RelationalAI

Logic & Algorithms in DB & AI – 2023

Worst-Case Optimal Joins

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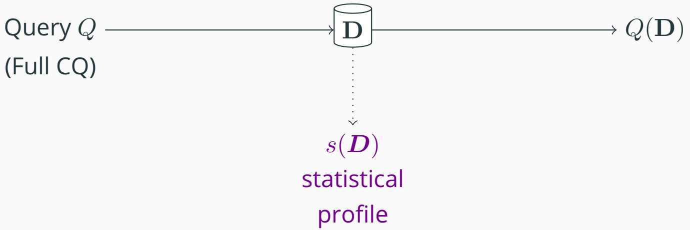
Outline

Worst-Case Optimal Join (WCOJ) Algorithm



Definition

A "worst-case optimal" join algorithm is an algorithm computing $Q(\boldsymbol{D})$ in time

$$\tilde{O}\left(|\boldsymbol{D}| + \sup_{\boldsymbol{D}' \models s(\boldsymbol{D})} |Q(\boldsymbol{D}')|\right)$$

$\tilde{O}$ hides log and query-dependent factors

Hierarchy of Join Algorithms

JAAT	Join at a time
VAAT	Variable at a time
IAAT	Inequality at a time

The Algorithm is in the Pudding

Proof $\Longrightarrow$ Algorithm!

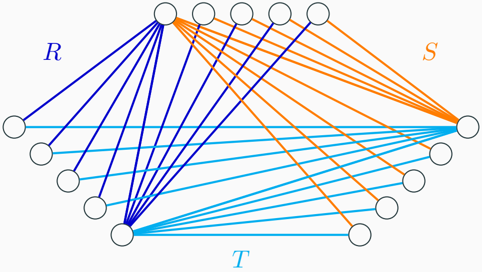
Hierarchy of Join Algorithms

$$
\begin{aligned}
\text{(not achievable ?)} \quad & \sup_{\boldsymbol{D}' \models s(\boldsymbol{D})} |Q(\boldsymbol{D}')| \\
\text{(not achievable ?)} \quad & \leq \text{entropic-bound}(Q, s) \\
\text{(IAAT)} \quad & \leq \text{polymatroid-bound}(Q, s) && \text{PANDA} \\
\text{(VAAT)} \quad & \leq \text{chain-bound}(Q, s, \sigma) && \text{NPRR, LFTJ, GJ} \\
\text{(VAAT)} \quad & \leq \text{agm-bound}(Q, s) && \text{NPRR, LFTJ, GJ} \\
\text{(JAAT)} \quad & \leq \text{integral-edge-cover}(Q, s) && \text{Binary Plans}
\end{aligned}
$$

Runtime = $\tilde{O}(2^{\text{bound}})$

Outline

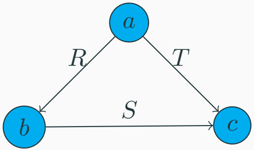
Example: JAAT Query Plans Are Sub-Optimal

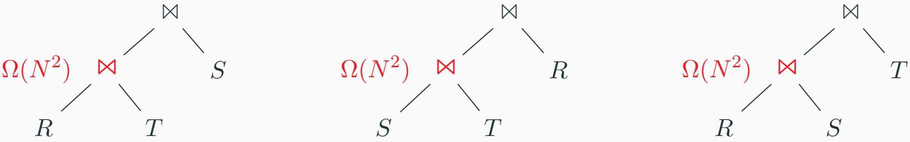


$$|R| = |S| = |T| = 2N - 1$$

$$Q_\Delta(a,b,c) = R(a,b) \wedge S(b,c) \wedge T(a,c)$$

$$\sup_{\mathbf{D}' \models \mathsf{DC}(\boldsymbol{D})} |Q_\Delta(\mathbf{D}')| = O(N^{1.5})$$





