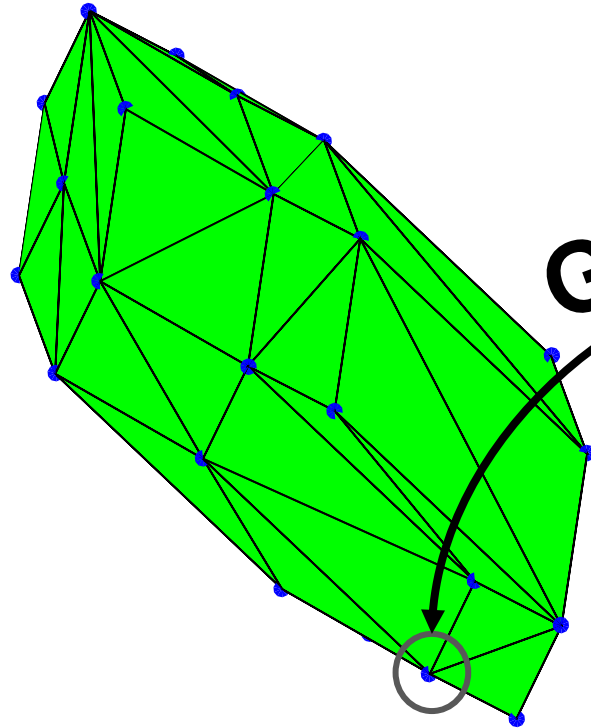
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$$\sigma_{r,s}^i = -1$$

“Node i is *not* in cluster r , but could be in cluster s ”

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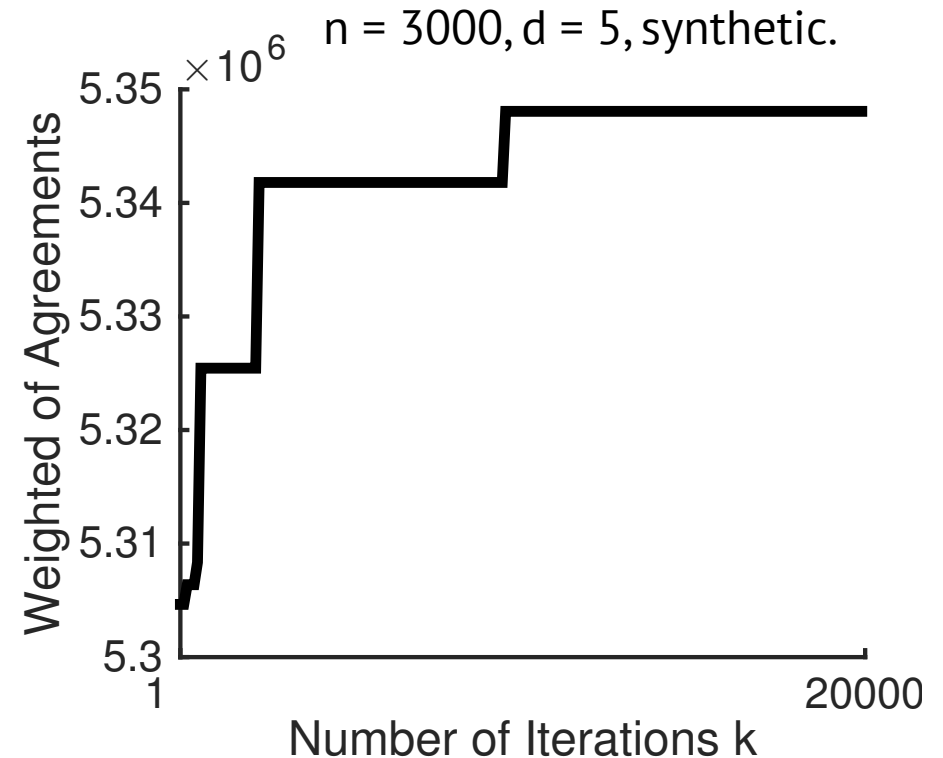
4. There are $O(n^{d^2-1})$ vertices

“Node i is *not* in cluster r , but could be in cluster s ”