Outline

Triangle Query

$$Q_\triangle(A,B,C) \leftarrow R(A,B), S(B,C), T(A,C)$$

AGM-bound for Q: (E.g. num triangles in a graph $\leq |E|^{3/2}$)

$$|Q_\triangle| \leq |R|^{\lambda_R} \cdot |S|^{\lambda_S} \cdot |T|^{\lambda_T}$$

whenever $\boldsymbol{\lambda} = (\lambda_R, \lambda_S, \lambda_T)$ is a fractional edge cover for the triangle:

$$\lambda_R + \lambda_S \geq 1$$
$$\lambda_R + \lambda_T \geq 1 \qquad \boldsymbol{\lambda} \geq \mathbf{0}$$
$$\lambda_S + \lambda_T \geq 1$$

Pick $\boldsymbol{\lambda}$ to minimize the bound.

Triangle Query: AGM Bound from Hölder Inequality

Consider a "section" of this query on a given value $a \in \text{Dom}(A)$:

$$Q_\triangle(a, B, C) \leftarrow R(a, B), S(B, C), T(a, C)$$

Need (b, c) in the intersection $S \cap (\sigma_{A=a}R \times \sigma_{A=a}T)$, thus

$$\begin{aligned} |\sigma_{A=a}Q_\triangle| &\leq \min\{|S|, |\sigma_{A=a}R| \cdot |\sigma_{A=a}T|\} \\ &\leq |S|^{\lambda_S} \cdot (|\sigma_{A=a}R| \cdot |\sigma_{A=a}T|)^{1-\lambda_S} \\ &\leq |S|^{\lambda_S} \cdot |\sigma_{A=a}R|^{\lambda_R} \cdot |\sigma_{A=a}T|^{\lambda_T} \end{aligned}$$

Triangle Query: AGM Bound Based on Hölder Inequality

Iterate over all possible values of a:

$$
\begin{aligned}
|Q_\triangle| &= \sum_a |\sigma_{A=a} Q_\triangle| \\
&\le \sum_a |S|^{\lambda_S} \cdot |\sigma_{A=a} R|^{\lambda_R} \cdot |\sigma_{A=a} T|^{\lambda_T} \\
&= |S|^{\lambda_S} \cdot \sum_a |\sigma_{A=a} R|^{\lambda_R} \cdot |\sigma_{A=a} T|^{\lambda_T} \\
\text{(Hölder)} &\le |S|^{\lambda_S} \cdot \Big(\sum_a |\sigma_{A=a} R|\Big)^{\lambda_R} \cdot \Big(\sum_a |\sigma_{A=a} T|\Big)^{\lambda_T} \\
&= |S|^{\lambda_S} \cdot |R|^{\lambda_R} \cdot |T|^{\lambda_T}
\end{aligned}
$$

$$Q_\triangle(A,B,C) \leftarrow R(A,B), S(B,C), T(A,C)$$

Algorithm 1: based on Hölder's inequality proof

for $a \in \pi_A R \cap \pi_A T$ **do**
  for $b \in \pi_B \sigma_{A=a} R \cap \pi_B S$ **do**
    for $c \in \pi_C \sigma_{B=b} S \cap \pi_C \sigma_{A=a} T$ **do**
      Report (a,b,c);

In English

- For each $a \in \pi_A R \cap \pi_A T$, enumerate $(a,b,c) \in \sigma_{A=a} Q_\triangle$

Triangle Query: VAAT is in the Pudding

Computing this "section" of the query on a given value $a \in \text{Dom}(A)$

$$Q_\triangle(a, B, C) \leftarrow R(a, B), S(B, C), T(a, C)$$

is to compute the intersection $S \cap (\sigma_{A=a}R \times \sigma_{A=a}T)$, which can be done in time

$$\tilde{O}\left(\min\{|S|, |\sigma_{A=a}R| \cdot |\sigma_{A=a}T|\}\right) \leq \tilde{O}\left(|\sigma_{A=a}R|^{\lambda_R} \cdot |\sigma_{A=a}T|^{\lambda_T} \cdot |S|^{\lambda_S}\right)$$

Overall, the algorithm runs in time (Modulo $\tilde{O}(N)$ pre-processing)

$$\tilde{O}\left(\sum_a |\sigma_{A=a}R|^{\lambda_R} \cdot |\sigma_{A=a}T|^{\lambda_T} \cdot |S|^{\lambda_S}\right) = \tilde{O}\left(|S|^{\lambda_S} \cdot |R|^{\lambda_R} \cdot |T|^{\lambda_T}\right)$$

Full Conjunctive Query

$$Q(V) \leftarrow \bigwedge_{S \in \mathcal{E}} R_S(S)$$

Query hypergraph $\mathcal{H} = (V, \mathcal{E})$

AGM-bound for Q:

Assuming only cardinality constraints $(\emptyset, S, |R_S|)$

$$|Q| \leq \prod_{S \in \mathcal{E}} |R_S|^{\lambda_S}$$

whenever $\boldsymbol{\lambda} = (\lambda_S)_{S \in \mathcal{E}}$ is a fractional edge cover for $\mathcal{H}$:

$$\forall v \in V \ : \ \sum_{S \in \mathcal{E}, v \in S} \lambda_S \geq 1 \qquad \boldsymbol{\lambda} \geq \mathbf{0}.$$

Pick $\boldsymbol{\lambda}$ to minimize the bound.

Full Conjunctive Query: AGM Bound from Hölder Inequality

Consider a "section" of this query on a given value $a \in \text{Dom}(A)$:

$$\sigma_{A=a} Q(V) \leftarrow \bigwedge_{S \in \mathcal{E}, A \notin S} R_S(S) \wedge \bigwedge_{S \in \mathcal{E}, A \in S} \sigma_{A=a} R_S(S)$$

Now iterate over all possible values of a:

$$|Q| = \sum_a |\sigma_{A=a} Q_\triangle| \leq \sum_a \prod_{S \in \mathcal{E}, A \notin S} |R_S|^{\lambda_S} \cdot \prod_{S \in \mathcal{E}, A \in S} |\sigma_{A=a} R_S|^{\lambda_S}$$

$$\text{(Hölder's inequality)} \leq \prod_{S \in \mathcal{E}, A \notin S} |R_S|^{\lambda_S} \cdot \prod_{S \in \mathcal{E}, A \in S} \left(\sum_a |\sigma_{A=a} R_S| \right)^{\lambda_S}$$

$$= \prod_{S \in \mathcal{E}} |R_S|^{\lambda_S}.$$

Full Conjunctive Query: VAAT Algorithm

[NPRR 12] [Veldhuizen 14]

$$Q(V) \leftarrow \bigwedge_{S \in \mathcal{E}} R_S(S)$$

Algorithm 2: based on Hölder's inequality proof

for $a \in \bigcap_{S \in \mathcal{E}, A \in S} \pi_A R_S$ **do**

  Recursively solve the query section $\sigma_{A=a} Q$ (on variables $V - \{A\}$);

  Report $\{a\} \times \sigma_{A=a} Q$

Runtime $\tilde{O}\left(\sum_{S \in \mathcal{E}} |R_S|\right) + \tilde{O}\left(\prod_{S \in \mathcal{E}} |R_S|^{\lambda_S}\right)$.

Proof: straightforward.