In practice we just sample vertices of the zonotope

ZonoCC Algorithm $O(nk)$

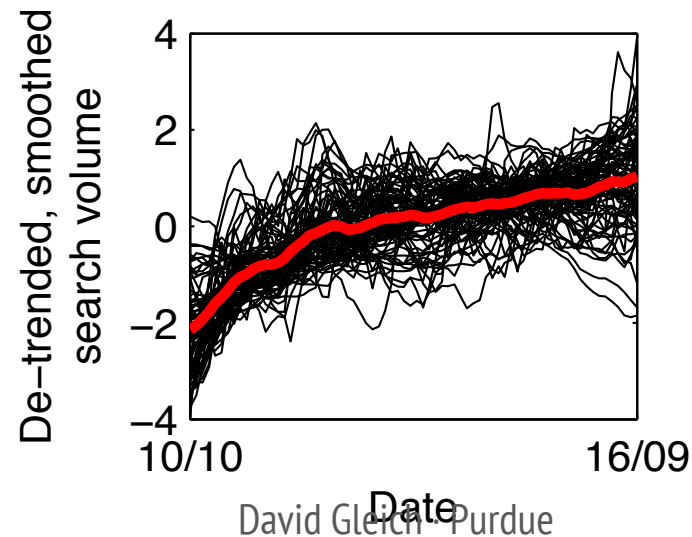
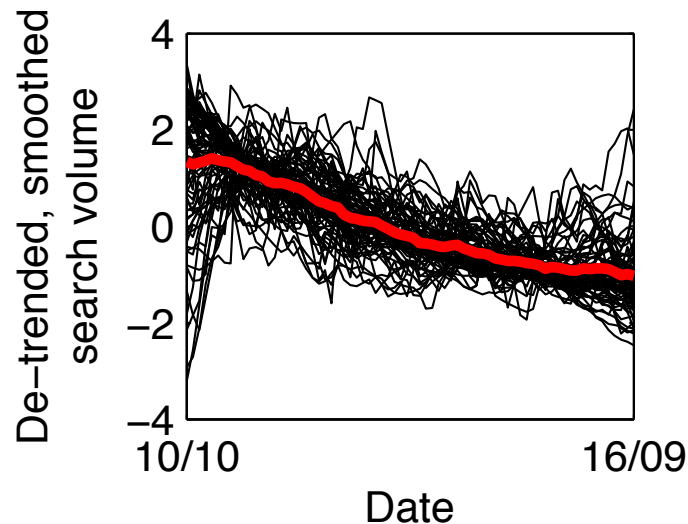
1. Generate k random extremal signings σ via random sampling
2. Find the clustering of each σ
3. Output clustering with highest objective score



For a rank-3 problem on clustering conferences

Table 2: Objective scores and runtimes in seconds for correlation clustering algorithms on two real-world datasets. Due to the size of the stocks dataset, we can run only ZONOCC and PIVOT on it.

| Dataset | | ZONOCC | PIVOT | CGW | ILP |
|-----------------------|------|--------|--------|--------|--------|
| CS Conf. $n = 157$ | Obj. | 7540.0 | 7540.0 | 7540.0 | 7540.0 |
| | Time | 7 | 1 | 1380 | 52 |
| Stocks $n = 497$ | Obj. | 5100.2 | 5099.5 | — | — |
| | Time | 40 | 20 | — | — |



Back to the question

Are there better ways of sampling these zonotopes?

$$\mathbf{G} \text{sign}(\mathbf{G}^T \mathbf{x}) \text{ where } x_i \in N(0, 1)$$