Full Conjunctive Query: Chain-Bound, VAAT Algorithm

For degree constraints *beyond* cardinality constraints

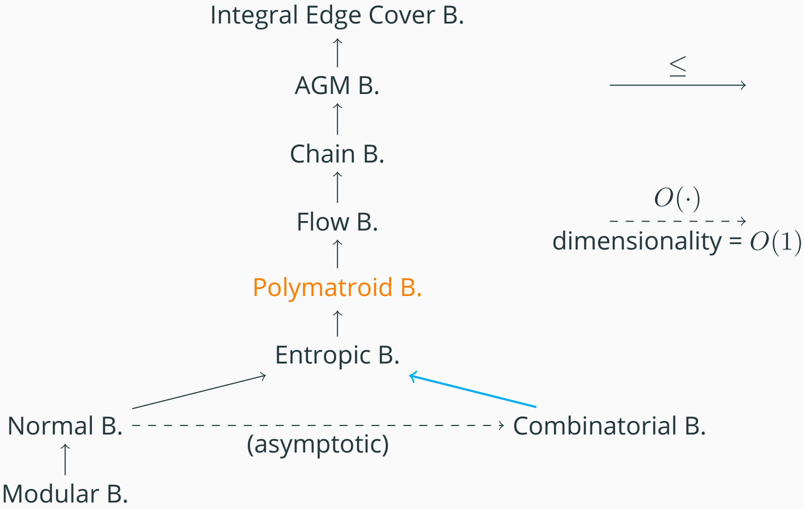
- The AGM bound does not apply.
- We use the chain-bound instead.

Find a variable ordering σ

- Arbitrary degree constraint set DC
- Runtime predicted by $\text{chain-bound}(\text{DC}, \sigma)$ Tight for acyclic DC
- VAAT algorithm meeting the chain-bound; similar analysis

Outline

Recall: Bound Hierarchy



Recall: The Entropic and Polymatroid Bounds

Γ_n is the set of n-dim polymatroids. $h(Y|X) = h(Y) - h(X), X \subseteq Y$.

Theorem

If $s(\boldsymbol{D})$ contains only degree constraints, then

$$\log \sup_{\boldsymbol{D}' \models s(\boldsymbol{D})} |Q(\boldsymbol{D}')| \leq \sup_{h \in \overline{\Gamma}^*_n \cap DC} h(V) \leq \max_{h \in \Gamma_n \cap DC} h(V)$$

where DC is the set of linear constraints of the form

$$h(Y|X) \leq \log N$$

for each degree constraint (X, Y, N).

More explicitly,

$$
\begin{array}{llll}
\max & h(V) & & \text{dual vars} \\
\text{s.t.} & h(Y) - h(X) \leq \log N, & (X, Y, N) \in \text{DC} & \delta_{Y|X} \\
& h(I \cup J|J) - h(I|I \cap J) \leq 0, & I \perp J & \sigma_{I,J} \\
& h(X) - h(Y) \leq 0, & \emptyset \neq X \subset Y \subseteq V & \mu_{Y|X} \\
& h(Z) \geq 0, & \emptyset \neq Z \subseteq V. &
\end{array}
$$

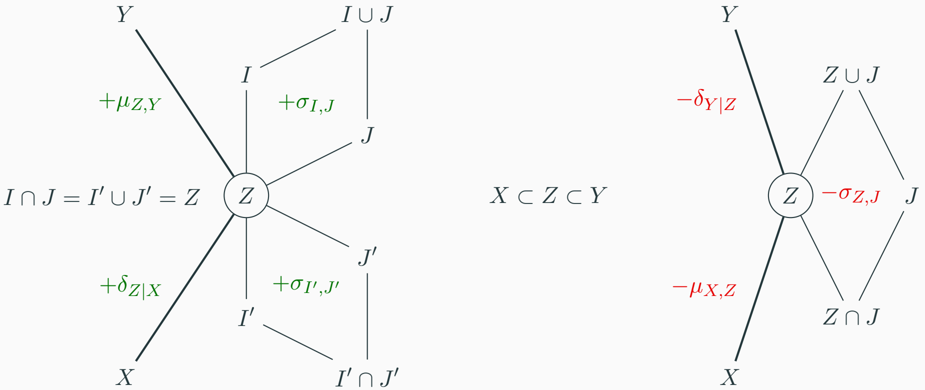
$I \perp J$ means $I \not\subseteq J$ and $J \not\subseteq I$.