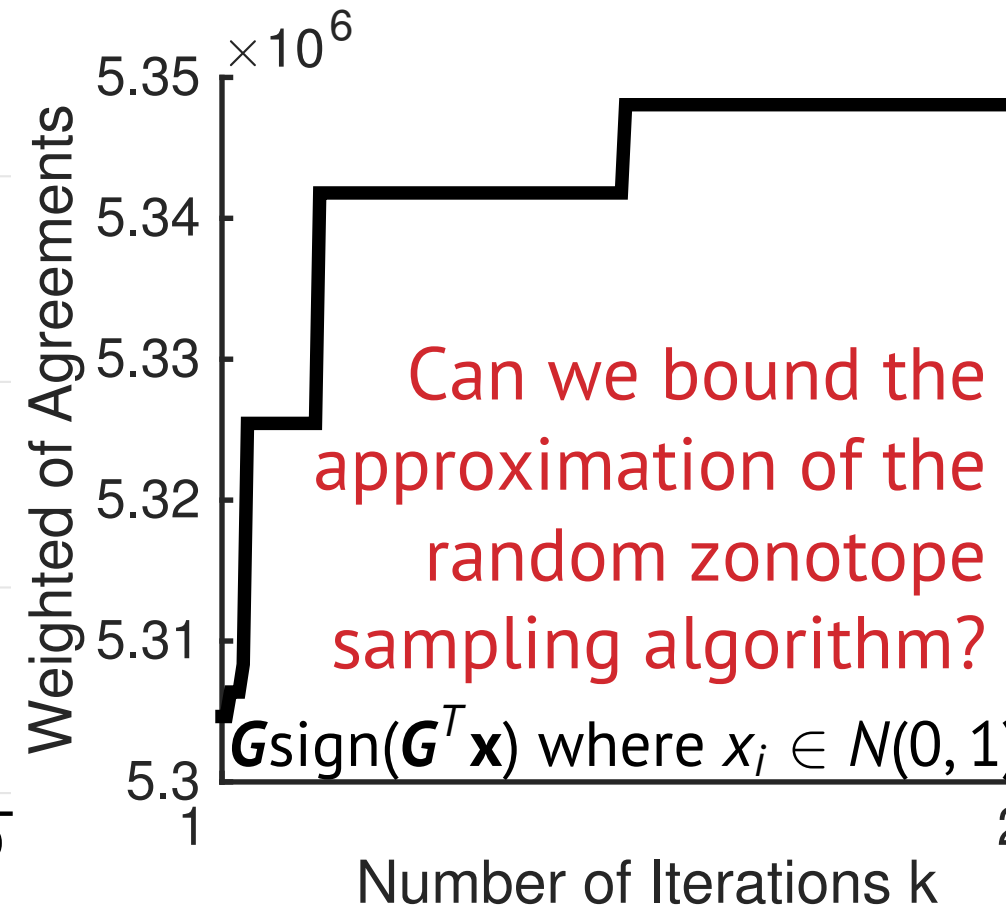
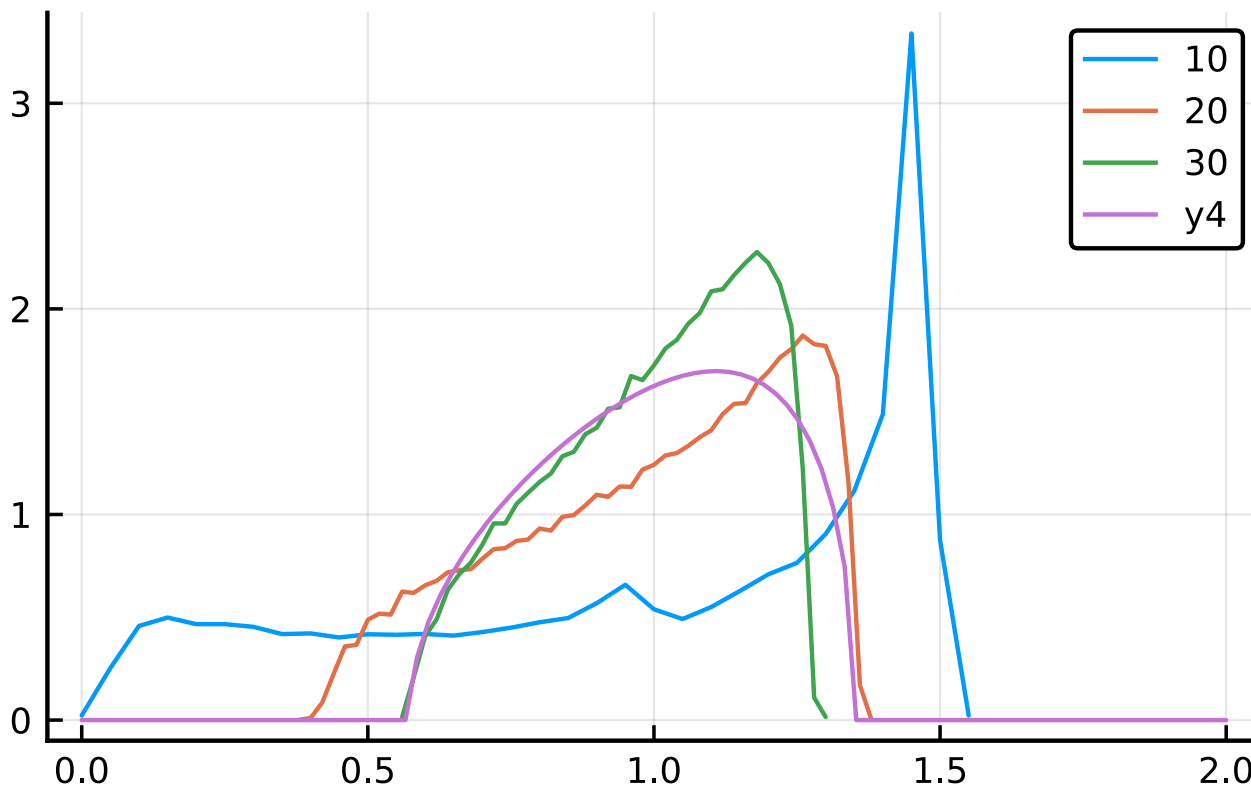
Is there anything like an approximation bound we can give for an algorithm based on sampling these in terms of the objective?

- [Ferrez, Fukuda, Libeling] reduce low-rank binary QP maximization to zonotope enumeration

$$\max \mathbf{x}^T \mathbf{A} \mathbf{x} \text{ such that } x_i \in \{0, 1\}$$

Summary. Here are my two questions!

What is this distribution??



Can we bound the approximation of the random zonotope sampling algorithm?

$\mathbf{G} \text{sign}(\mathbf{G}^T \mathbf{x})$ where $x_i \in N(0, 1)$

Code to reproduce <https://gist.github.com/dgleich/4d4becc858e4a7d7952af6c66c99e7b9>