$$
\begin{aligned}
\min \quad & \sum_{(X,Y,N)\in \mathsf{DC}} \log N \cdot \delta_{Y|X} \\
\text{s.t.} \quad & \mathsf{excess}(V) \geq 1, \\
& \mathsf{excess}(Z) \geq 0, \qquad \emptyset \neq Z \subseteq V. \\
& (\boldsymbol{\delta}, \boldsymbol{\sigma}, \boldsymbol{\mu}) \geq \mathbf{0}.
\end{aligned}
$$

where, for any $\emptyset \neq Z \in 2^V$, the quantity $\mathsf{excess}(Z)$ is defined by

$$
\begin{aligned}
\mathsf{excess}(Z) := & \sum_{X:(X,Z)\in \mathsf{DC}} \delta_{Z|X} - \sum_{Y:(Z,Y)\in \mathsf{DC}} \delta_{Y|Z} + \sum_{\substack{I \perp J \\ I \cap J = Z}} \sigma_{I,J} \\
& + \sum_{\substack{I \perp J \\ I \cup J = Z}} \sigma_{I,J} - \sum_{J:J \perp Z} \sigma_{Z,J} - \sum_{X:X \subset Z} \mu_{X,Z} + \sum_{Y:Z \subset Y} \mu_{Z,Y}.
\end{aligned}
$$

Contributions of coefficients to $\text{excess}(Z)$



Definition

Given $\boldsymbol{\delta} \geq 0$, the following is a Shannon-flow inequality if it holds for all $h \in \Gamma_n$:

$$h(V) \leq \sum_{(X,Y,N) \in \mathsf{DC}} \delta_{Y|X} \cdot (h(Y) - h(X))$$

- $\boldsymbol{\delta}$ defines a Shannon-flow inequality iff $\exists(\boldsymbol{\sigma}, \boldsymbol{\mu})$ s.t. $(\boldsymbol{\delta}, \boldsymbol{\sigma}, \boldsymbol{\mu})$ is dual-feasible.
- If DC contains only cardinality constraints $(X = \emptyset, Y, N)$, then $\boldsymbol{\delta}$ defines a Shannon-flow inequality iff it is a fractional edge cover of the query hypergraph.

Shearer's Lemma!

Example: Shannon-Flow Inequality for Triangle Query

$$h(A,B,C) \leq \frac{1}{2}(h(A,B)+h(B,C)+h(A,C))$$

A step-by-step proof: [Radhakrishnan 2003]

$$\begin{aligned}
& h(A,B)+h(A,C)+h(B,C) \\
\text{(decomposition)} &= h(A)+h(B|A)+h(B,C)+h(A,C) \\
\text{(sub-modularity)} &\geq (h(A|B,C)+h(B,C))+(h(B|A)+h(A,C)) \\
\text{(composition)} &= h(A,B,C)+(h(B|A)+h(A,C)) \\
\text{(sub-modularity)} &\geq h(A,B,C)+(h(B|A,C)+h(A,C)) \\
\text{(composition)} &= h(A,B,C)+h(A,B,C)
\end{aligned}$$

Example: Another Shannon-Flow Inequality

$$h(ABCD) \leq \frac{1}{2}[h(AB) + h(BC) + h(CD) + h(D|AC) + h(A|BD)],$$

$$\begin{aligned}
& h(AB) + h(BC) + h(CD) + h(D|AC) + h(A|BD) \\
\text{(decomposition)} &= h(AB) + h(B) + h(C|B) + h(CD) + h(D|AC) + h(A|BD) \\
\text{(sub-modularity)} &\geq h(AB) + h(B) + h(C|B) + h(CD|B) + h(D|AC) + h(A|BD) \\
\text{(composition)} &= h(AB) + h(C|B) + h(BCD) + h(D|AC) + h(A|BD) \\
\text{(sub-modularity)} &\geq h(AB) + h(C|B) + h(BCD) + h(D|AC) + h(A|BCD) \\
\text{(composition)} &= h(AB) + h(C|B) + h(D|AC) + h(ABCD) \\
\text{(sub-modularity)} &\geq h(AB) + h(C|AB) + h(D|AC) + h(ABCD) \\
\text{(composition)} &= h(ABC) + h(D|AC) + h(ABCD) \\
\text{(sub-modularity)} &\geq h(ABC) + h(D|ABC) + h(ABCD) \\
\text{(composition)} &= h(ABCD) + h(ABCD).
\end{aligned}$$

From LP-duality, there exists $\boldsymbol{\delta} \geq \mathbf{0}$ s.t.

$$\text{polymatroid-bound} := \max\{h(V) \mid h \in C \cap \Gamma_n\} = \sum_{(X,Y,N) \in \text{DC}} \delta_{Y|X} \log N,$$

and for these $\boldsymbol{\delta}$, from Farkas's lemma we have

$$h(V) \leq \sum_{(X,Y,N) \in \text{DC}} \delta_{Y|X} \cdot h(Y|X), \quad \forall h \in \Gamma_n$$

Proof Sequence for a Shannon-Flow Inequality

$$h(V) \leq \sum_{(X,Y,N)\in \text{DC}} \delta_{Y|X} \cdot h(Y|X)$$

A proof sequence is a conversion from RHS to LHS using a sequence of steps of the form

(In)equality	Steps ($X \subseteq Y$)
$h(X) + h(Y\|X) = h(Y)$	$h(X) + h(Y\|X) \to h(Y)$
$h(Y) = h(X) + h(Y\|X)$	$h(Y) \to h(X) + h(Y\|X)$
$h(Y) \geq h(X)$	$h(Y) \to h(X)$
$h(Y\|X) \geq h(Y \cup Z\|X \cup Z)$	$h(Y\|X) \to h(Y \cup Z\|X \cup Z)$

Existence of Proof Sequence

Lemma (ANS 2017)

There is a proof sequence for every Shannon-flow inequality. (The length is at most doubly exponential in $|V|$).

The Shannon-flow inequality is a linear combination of dual constraints; the proof sequence is more stringent than that.

Outline

Overview

JAAT Algorithm

VAAT Algorithm

Shannon-Flow Inequalities

IAAT Algorithm

Open Problems – Research Directions

One Inequality At A Time (IAAT)

There is an algorithm (called PANDA) that converts a proof sequence $\rightarrow$ an efficient algorithm to answer the original query

Steps ($X \subseteq Y$)	Relational Operator
$h(X) + h(Y\|X) \rightarrow h(Y)$	(join)
$h(Y) \rightarrow h(X) + h(Y\|X)$	(data partition)
$h(Y) \rightarrow h(X)$	(projection)
$h(Y\|X) \rightarrow h(Y \cup Z\|X \cup Z)$	(NOP)

Theorem

PANDA solves any conjunctive query Q in time $\tilde{O}(N + \text{poly}(\log N) \cdot 2^{\textit{polymatroid bound}})$

Example: PANDA for Triangle Query

$$Q(A,B,C) \leftarrow R(A,B), S(B,C), T(A,C)$$

$$R^{\text{heavy}}(A,B) = \{(a,b) \ : \ |\sigma_{A=a}R| > \sqrt{N}\}$$

$$R^{\text{light}}(A,B) = \{(a,b) \ : \ |\sigma_{A=a}R| \leq \sqrt{N}\}$$

Algorithm is in the pudding!

$$\begin{aligned}
h(A,B) + h(A,C) + h(B,C) & \quad R(A,B), S(B,C), T(A,C) \\
= h(A) + h(B|A) + h(B,C) + h(A,C) & \quad R^{\text{heavy}}(A,B), R^{\text{light}}(A,B), S(B,C), T(A,C) \\
\geq (h(A|B,C) + h(B,C)) + (h(B|A) + h(A,C)) & \quad R^{\text{heavy}}(A,B), R^{\text{light}}(A,B), S(B,C), T(A,C) \\
= h(A,B,C) + (h(B|A) + h(A,C)) & \quad I^{\text{heavy}}(A,B,C), R^{\text{light}}(A,B), T(A,C) \\
\geq h(A,B,C) + (h(B|A,C) + h(A,C)) & \quad I^{\text{heavy}}(A,B,C), R^{\text{light}}(A,B), T(A,C) \\
= h(A,B,C) + h(A,B,C) & \quad I^{\text{heavy}}(A,B,C), I^{\text{light}}(A,B,C).
\end{aligned}$$

Example: PANDA for Triangle Query

The real query plan:

$$
\begin{aligned}
& R(A, B) \wedge S(B, C) \wedge T(A, C) \\
& = (R^{\text{heavy}}(A, B) \vee R^{\text{light}}(A, B)) \wedge S(B, C) \wedge T(A, C) \\
& = (R^{\text{heavy}}(A, B) \wedge S(B, C) \wedge T(A, C)) \vee (R^{\text{light}}(A, B) \wedge S(B, C) \wedge T(A, C)) \\
& = (R^{\text{heavy}}(A, B) \wedge S(B, C) \wedge T(A, C)) \vee (R^{\text{light}}(A, B) \wedge S(B, C) \wedge T(A, C)) \\
& = I^{\text{heavy}}(A, B, C) \wedge T(A, C) \vee I^{\text{light}}(A, B, C) \wedge S(B, C).
\end{aligned}
$$

- Note that $|I^{\text{heavy}}(A, B, C)| \leq N^{3/2}$ and $|I^{\text{light}}(A, B, C)| \leq N^{3/2}$.
- Overall runtime is $\tilde{O}(N^{3/2})$.

Example: PANDA for a More Interesting Example [ANS 17]

$$Q(A,B,C,D) \leftarrow R(A,B) \wedge S(B,C) \wedge T(C,D) \wedge \text{hash}(A,C) = D \wedge \text{hash}(B,D) = A$$

From the Shannon-flow inequality:

$$h(ABCD) \leq \frac{1}{2}[h(AB) + h(BC) + h(CD) + h(D|AC) + h(A|BD)],$$

we know

$$\log_2 |Q| \leq \frac{1}{2}[\log_2 |R| + \log_2 |S| + \log_2 |T| + 0 + 0]$$

or

$$|Q| \leq \sqrt{|R||S||T|}$$

Example: Another Shannon-Flow Inequality

$$h(ABCD) \leq \frac{1}{2}[h(AB) + h(BC) + h(CD) + h(D|AC) + h(A|BD)],$$

$$\begin{aligned}
& h(AB) + h(BC) + h(CD) + h(D|AC) + h(A|BD) \\
\text{(decomposition)} &= h(AB) + h(B) + h(C|B) + h(CD) + h(D|AC) + h(A|BD) \\
\text{(sub-modularity)} &\geq h(AB) + h(B) + h(C|B) + h(CD|B) + h(D|AC) + h(A|BD) \\
\text{(composition)} &= h(AB) + h(C|B) + h(BCD) + h(D|AC) + h(A|BD) \\
\text{(sub-modularity)} &\geq h(AB) + h(C|B) + h(BCD) + h(D|AC) + h(A|BCD) \\
\text{(composition)} &= h(AB) + h(C|B) + h(D|AC) + h(ABCD) \\
\text{(sub-modularity)} &\geq h(AB) + h(C|AB) + h(D|AC) + h(ABCD) \\
\text{(composition)} &= h(ABC) + h(D|AC) + h(ABCD) \\
\text{(sub-modularity)} &\geq h(ABC) + h(D|ABC) + h(ABCD) \\
\text{(composition)} &= h(ABCD) + h(ABCD).
\end{aligned}$$

Example : PANDA for a More Interesting Example

$$
\begin{aligned}
& R(A,B) \wedge S(B,C) \wedge T(C,D) \wedge \text{hash}(A,C) = D \wedge \text{hash}(B,D) = A \\
&= R(A,B) \wedge S^{\text{heavy}}(B,C) \wedge T(C,D) \wedge \text{hash}(A,C) = D \wedge \text{hash}(B,D) = A \\
&\vee R(A,B) \wedge S^{\text{light}}(B,C) \wedge T(C,D) \wedge \text{hash}(A,C) = D \wedge \text{hash}(B,D) = A \\
&= R(A,B) \wedge S^{\text{heavy}}(B,C) \wedge T(C,D) \wedge \text{hash}(A,C) = D \wedge \text{hash}(B,D) = A \\
&\vee R(A,B) \wedge S^{\text{light}}(B,C) \wedge T(C,D) \wedge \text{hash}(A,C) = D \wedge \text{hash}(B,D) = A \\
&= R(A,B) \wedge I^{\text{heavy}}(B,C,D) \wedge \text{hash}(A,C) = D \wedge \text{hash}(B,D) = A \\
&\vee I^{\text{light}}(A,B,C) \wedge T(C,D) \wedge \text{hash}(A,C) = D \wedge \text{hash}(B,D) = A \\
&= R(A,B) \wedge I^{\text{heavy}}(B,C,D) \wedge \text{hash}(A,C) = D \wedge \text{hash}(B,D) = A \\
&\vee I^{\text{light}}(A,B,C) \wedge T(C,D) \wedge \text{hash}(A,C) = D \wedge \text{hash}(B,D) = A \\
&= R(A,B) \wedge J^{\text{heavy}}(A,B,C,D) \wedge \text{hash}(A,C) = D \\
&\vee J^{\text{light}}(A,B,C,D) \wedge T(C,D) \wedge \text{hash}(B,D) = A.
\end{aligned}
$$

Example : PANDA for a More Interesting Example

Main question

How to define S^{heavy} and S^{light} so that runtime is $\tilde{O}(2^{h^*(A,B,C,D)})$

$$S^{\text{heavy}}(B,C) = \{(b,c) \ : \ |\sigma_{C=c}S| > 2^{h^*(B,C)-h^*(C)}\}$$
$$S^{\text{light}}(B,C) = \{(b,c) \ : \ |\sigma_{C=c}S| \leq 2^{h^*(B,C)-h^*(C)}\}$$

Assuming h^* and $(\boldsymbol{\delta}^*, \boldsymbol{\sigma}^*, \boldsymbol{\mu}^*)$ are primal-dual optimal: ($\delta^*_{CD|\emptyset} > 0$ and $\delta^*_{AB|\emptyset} > 0$)

$$|S^{\text{light}}(B,C) \wedge T(C,D)| \leq 2^{h^*(B,C)-h^*(C)} \cdot 2^{h^*(C,D)} = 2^{h^*(B,C,D)} \leq 2^{h^*(A,B,C,D)}$$
$$|S^{\text{heavy}}(B,C) \wedge R(A,B)| \leq 2^{h^*(C)} \cdot 2^{h^*(A,B)} = 2^{h^*(A,B,C)} \leq 2^{h^*(A,B,C,D)}.$$

$=$ holds because SFI holds with $=$ for h^*.

The Actual PANDA Algorithm (Needs Another Talk)

More complicated because:

- Couldn't prove that every heavy / light copy reaches $h(V)$ eventually.
- Couldn't prove that in the proof sequence we won't ever compose terms which were decomposed in an earlier step

Main ideas to push through:

- A decomposition $h(A, B) \to h(A) + h(B|A)$ corresponds to partitioning R into logarithmically many "uniform" parts.
- Essentially, each each part, both the heavy condition and the light condition are satisfied.
- Induct on logarithmically many subproblems, including constructing a new proof sequence for each of them

The Actual PANDA Algorithm

PANDA runs in Time

$$\tilde{O}(N + \text{poly}(\log N) \cdot 2^{\text{polymatroid bound for } Q}) = \tilde{O}(N + \text{poly}(\log N) \cdot \sup_{\boldsymbol{D'} \models s(\boldsymbol{D})} |Q(\boldsymbol{D'})|)$$

Outline

Open Problems

- Bound the minimum proof sequence length
  - Educated conjecture: polynomial($|DC|$)
- Remove polylog factor from runtime of PANDA
- Is there a more natural algorithm?
- Exploiting conditional independence

Main References

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Many Thanks!